



Resilient Monitoring and Control of Distributed Cyber-Physical Systems

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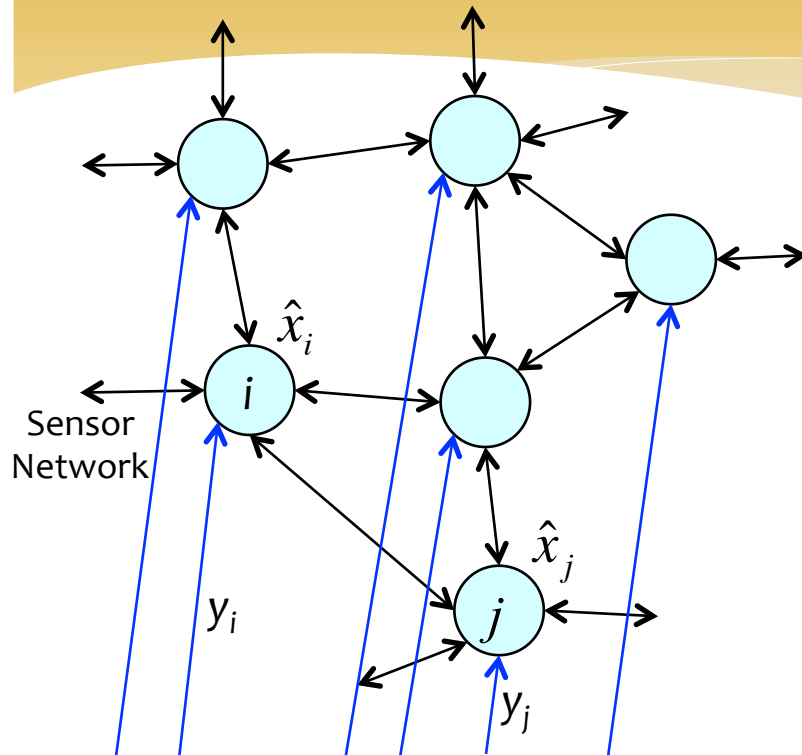
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Distributed Parameter Estimation



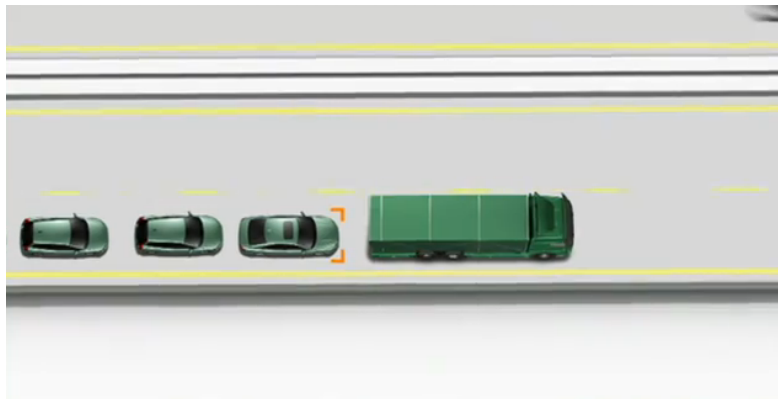
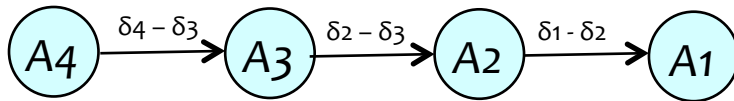
- All sensors measure independently some physical phenomenon with some error due to noise
 $y_i = \theta + v_i, v_i \sim N(0, \sigma_i^2), i = 1, 2, \dots, n$
- The sensors improve their estimate by averaging the measurements
- Minimum variance estimate

$$\hat{\theta}_{MV} = \frac{\frac{1}{n} \sum_{i=1}^n \frac{1}{\sigma_i^2} y_i}{\frac{1}{n} \sum_{j=1}^n \frac{1}{\sigma_j^2}}$$

- It can be asymptotically computed in a distributed fashion using two average consensus algorithms in parallel



Distributed Control of Multi-Agent Systems



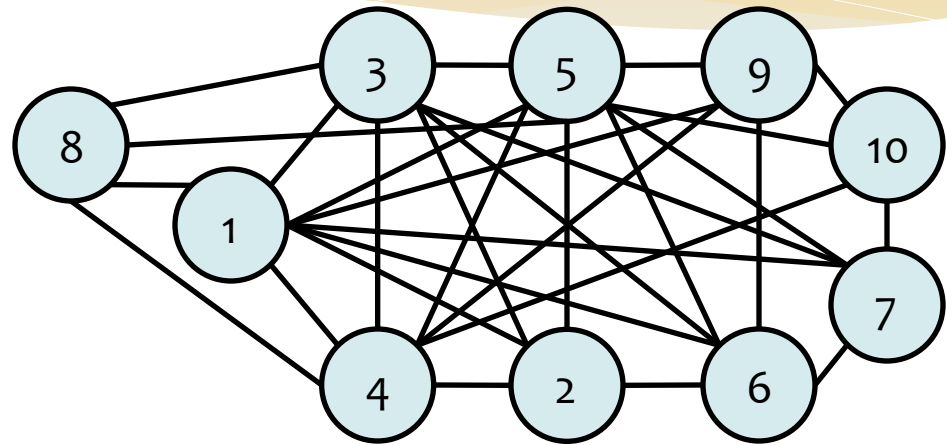
Consensus-based Formation Control

- * Distributed consensus
 - * Each vehicle updates its state based on the states of its local neighbors
 - * The final state of each vehicle converges to a common value

- * Distributed Consensus Applications in CPS
 - * Vehicle rendezvous
 - * Formation control
 - * Parameter estimation
 - * Least squares data regression
 - * Sensor calibration
 - * Time synchronization
 - * Kalman filtering

Resilient Consensus in the Presence of Adversaries

(3,2)-robust graph: resilient consensus in the presence of 1 adversary



- * Adversarial Consensus Protocol
- * Adversary models
 - * Threat
 - * Scope
- * Robust network topologies
 - * Local redundancy
- * Resilience requires high degree of redundancy
- * Can we relax the redundancy requirements?

