



# Robust Convergence of Distributed Routing with Heterogeneous Population Dynamics

Walid Krichene

Alex Bayen

Dept. of Electrical Engineering & Computer Sciences,  
UC Berkeley, CA, USA



# Outline

## 1 Introduction

- Routing Game
- Motivation
- Existing results

## 2 Convergence

- Model
- Convergence of averages
- Convergence using Stochastic Approximation
- Convergence using Stochastic Mirror Descent (SMD)

## 3 Simulations

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# Routing game

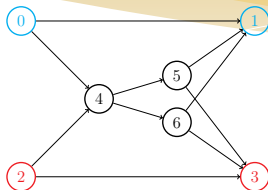


Figure: Example network

- Directed graph  $(V, E)$
- Population  $k$ : paths  $\mathcal{P}_k$
- Population distribution over paths  $x_{\mathcal{P}_k} \in \Delta^{\mathcal{P}_k}$
- Loss on path  $p$ :  $\ell_p(x)$

# Routing game

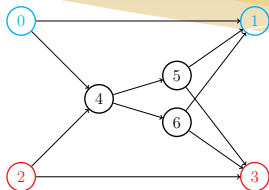


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# Routing game

## Equilibrium

$x^* = (x_{P_1}^*, \dots, x_{P_k}^*)$  is an equilibrium if  $\forall k$ ,

$$\langle \ell_{P_k}(x^*), x_{P_k}^* \rangle \leq \langle \ell_{P_k}(x^*), x_{P_k} \rangle$$

Losses are minimal on the support of  $x_{P_k}^*$

# The routing game

## One-shot routing game

- Well understood
- Useful for characterizing **'steady-state' behavior**
  - Network performance (price of anarchy)
  - System optimal tolling
  - Other applications

### Why study dynamics?

- How do players arrive at equilibrium?
- How fast?
- Stability?
- Robustness (noisy measurements)?

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# Applications

	Transportation networks	Packet routing	Load balancing
Time scale	Day	minute/second	minute/second
Measurements	Route delays	Route delays	Job completion
Decision model	Distributed	Distributed	Can be centralized

# Convergence rate

How fast does the system reconverge to equilibrium?

- Catastrophic failure: Mississippi river bridge collapse (2005)



# Convergence rate

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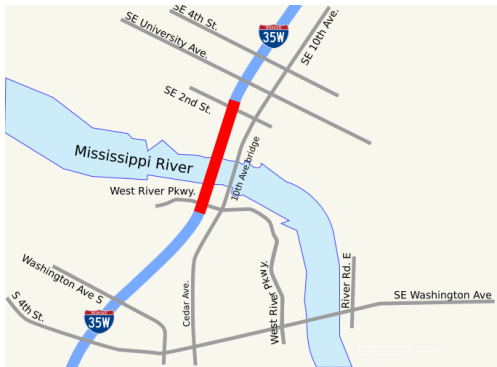
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# Convergence rate

How fast does the system reconverge to equilibrium?

- Catastrophic failure: Mississippi river bridge collapse (2005)



# Convergence rate

How fast does the system reconverge to equilibrium?

- Incident response: closure of I15 after fire on bridge during construction.



# Convergence rate

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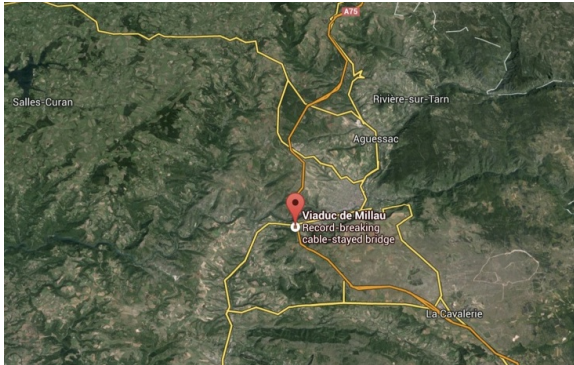
- Adding a link to the network: construction of the Millau Viaduct (2004)



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- Adding a link to the network: construction of the Millau Viaduct (2004)



# Convergence rate

How fast does the system reconverge to equilibrium?

- Tolling: Electronic Road Pricing (ERP) in Singapore.





## Existing results

### Continuous time:

- General case of potential games, under a positive correlation condition [9]
- Special case of routing games, under replicator dynamics [5]

### Discrete time:

- General class of no-regret dynamics, limited convergence result [2]

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[9] William H Sandholm. [Potential games with continuous player sets](#). *Journal of Economic Theory*, 97(1):81–108, 2001

[5] Simon Fischer and Berthold Vöcking. [On the evolution of selfish routing](#). In *Algorithms–ESA 2004*, pages 323–334. Springer, 2004

[2] Avrim Blum, Eyal Even-Dar, and Katrina Ligett. [Routing without regret: on convergence to nash equilibria of regret-minimizing algorithms in routing games](#). In *Proceedings of the twenty-fifth annual ACM symposium on Principles of distributed computing*, PODC '06, pages 45–52, New York, NY, USA, 2006. ACM

