



Network Calculus Based Approach to DDoS Detection and Mitigation in CPS

William Emfinger

04-11-2015



Outline

- * Introduction
- * Network Calculus fundamentals
- * Network Performance Analysis for CPS Applications : PNP²
- * PNP² DDoS Detection
- * DDoS Mitigation
- * Conclusions

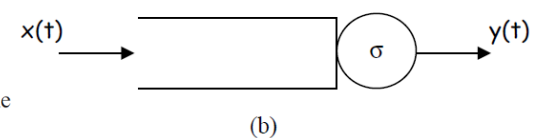
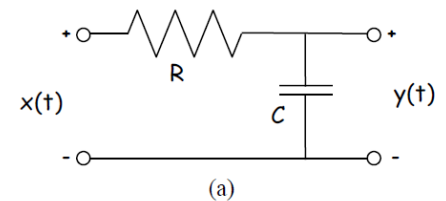
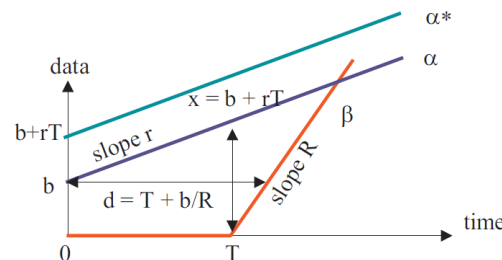
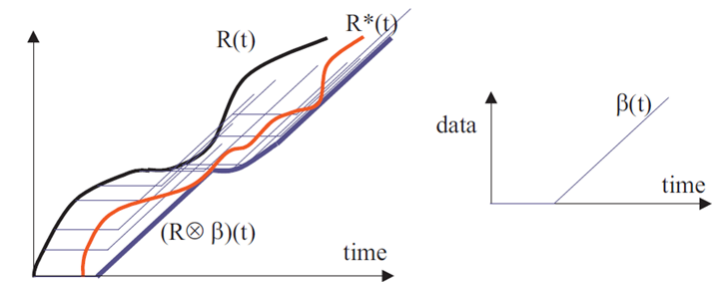
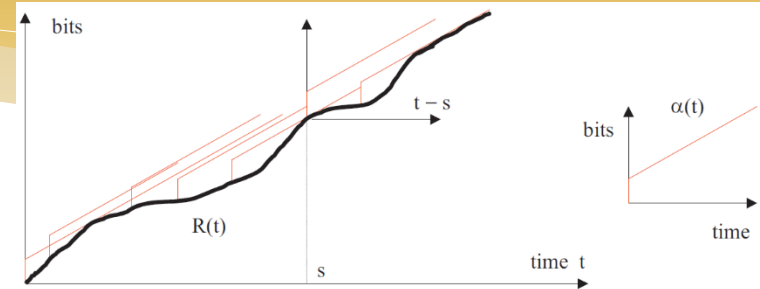
Introduction

- * Cyber-Physical Systems (CPS):
 - * According to NIST: *CPS are smart systems that include co-engineered interacting networks of physical and computational components*
 - * Embedded computers, sensing/actuation software, tightly coupled physical system
 - * Communications network (wireless) acts as backbone for cooperation between nodes
 - * Network resources **provided as a service** to users
- * Network resource availability affected by both application load and system's environment
 - * Bandwidth (bits/sec) required
 - * Buffering delay and transmission delay incurred
 - * Availability of network resources
 - * Some systems may have **deterministic, time-varying network characteristics** and/or application **network load**
- * Design-time analysis is critical for applications that require tight and/or real-time guarantees
 - * Use the analysis to perform **test for admission** into the system



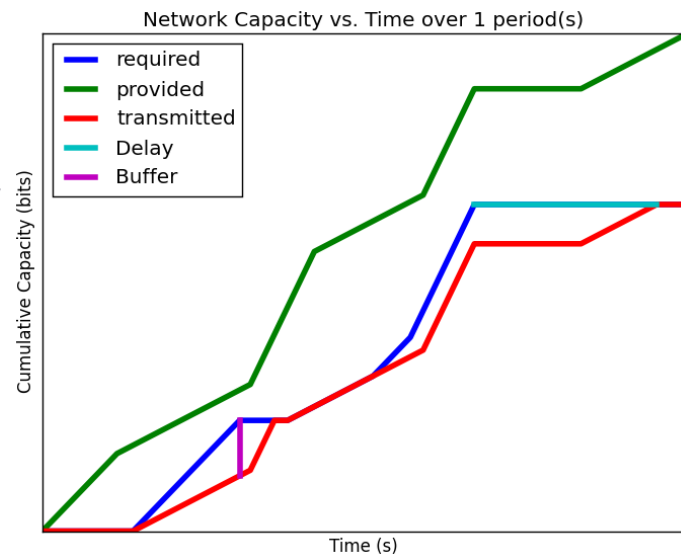
Network Calculus Fundamentals

- * Based on $(\min,+)$ -calculus
- * Abstracts network traffic and computing nodes as *arrival curves* and traffic shaping *service curves*
- * Analytically predicts worst-case buffer requirements and application buffering delay
- * Applies system-theory concepts to computer network analysis
 - * Input function (data input flow) convolved with a system function (shaper) to produce the output function (data output flow)
- * Allows the calculation of deterministic performance bounds on
 - * Buffering delay
 - * Buffer backlog