Overview

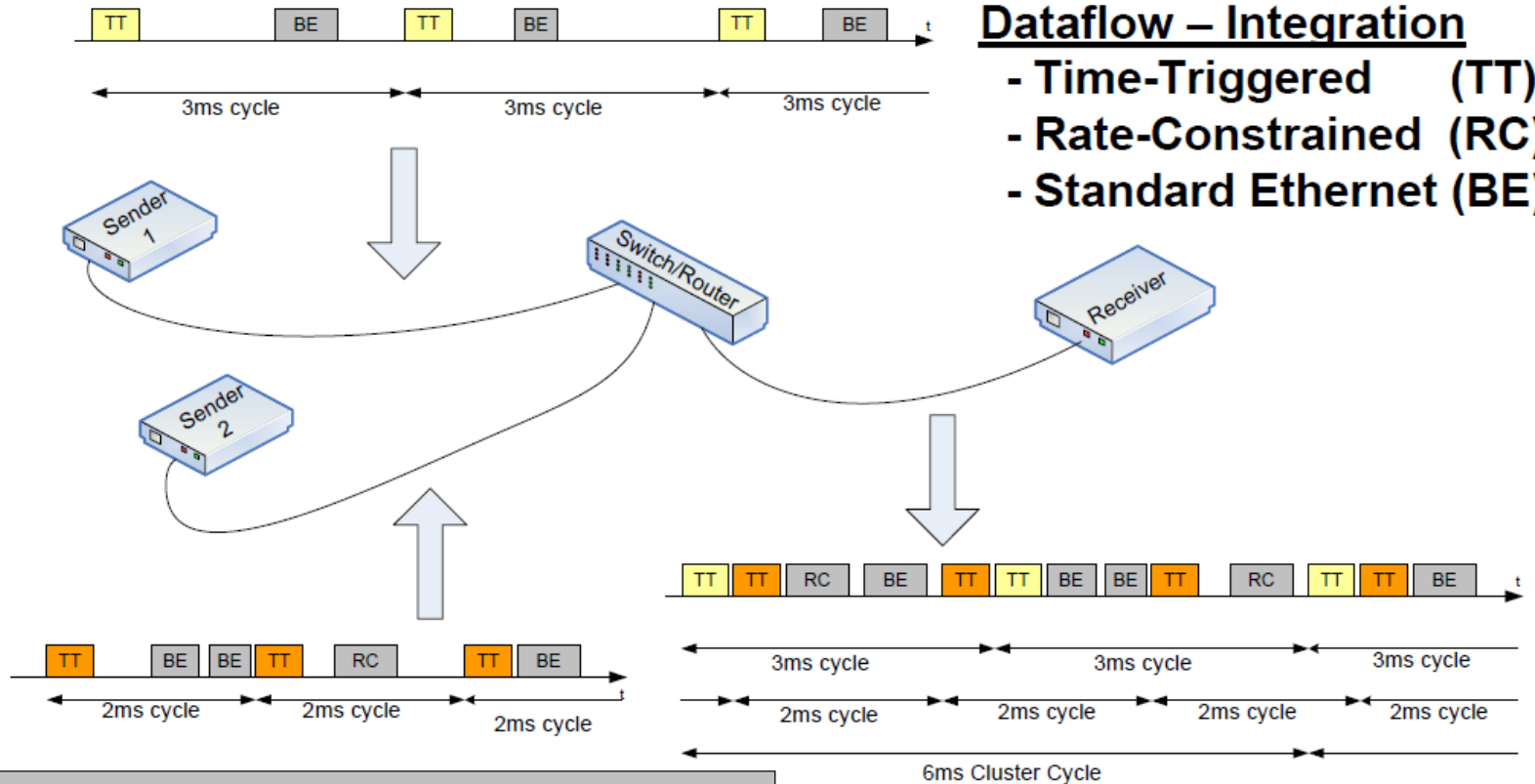
- * Performance Impact of Authentication in Time-Triggered Networked Control Systems
 - * Theoretical analysis of performance impact
 - * Experimental validation
- * Resilient Consensus Protocols with Trusted Nodes
 - * Connected Dominating Set
 - * Trusted Nodes and Network Robustness
- * Stochastic Message Authentication
 - * Game Theoretic Model
 - * Trade-off Between Computation and Security
- * Conclusions

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Time-Triggered Ethernet Overview

