

# Go-RealTime: Knowledge and Control of Time in High Level Programming Language



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## Introduction

**Go-RealTime** is a new multiprocessor real-time framework, based on Go language. It integrates key functionalities of programming **time sensitive applications** into language runtime: **Quality of Time (QoT)**, **Real-Time Scheduler** and **Parallel Programming**.

This work provides the software/programming language level support for the **RoseLine QoT Stack**.

## APIs

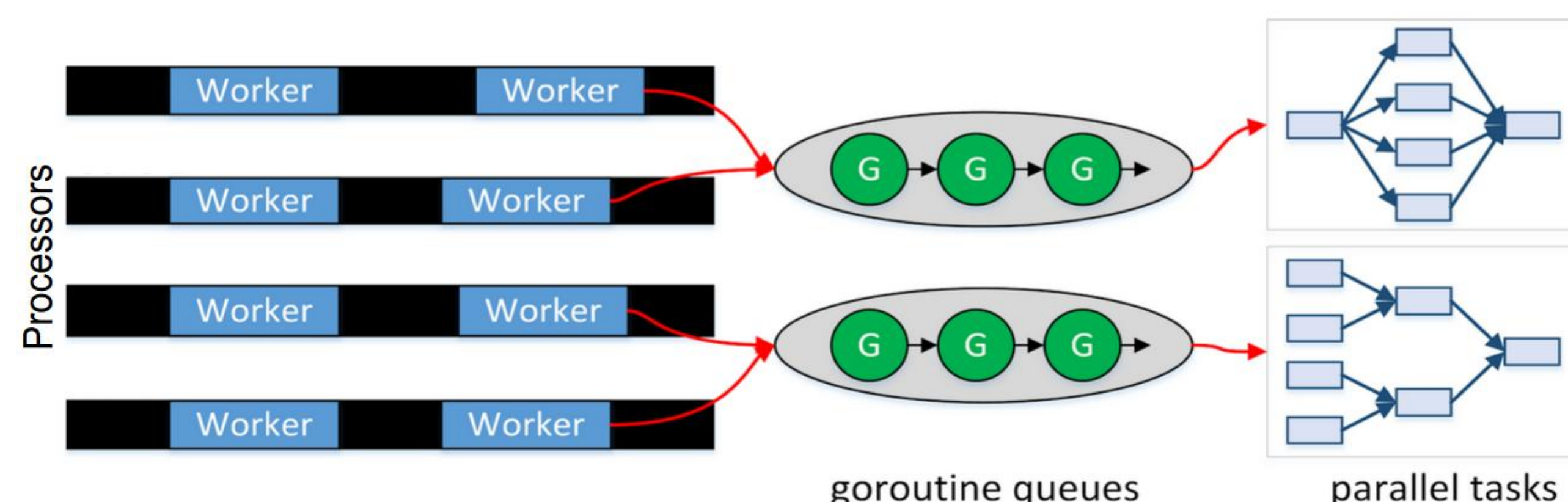
We import OS APIs (thread scheduler, CPU affinity and timer) into Go, modify Go runtime to enable **direct goroutine switch**, and build external Go packages.

Type	Method	Description
Resource	<code>SetSched(thread, policy, priority)</code>	Linux thread scheduler
	<code>BindCPU(CPU)</code>	Processor affinity
	<code>SetTimerFd(file_descriptor, time)</code>	Linux timer
Scheduling	<code>GoSwitch(goroutine)</code>	Switch to a goroutine
Task	<code>NewWorker(Nw)</code>	Create <i>Nw</i> workers
	<code>task.SetTimeSpec(ts, tp, td, tr, QoT)</code>	Set timing specification
	<code>task.Run(func, arg)</code>	Start a RT task