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- **Model**
- Convergence of averages
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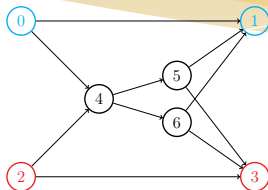


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# Online learning model

## Online learning model

At iteration  $t$

- Players of population  $k$  choose routes. Distribution  $x^{(t)}$ .
- $\ell_{\mathcal{P}_k}(x^{(t)})$  is revealed to players of population  $k$ .
- Players update their distribution.

$$x_{\mathcal{P}_k}^{(t+1)} = u_k(x_{\mathcal{P}_k}^{(t)}, \text{history})$$

## Main problem

Define a class  $\mathcal{C}$  of algorithms (update rules) such that

$$u_k \in \mathcal{C} \forall k \Rightarrow x^{(t)} \rightarrow \mathcal{N}$$

Extension: Losses are noisy  $\hat{\ell}_{\mathcal{P}_k}(x^{(t)})$  with

$$\mathbb{E}[\hat{\ell}_{\mathcal{P}_k}(x^{(t)}) | x^{(t)}] = \ell_{\mathcal{P}_k}(x^{(t)})$$

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# Convex potential

## Rosenthal potential

$f(x)$  Convex

$$\nabla_{x_{\mathcal{P}_k}} f(x) = \ell_{\mathcal{P}_k}(x)$$

$$\mathcal{N} = \arg \min_{x \in \Delta^{\mathcal{P}_1} \times \dots \times \Delta^{\mathcal{P}_K}} f(x)$$

Optimality conditions:

$$\langle \ell(x^*), x - x^* \rangle \geq 0 \quad \forall x \quad \Leftrightarrow \quad \forall k, \forall x_{\mathcal{P}_k}, \langle \ell_{\mathcal{P}_k}(x_{\mathcal{P}_k}^*), x_{\mathcal{P}_k} - x_{\mathcal{P}_k}^* \rangle \geq 0$$

- Continuous time:  $f$  used as a Lyapunov function.
- Discrete time: regret.

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# Regret

Instantaneous regret:

$$r^{(t)}(x) = \langle \ell(x^{(t)}), x^{(t)} - x \rangle$$

## Equilibrium

$$x^{(t)} \rightarrow \mathcal{N} \Leftrightarrow \limsup_t \sup_x r^{(t)}(x) \leq 0$$



# Regret

Average cumulative regret

$$R^{(t)}(x) = \frac{1}{t} \sum_{\tau \leq t} r^{(\tau)}(x)$$

## Equilibrium

$$\bar{x}^{(t)} = \frac{1}{t} \sum_{\tau \leq t} x^{(\tau)} \rightarrow \mathcal{N} \Leftrightarrow \limsup_t \sup_x R^{(t)}(x) \leq 0$$

By convexity of  $f$ ,

$$f\left(\frac{1}{t} \sum_{\tau \leq t} x^{(\tau)}\right) - f(x) \leq \frac{1}{t} \sum_{\tau \leq t} f(x^{(\tau)}) - f(x) \leq \frac{1}{t} \sum_{\tau \leq t} \langle \ell(x^{(t)}), x^{(t)} - x \rangle = R^{(t)}(x)$$

# Regret

- Regret first defined by Hannan (1957) in the context of repeated games [6]
- Large classes of algorithms have “no regret” guarantees, e.g. [3]
- However, **only guarantees convergence of  $\bar{x}^{(t)}$ , not  $x^{(t)}$**
- Seek additional conditions to guarantee  $x^{(t)} \rightarrow \mathcal{N}$ .

## Observation

If  $f(x^{(t)})$  is eventually monotone, then  $f(x^{(t)}) \rightarrow f^*$ .

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[6] James Hannan. Approximation to Bayes risk in repeated plays. *Contributions to the Theory of Games*, 3:97–139, 1957

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# Replicator dynamics

## Replicator equation

$$\forall p \in \mathcal{P}_k, \frac{dx_p^k}{dt} = x_p^k \left( \langle \ell_{\mathcal{P}_k}(x), x_{\mathcal{P}_k} \rangle - \ell_p^k(x) \right) \quad (1)$$

Also in evolutionary game theory, Weibull [10].

Theorem: Fischer and Vöcking [5]

Every solution of the ODE (1) converges to the set of its stationary points.

Proof:  $f$  is a Lyapunov function.