Network Performance Analysis for CPS Applications

- * To address the issues described above, we have developed a network modeling paradigm similar to Network Calculus' arrival curves and service curves, called **Precise Network Performance Prediction (PNP²)**
- * In contrast to Network Calculus, we precisely model the network traffic as a function of time, i.e. bits produced or serviced as functions of time – called *profiles*
- * Convolve the application data profile with the network capacity profile to obtain transmitted data profile
 - * $r[t]$: application (*required*) data profile
 - * $p[t]$: network (*provided*) capacity profile
 - * $o[t]$: the transmitted (*output*) data profile
- * Assumptions:
 - * Know at design-time how the application and the network is expected perform
 - * Time-Synchronization between the nodes
 - * State-less transmission protocols, e.g. UDP
 - * No packet loss or transmission errors



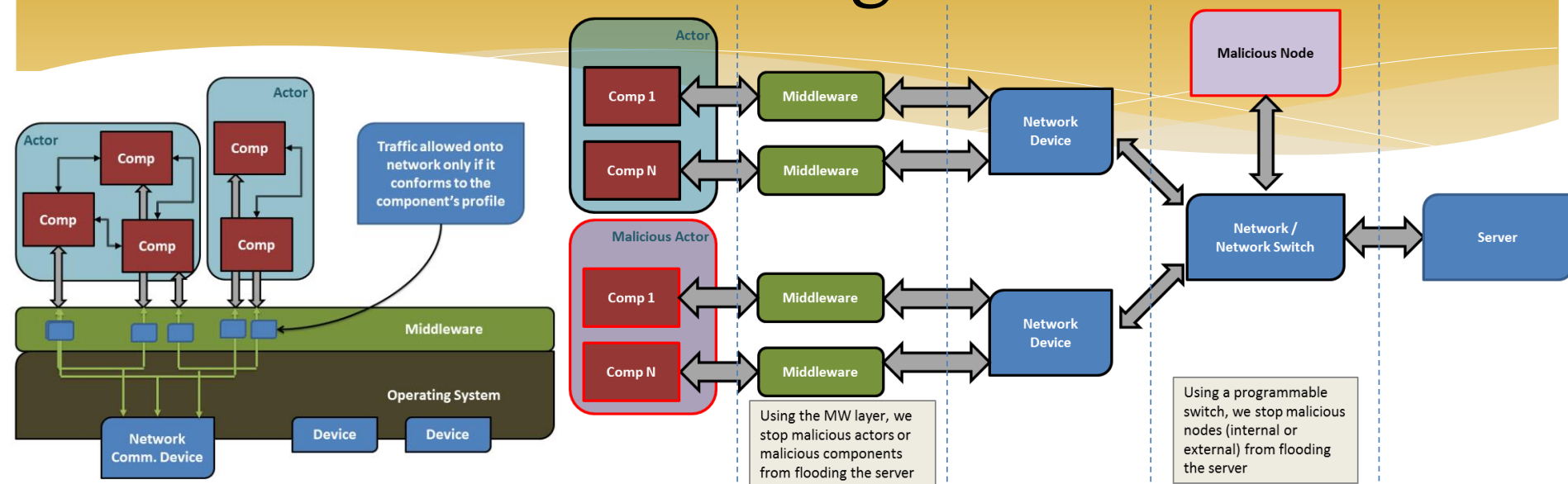
$$y = o[t] = (r \otimes p)[t]$$

$$= \min(r[t], p[t] - (p[t-1] - o[t-1]))$$

$$buffer = \sup\{r[t] - o[t] : t \in \mathbb{N}\}$$

$$delay = \sup\{o^{-1}[y] - r^{-1}[y]\}$$

DDoS Mitigation



- * Using PNP²:
 - * Precisely model application network behavior, with some bounds on deviation
 - * Assume infrastructure is controlled and verified
 - * only applications may be compromised
 - * Middleware detects that application traffic production deviates from model
 - * Use out-of-band communication between server and clients
 - * Server sees multiple clients simultaneously producing more data than normal
 - * Informs client-side middleware to throttle clients and prevent denial of service

Conclusions

- * We have developed a foundation for precisely modeling and analyzing at design-time applications and CPS networks
- * Implementing the modeling semantics in a trusted middleware allows management of the network performance of the CPS at run-time
- * Since these models capture the behavior of well-behaved applications, the middleware can detect deviations from this behavior and act accordingly
 - * Sender side detection: data produced far exceeds the model; throttle
 - * Receiver side detection: each sender produces slightly more data than the model and is allowed on the network; receiver cannot handle aggregate traffic so must discriminate good senders from bad
- * Using capabilities in the middleware, the misbehaving senders can be throttled at the source



Thank you!