## Design

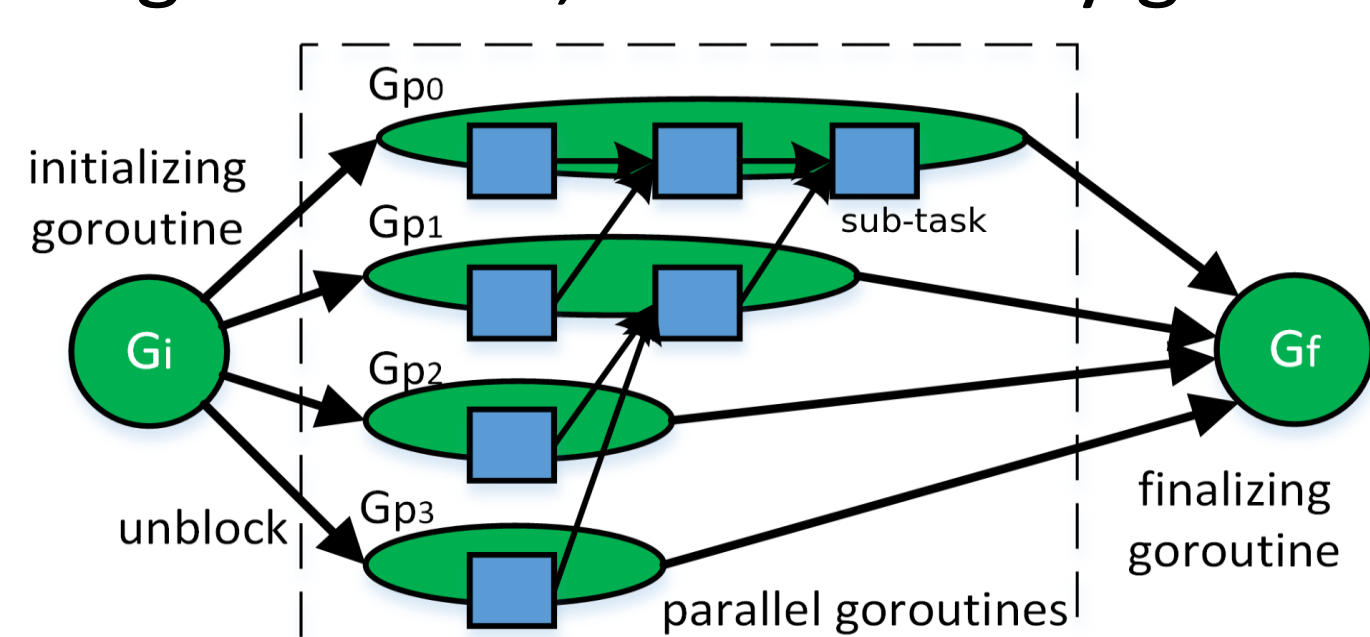
### System Model

**Worker** is the abstraction of thread with a processor resource reservation. Goroutines are sorted in a few priority queues. A sequential task is associated with a goroutine.