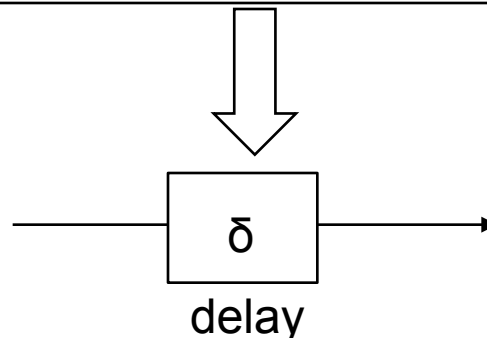
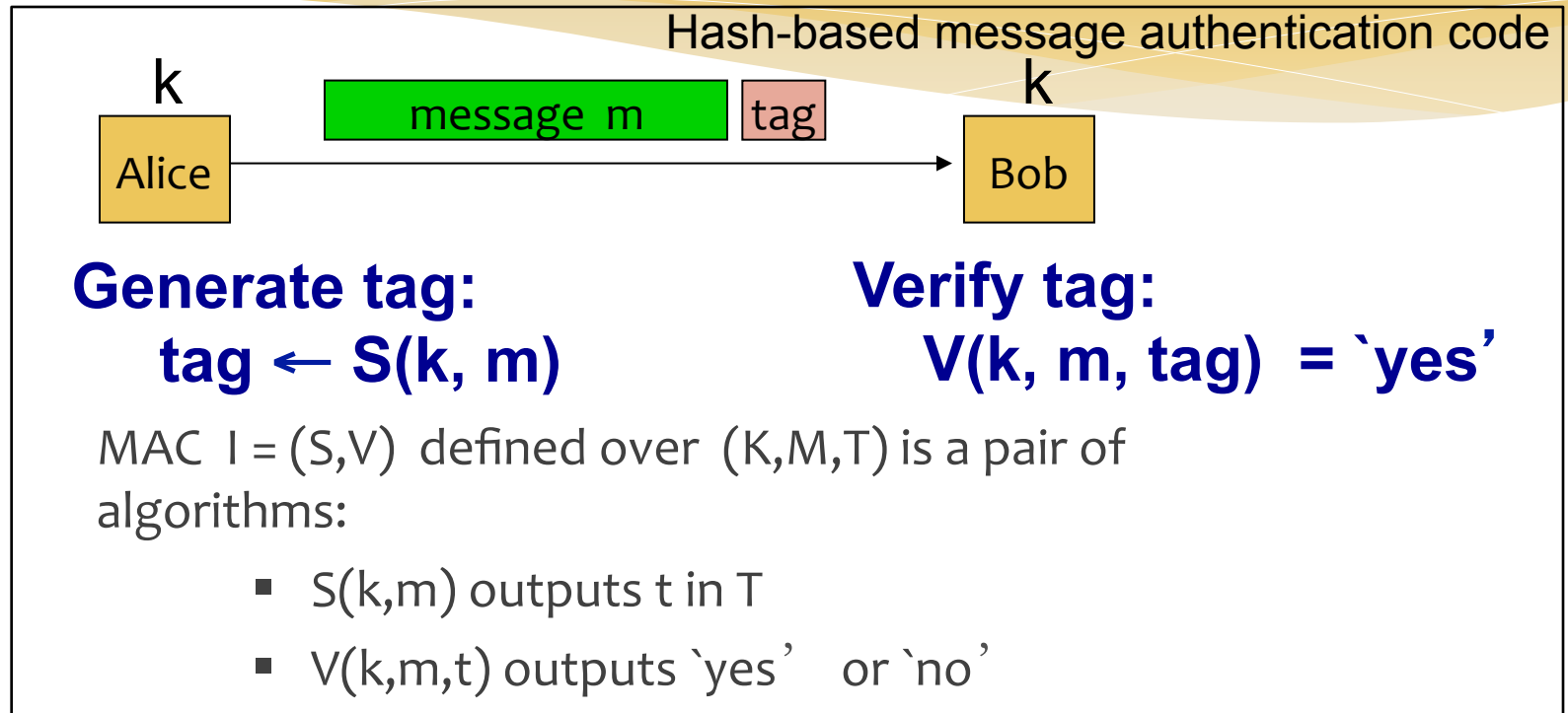
Integrated Dataflow Example



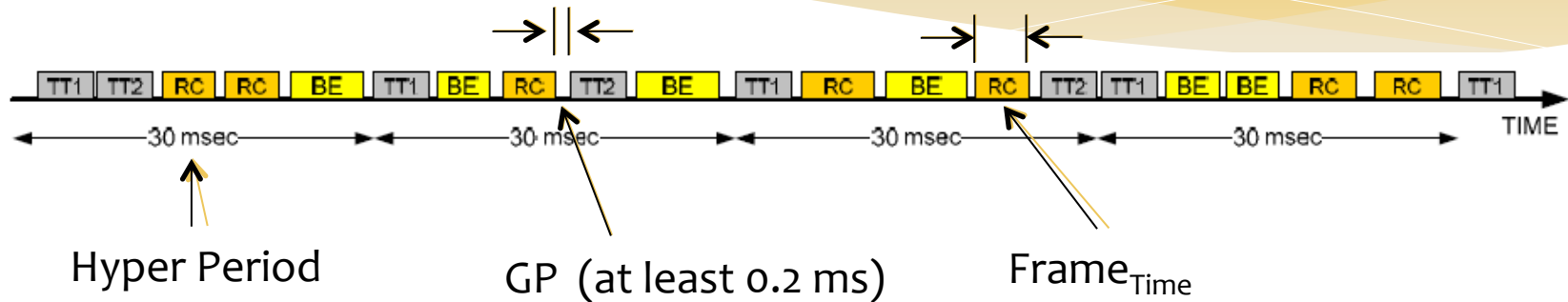
In parallel, two types of Ethernet communications:

- Synchronous (TDMA-style) Communication: TT
- Asynchronous (event-triggered style): RC + BE

Protected Message Transmission



Analysis of Performance Impact



- Max Number of Frames per Hyper Period (NF_{Max})