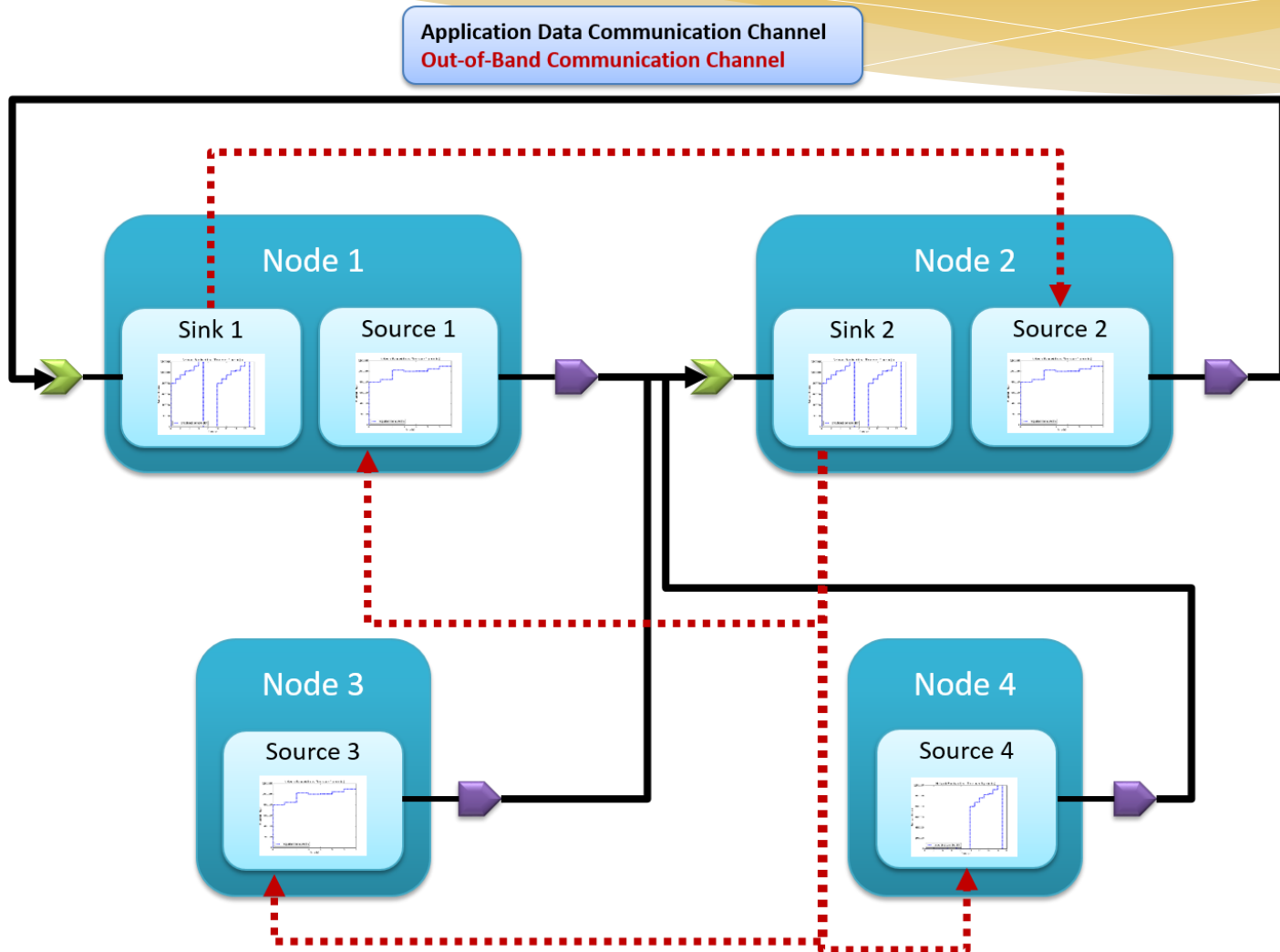
Questions?



Algorithm

```
receiver::limit_ddos( t_start, t_end )
{
    for sender in senders
    {
        d_start = sender.received_data(t_start)
        d_end = sender.received_data(t_end)
        profile_d_start =
            sender.profile(t_start)
        profile_d_end =
            sender.profile(t_end)
        allowed_data = profile_d_end - profile_d_start
        actual_data = d_end - d_start
        if actual_data > allowed_data
        {
            sender.disable()
        }
    }
}
```

Example



PNP² DDoS Detection

- * DDoS attacks are generally classified as excessive traffic from a large amount of (possibly heterogeneous) sources targeted towards a single point or a single group
- * The receiver has a finite receive buffer and a specified profile governing the rate at which it will pull data from the buffer:
 - * Senders who send too much data may cause the buffer to fill up, causing data from well-behaved senders to be lost
 - * If we can detect which senders are misbehaving and prevent them from sending, well-behaved senders will not see any degradation of service
- * Developed an out-of-band (OOB) communications channel from the sender middleware to the receiver middleware:
 - * Invisible to the application
 - * Allows the receiver to communicate with and control the sender's middleware
- * Works because we know precisely the correct behavior of applications and we know all of the communication patterns and we have a trusted middleware