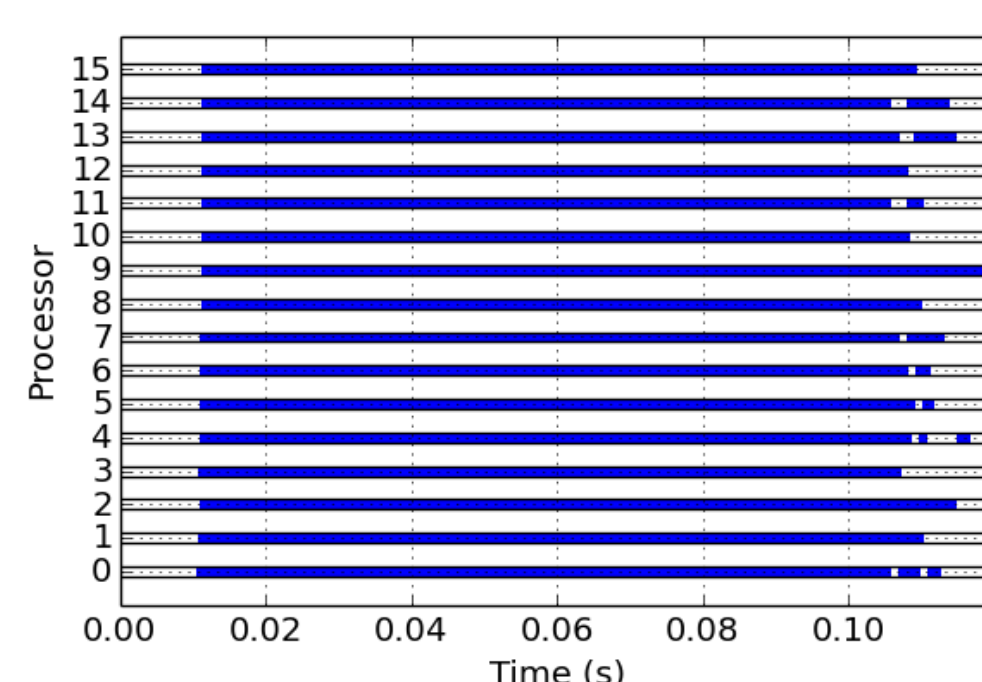


### Parallel Programming

Go-RealTime models a **parallel task** as a **DAG of sub-tasks**, executed on a group of parallel goroutines, scheduled by global EDF algorithm.