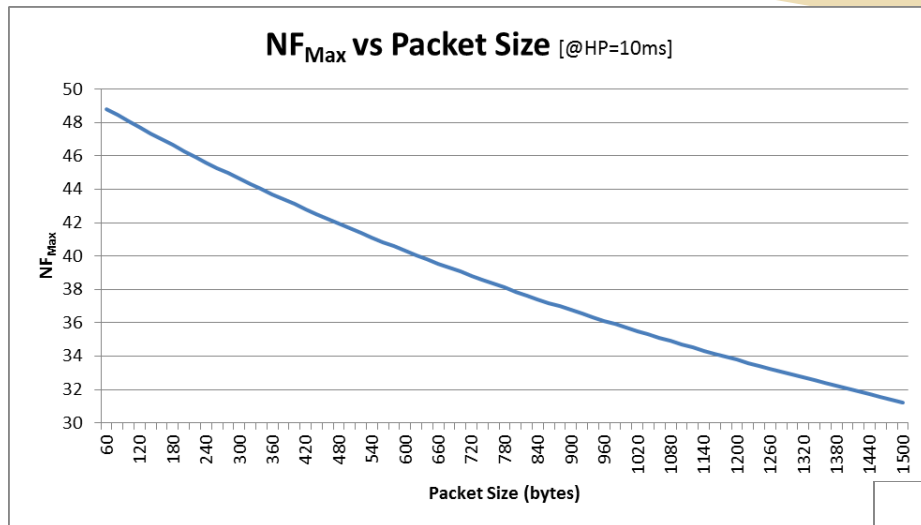
$$NF_{Max} = \frac{HP}{(Frame_{Time} + GP)}$$

$$Frame_{Time} = \frac{Packet_{Size}}{Transmission_{Rate}}$$

- Example:

$$NF_{Max} = \frac{10 (ms)}{(Frame_{Time} + 0.2 (ms))} \cong 48 \quad Frame_{Time} = \frac{60 (bytes)}{[(100(Mbits)/8)] \left(\frac{bytes}{s} \right)}$$

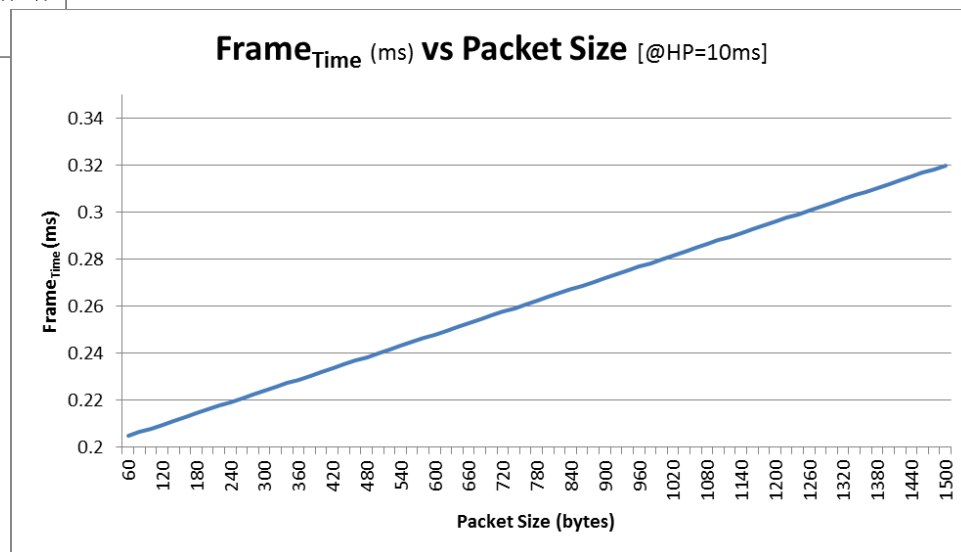
Effect of Packet Size



Guard Period = 0.2 ms



HP = 10 ms



Hardware Platform: IBX-530W

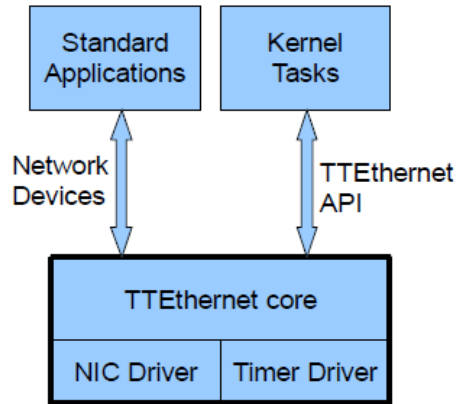
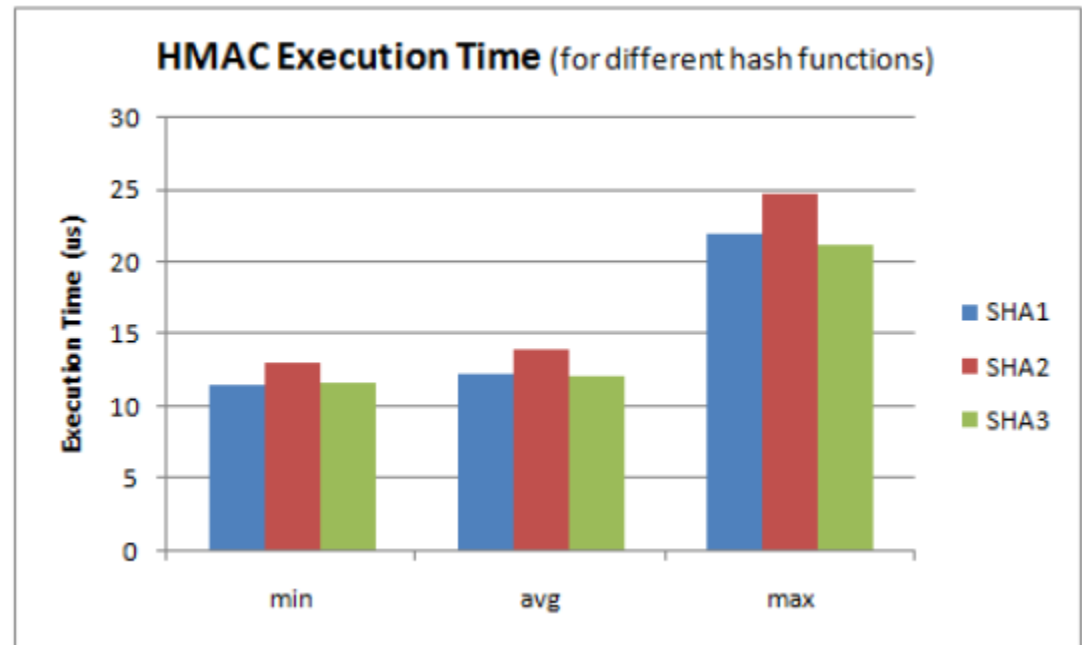