Synchronization of parallel goroutines via channels



Parallel dynamic programming on 16 cores

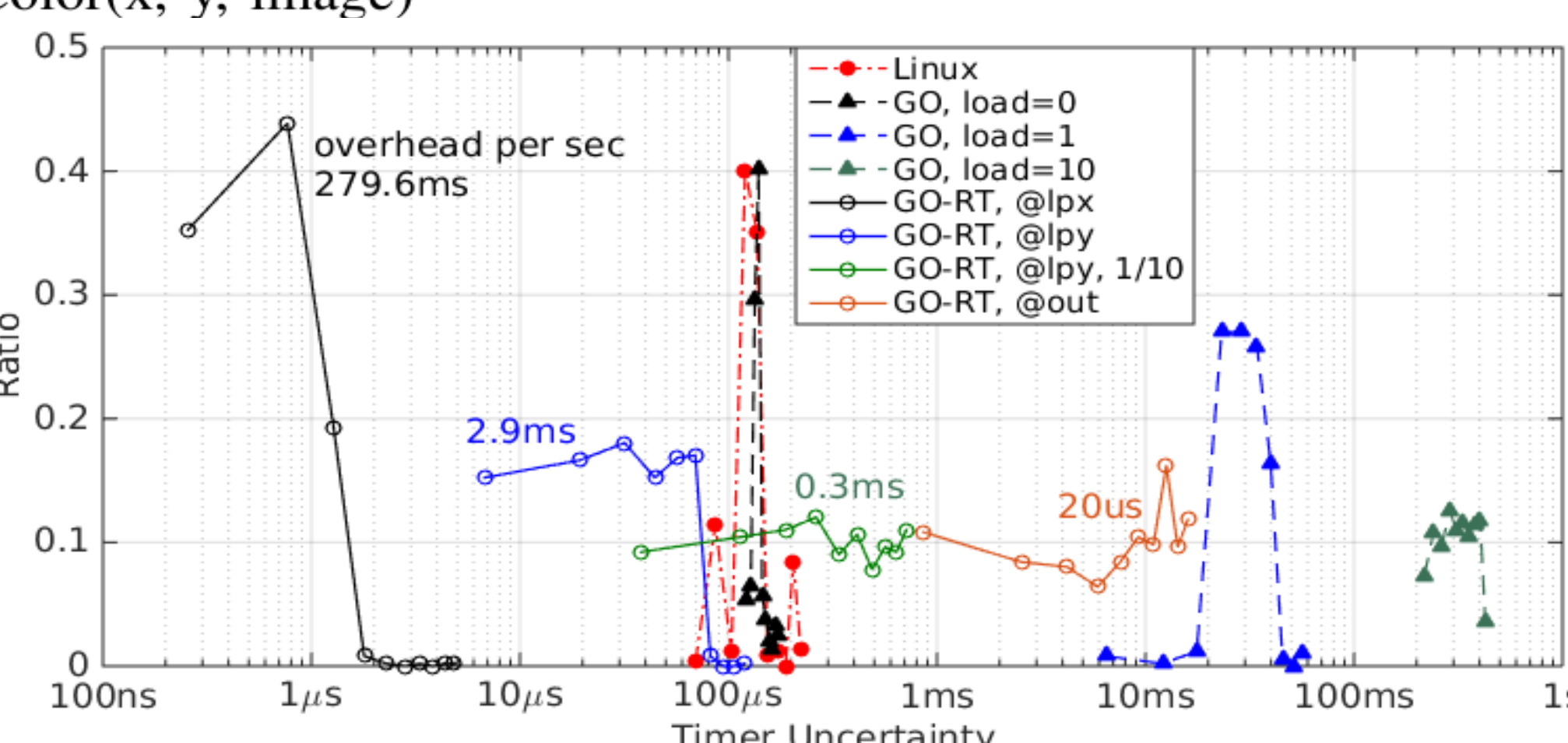
### Asynchronous Events and QoT

Go-RealTime handles **asynchronous events** via '**check\_async**' method. It checks all asynchronous events (timeout, message, etc.) and responds to the events which should have occurred.

Source code of Go-RealTime programs is instrumented with **check\_async** calls by Go parser library. The locations of calls are adjusted ahead according to measured timing profile. The calls are dynamically enabled/disabled to satisfy **QoT** requirement at runtime.

```

1: @out: check_async()
2: for loop over y coordinate do
3:   @lpy: check_async()
4:   for loop over x coordinate do
5:     @lpx: check_async()
6:     SetGrayColor(x, y, image)
7:   end for
8: end for
Instrumented
convert_gray
function
    
```



Timer accuracy of Linux, Go and Go-RealTime

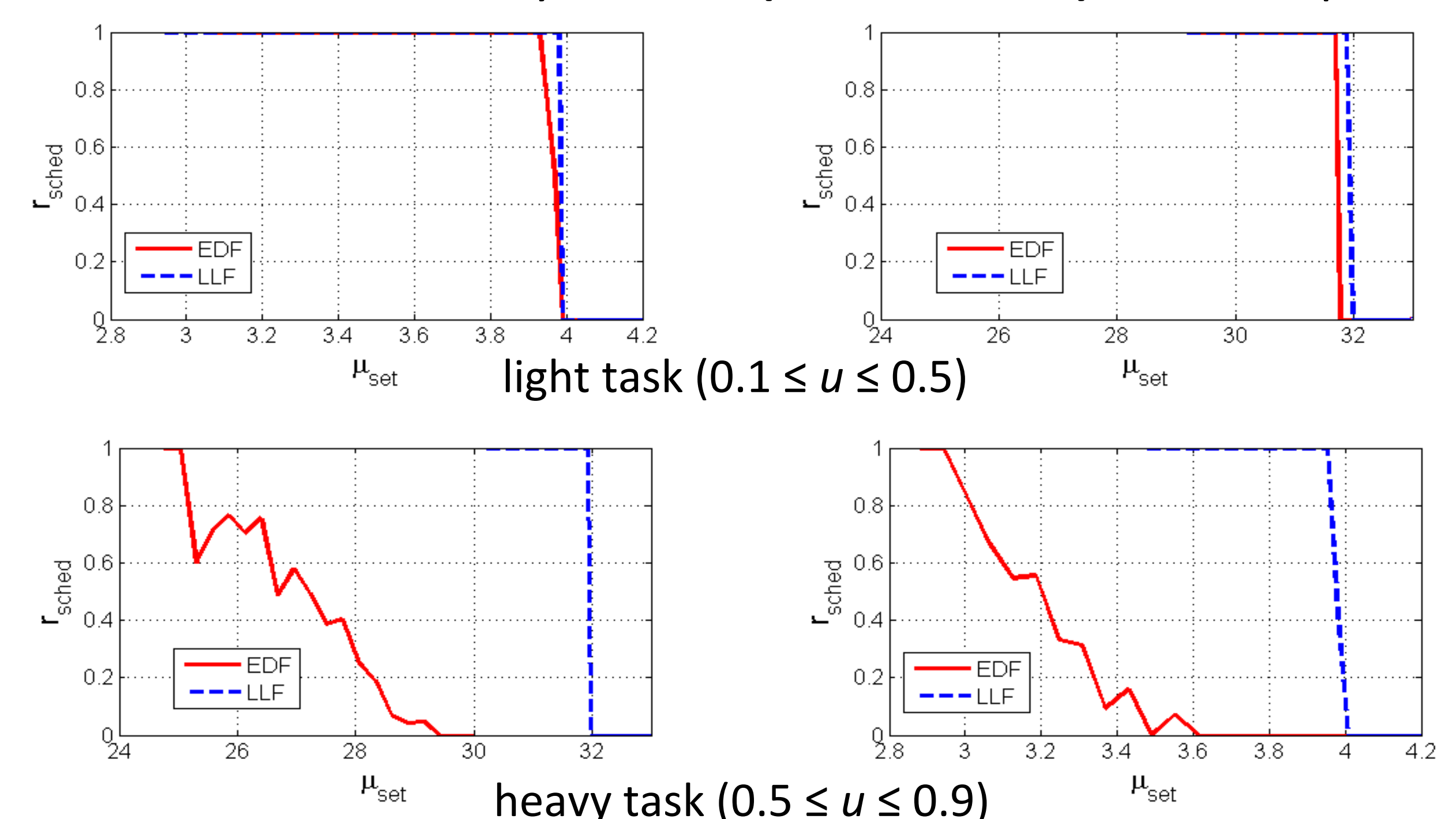
## Features

- **Goroutine** is the concurrency unit on which sequential and parallel tasks are running, controlled by Real-Time Schedulers.
- **OS thread** is the unit of **processor resource reservation**. Linux thread schedulers (FIFO, RR, etc.) are imported into Go.
- **Task** describes a real-time program with a **timing specification**: {start *ts*, period *tp*, deadline *td*, worst case budget *tr*, utilization  $u = tr/tp$ , QoT}
- Adaptive timing accuracy of asynchronous events (e.g. timer, task switching) upon **QoT** requirements of tasks.
- **Parallel tasks** are programmed as Directed Acyclic Graph (DAG) of sub-tasks and executed on multiprocessor.
- Different types of tasks are scheduled by the suitable algorithm: Earliest Deadline First (EDF), Least Laxity First (LLF).
- The total processor resources are partitioned dynamically among all schedulers according to utilization of tasks.