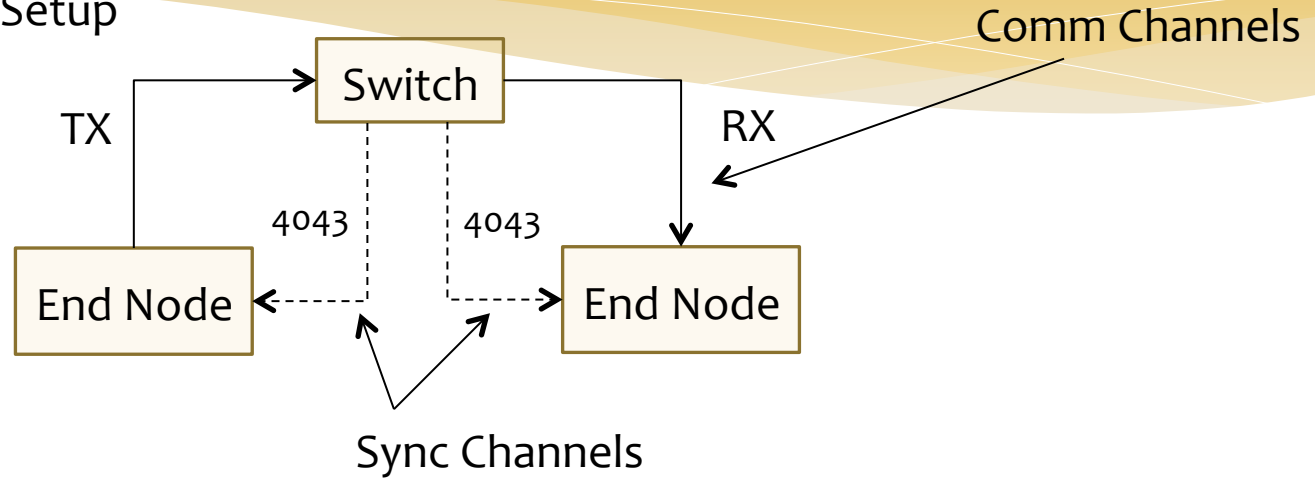
Figure 1: Structure of a TTEthernet end system

- * Intel Atom Processor, 1.6GHz
- * Linux 2.6.24-24-rt kernel
- * Crypto library

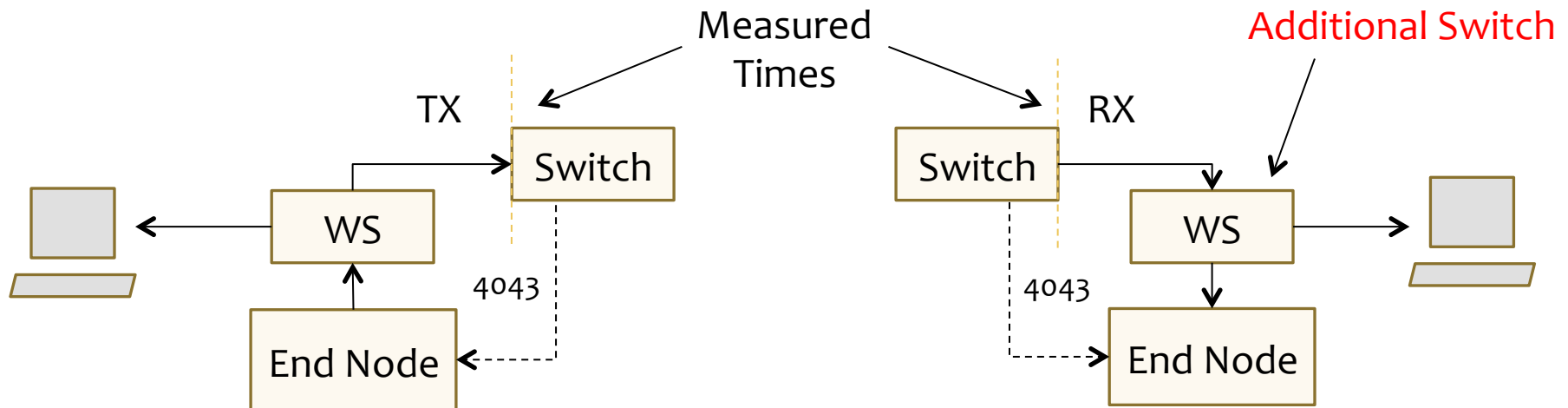


Impact on System Performance

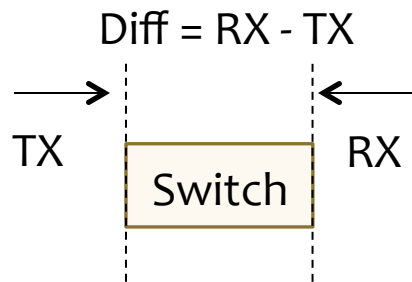
Original Physical Setup



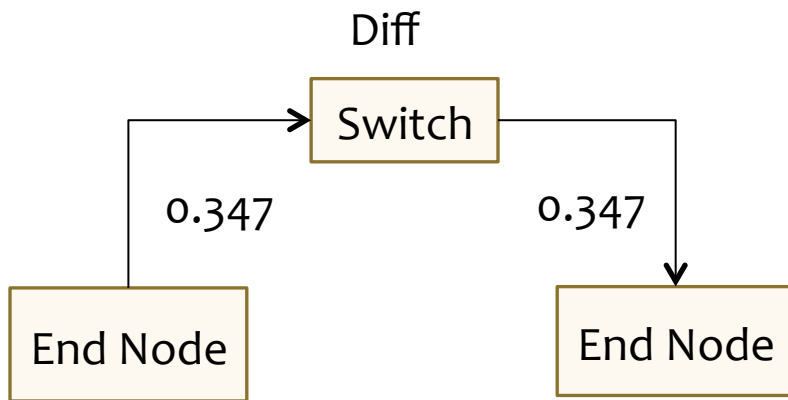
WireShark (WS) Physical Setup



Experimental Results (WireShark)