## Real-Time Scheduler

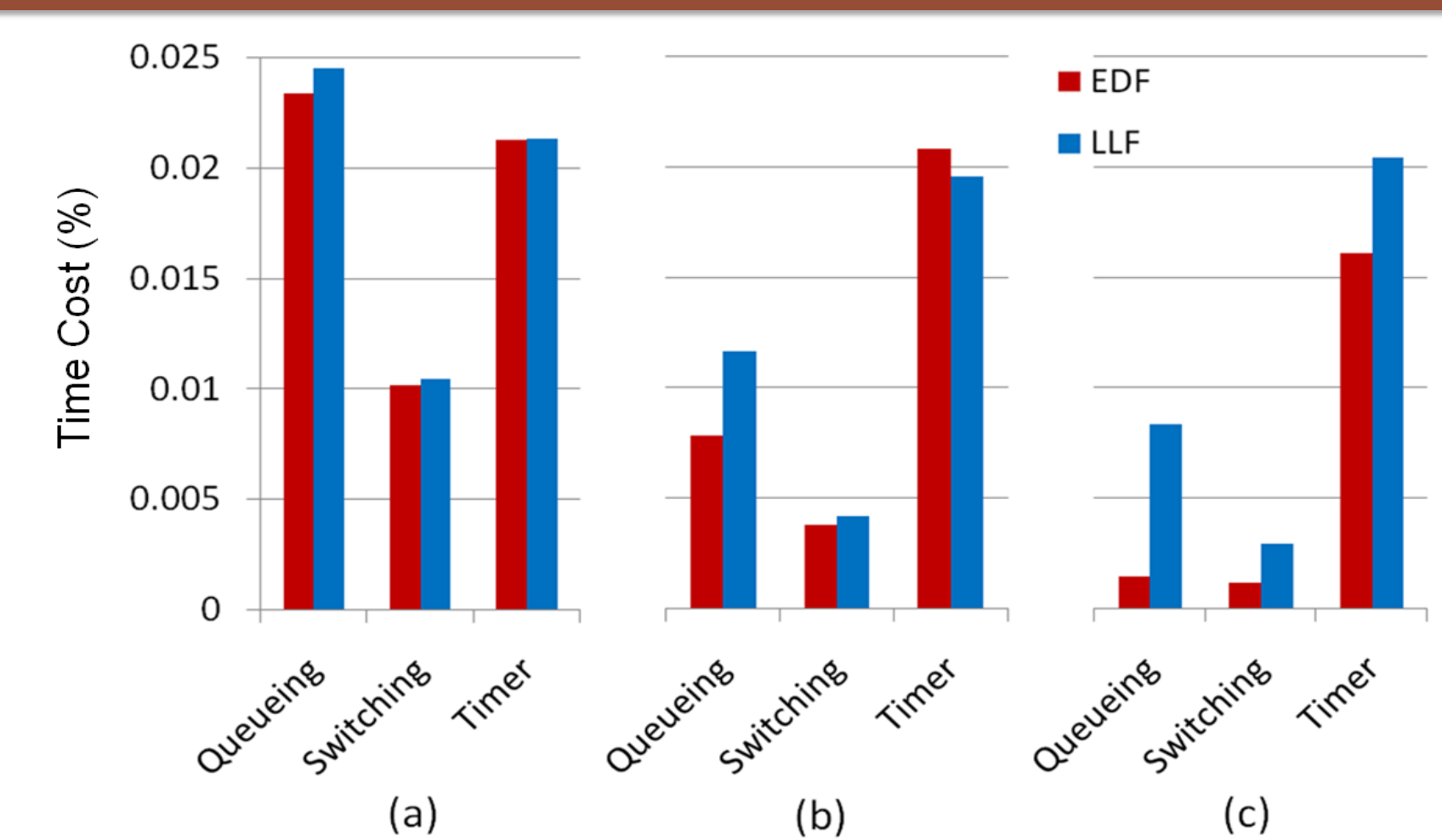
### Algorithm

**EDF/LLF** schedulers are implemented for sequential tasks. For **LLF**, Go-RealTime compares laxities of the running task with the head task in queue via **check\_async** method. A switch is allowed only after **tlif (default 10ms)** since the previous.



### System Cost

- **Queueing** cost: time consumed by scheduler code
- **Switching** cost: direct cost of goroutine switch
- **Timer** cost: cost of operations on timer queue such as insertion
- **Indirect** cost: cache pollution due to switching



Period: 100ms ≤ *tp* ≤ 300ms

(a) ultra light task ( $0.01 \leq u \leq 0.1$ ); (b) light task; (c) heavy task

### Scheduling Algorithm Change

In the example, when the total utilization increases, the scheduler changes from EDF to LLF in order to avoid deadline miss.