	60 bytes			1514 bytes		
	Min	Avg	Max	Min	Avg	Max
Tx (ms)	0	0.008	0.115	0.110	0.222	0.347
Rx (ms)	0.274	0.386	0.673	0.494	0.594	0.826
Diff	0.274	0.378	0.558	0.384	0.372	0.479



□ Max Total Transmission Time (Max_{TTT})

$$Max_{TTT} = (2 * Frame_{Time}) + Diff$$

$$(0.347 * 2) + 0.375 = 1.069 \text{ ms}$$

With WirelessShark Switch

$$(0.347 * 2) + 0.2 = 0.894 \text{ ms}$$

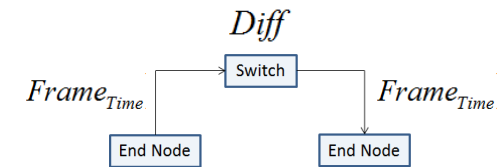
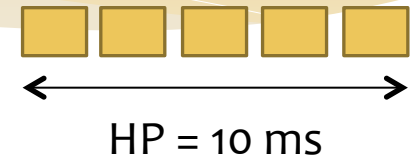
With assumed GP

Comparison

TTTech – TDMA Theoretical Results

Theoretical Values			
	60 Bytes	80 Bytes	1514 Bytes
NF_{Max}	48	48	31
$Frame_{Time}$ (ms)	0.0048	0.0064	0.12
Max_{TTT} (ms)	0.2096	0.2128	0.44

Guard Period = 0.2 ms



TTTech – TDMA Hardware Results

$$Max_{TTT} = (2 * Frame_{Time}) + Diff$$

GP

Hardware Values			
	60 Bytes	80 Bytes	1514 Bytes
NF_{Max}	23	20	11
$Frame_{Time}$ (Tx_{Max}) (ms)	0.115	0.150	0.347
Max_{TTT} (ms)	0.43	0.5	0.894

Experimental Analysis: Conclusions

- * The overhead time introduced by the kernel module implementing HMAC reduces the effective number of frames per hyper-period (HP)
- * There is a small impact on the maximum number of frames per HP by increasing the packet size from 60 to 80 bytes (tag)
- * Experimental results are consistent with the theoretical analysis
 - * Overhead time spent by the kernel module to transmit data to the physical medium is not considered by the theoretical analysis

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- * **Resilient Consensus Protocols with Trusted Nodes**
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Resilient Consensus Protocol with Trusted Nodes

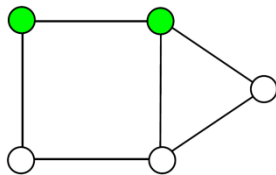
Under RCP-T, consensus is always achieved in the presence of *arbitrary number of adversaries* iff there exists a set of trusted nodes that form a **connected dominating set**.

Under RCP-T

- Any number of attacks can be handled
- Sparse networks can be made resilient

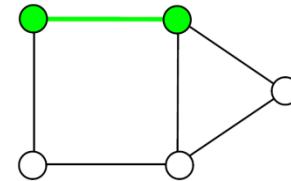
Dominating Set:

$$D \subseteq V, \quad \text{s.t.} \quad \bigcup_{v_i \in D} \mathcal{N}[v_i] = V$$



Connected Dominating Set:

Nodes in the dominating set induce a **connected** subgraph



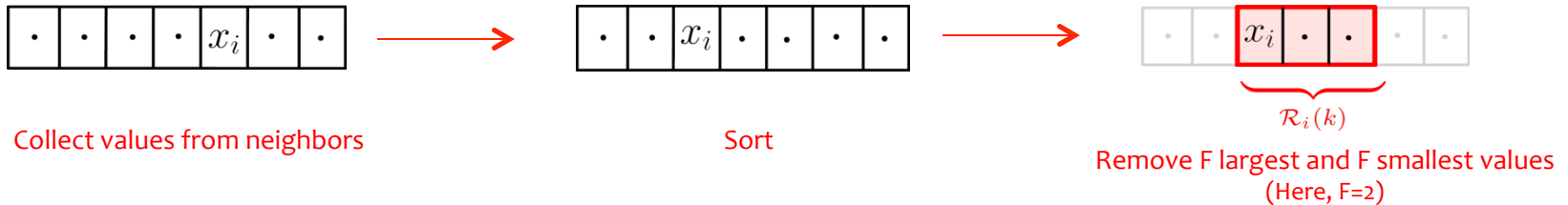
Resilient Consensus Protocol with Trusted Nodes (RCP-T)

Each normal node updates its value according to the following update rule

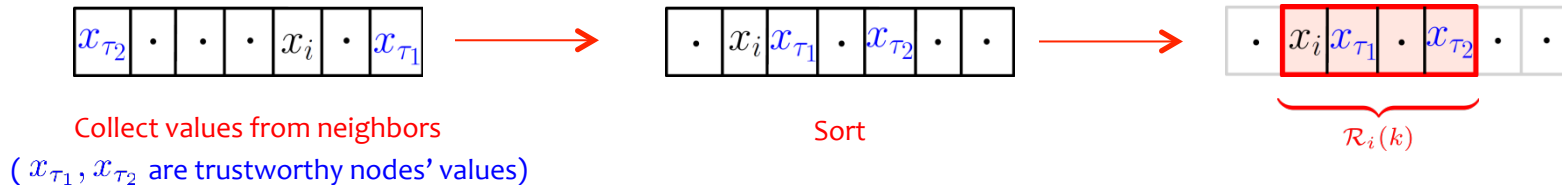
$$x_i(k+1) = \sum_{j \in \mathcal{R}_i(k)} w_{ij} x_j(k)$$

What is $\mathcal{R}_i(k)$?

- if node i is **not connected** to any trusted node
(F is the total number of attacks that can happen within the network)

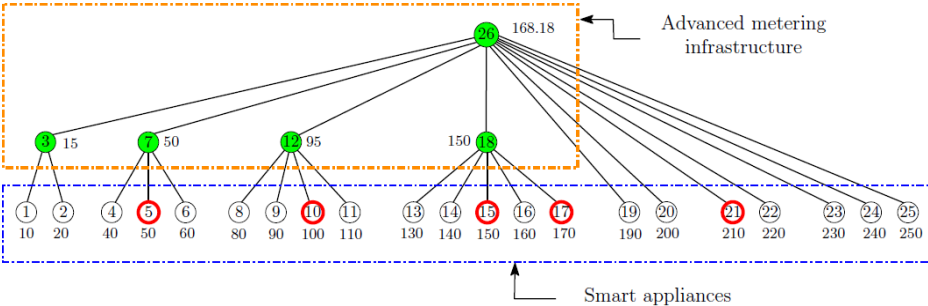


- if node i is **connected** to at least one **trusted node**

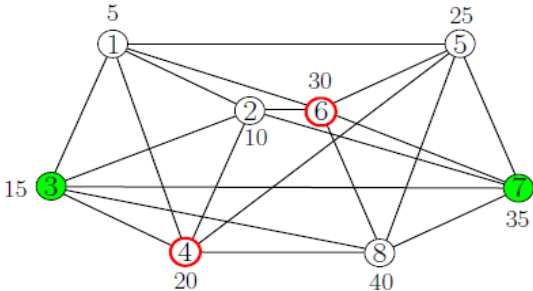


Examples

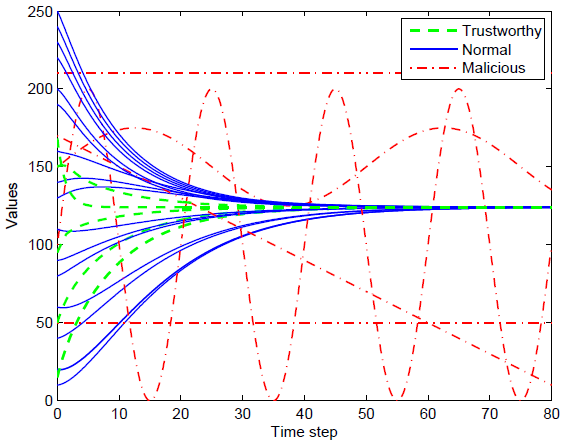
Example 1: (Tree – Sparse network)



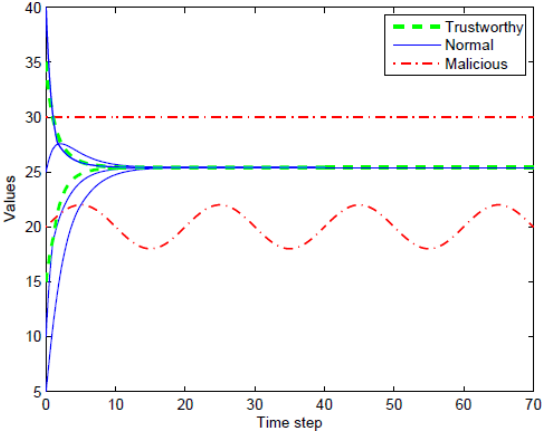
Example 2: (2,2) Robust graph



RCP-T



RCP-T



[Abbas et al., ISRCS 2014, Submitted]

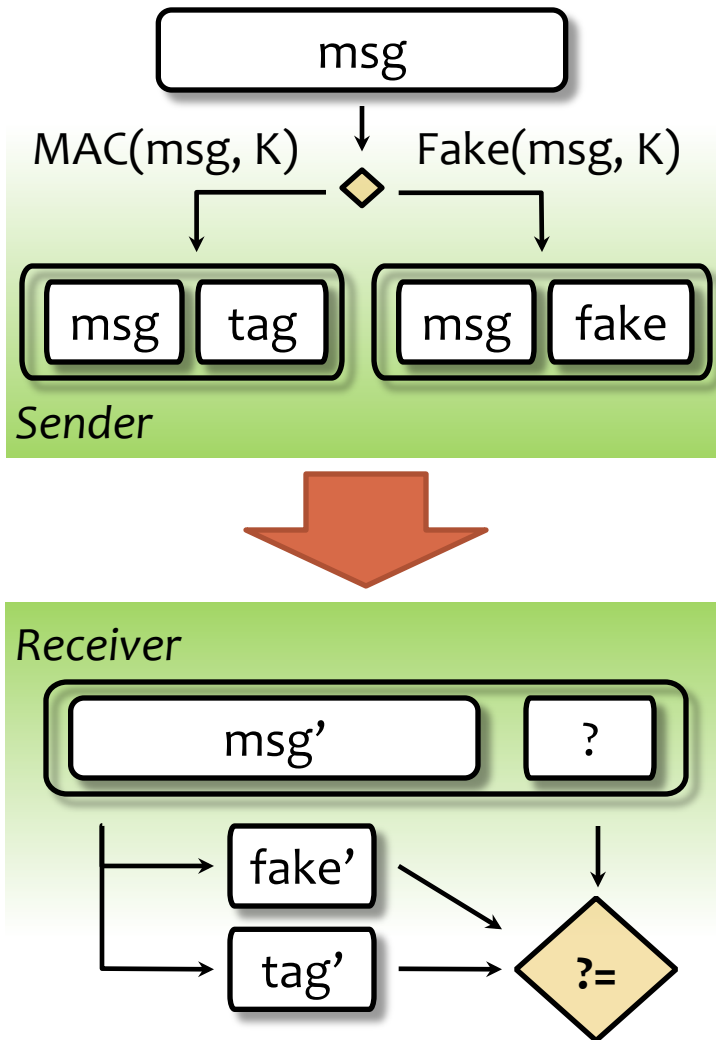
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- * **Stochastic Message Authentication**
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- * **Conclusions**

Motivation

- * Computational demand of cryptographic primitives can be too high for **resource-bounded** devices
 - * legacy devices in supervisory control systems
 - * embedded or battery-powered devices (RFID tags, sensors)
- * "Lightweight" cryptographic primitives
 - * Decision to secure a system is still **binary**: either security is employed, incurring some fixed overhead, or it is not
- * Our approach: General-purpose framework for trading off security and computational resources using an existing MAC scheme

Stochastic Message Authentication



- * For some messages, the sender computes a "fake tag", which is computationally less demanding, but does not protect integrity
- * Adversary cannot distinguish fake tags from correct tags
- * Receiver can verify if a message has a fake or a correct tag efficiently
→ detect attacks with high probability

Game-Theoretic Model

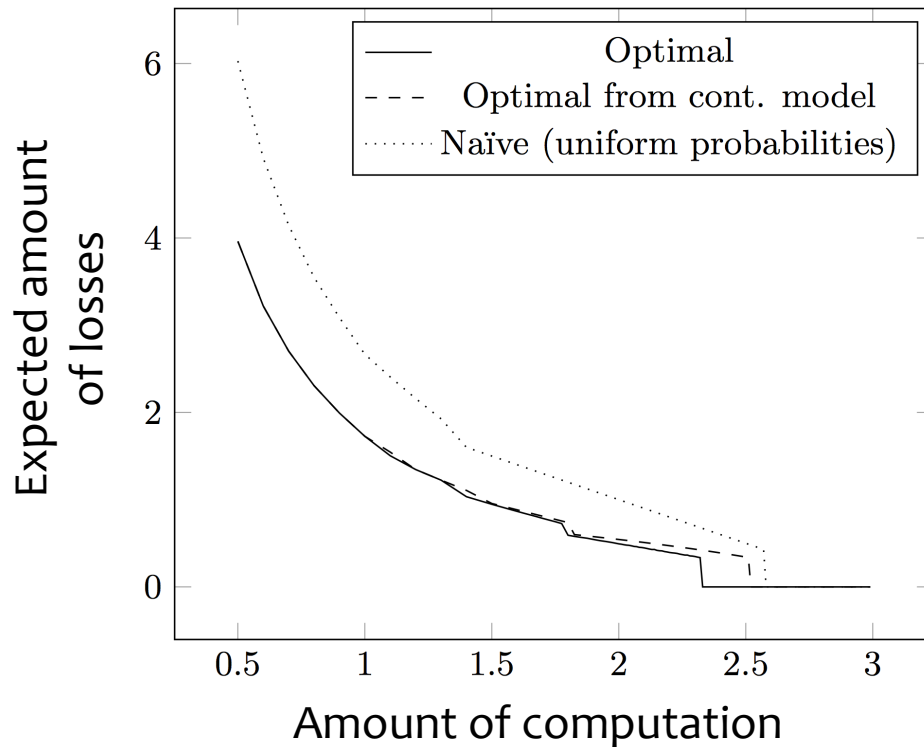
- * Stackelberg security game

- * Divide messages into C classes based on their potential to cause damage

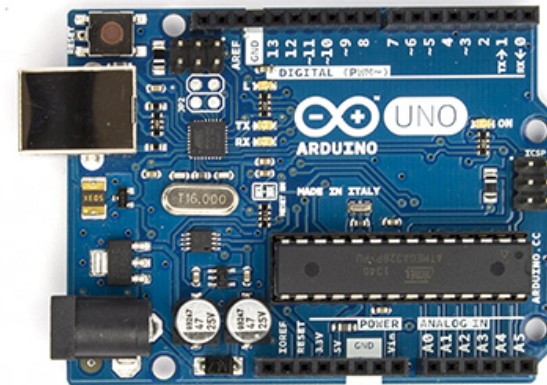
		Defender	Attacker
Strategy choice		for each class c , the probability of authentication p_c	for each class c , the number of modified / inserted messages a_c
Detection probability		$1 - \prod(1 - p_c)^{a_c}$	
Payoff	attack undetected	- (amount of total damage $\sum a_c L_c$)	amount of total damage $\sum a_c L_c$
	attack detected	zero	“punishment” $-F$

Results

Trade-off between computation and security



Proof-of-concept implementation using SHA-1 HMAC on an ATmega328P microcontroller



[Laszka et al., CCS 2014, Submitted]

Conclusions

- * Resilient Distributed Consensus Protocols in the Presence of Adversaries
 - * Exploit local information redundancy
- * Performance Impact of Authentication Mechanisms
 - * Theoretical analysis and experimental validation
- * Resilient Distributed Consensus Protocols with Trusted Nodes
 - * Trusted nodes form a connected dominating set
- * Stochastic Message Authentication
 - * Trade-off between computation and security