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# Approximate REplicator update

## Discretization of the continuous-time replicator dynamics

$$x_p^{(t+1)} - x_p^{(t)} = \eta_t x_p^{(t)} \left( \langle \ell^k(x^{(t)}), x_{\mathcal{P}_k}^{(t)} \rangle - \ell^k(x^{(t)}) \right) + \eta_t U_p^{(t+1)}$$

$(U^{(t)})_{t \geq 1}$  perturbations that satisfy for all  $T > 0$ ,

$$\lim_{\tau_1 \rightarrow \infty} \max_{\tau_2: \sum_{t=\tau_1}^{\tau_2} \eta_t < T} \left\| \sum_{t=\tau_1}^{\tau_2} \eta_t U^{(t+1)} \right\| = 0$$

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Michel Benaïm. [Dynamics of stochastic approximation algorithms](#).  
In *Séminaire de probabilités XXXIII*, pages 1–68. Springer, 1999

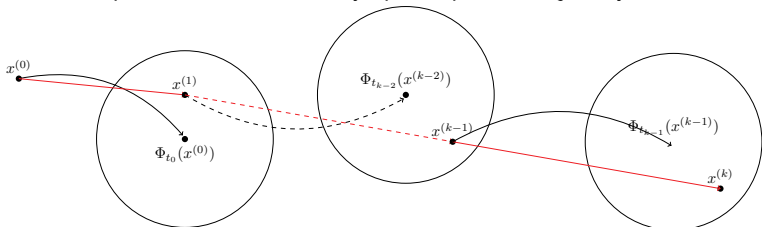
# Convergence to Nash equilibria

## Theorem Krichene et al. [7]

Under AREP updates, if  $\eta_t \downarrow 0$  and  $\sum \eta_t = \infty$ , then

$$x^{(t)} \rightarrow \mathcal{N}$$

- Affine interpolation of  $x^{(t)}$  is an asymptotic pseudo trajectory.



- $f$  is a Lyapunov function for Nash equilibria in the continuous system.

However, **No convergence rates.**

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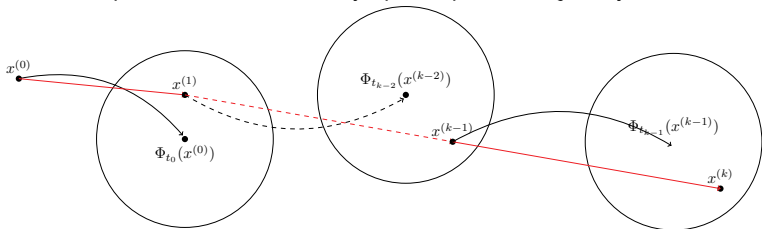
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# Stochastic Mirror Descent

minimize  $f(x)$  convex function  
subject to  $x \in \mathcal{X} \subset \mathbb{R}^d$  convex, compact set

At iteration  $t$

- have a stochastic subgradient  $\hat{g}^{(t)}$
- $\hat{g}^{(t)}$  unbiased:  $\mathbb{E} [\hat{g}^{(t)} | \mathcal{F}_t] = g^{(t)} \in \partial f(x^{(t)})$  ( $\mathcal{F}_t$  natural filtration of  $(x^{(t)})$ )

---

**Algorithm 1** SMD Method with learning rates  $(\eta_t)$

---

- 1: **for**  $t \in \mathbb{N}$  **do**
  - 2:      $\hat{g}^{(t)} \in \mathcal{F}_{t+1}$
  - 3:      $x^{(t+1)} = \arg \min_{x \in \mathcal{X}} \langle \hat{g}^{(t)}, x \rangle + \frac{1}{\eta_t} D_{\psi_t}(x, x^{(t)})$
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# Mirror Descent

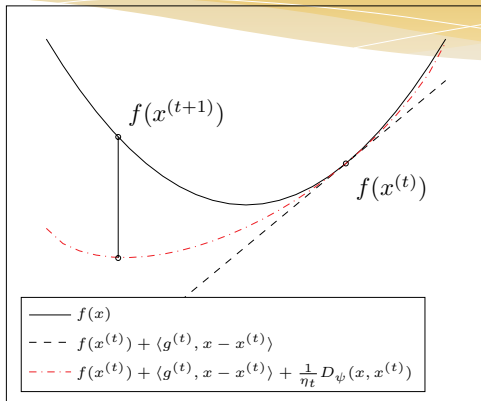


Figure: Mirror Descent iteration

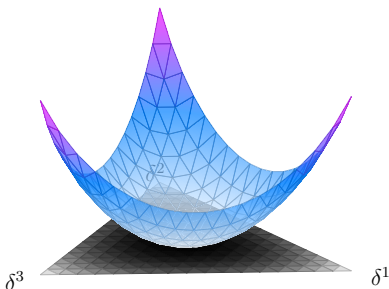
# Bregman Divergence

## Bregman Divergence

Strongly convex function  $\psi$

$$D_\psi(x, y) = \psi(x) - \psi(y) - \langle \nabla \psi(y), x - y \rangle$$

- $\psi(x) = \frac{1}{2}\|x\|_2^2$ ,  $D_\psi(x, y) = \frac{1}{2}\|x - y\|_2^2$  (projected gradient)
- $\psi(x) = -H(x) = \sum_{i=1}^d x_i \ln x_i$ ,  $D_\psi(x, y) = D_{KL}(x, y) = \sum_{i=1}^d x_i \ln \frac{x_i}{y_i}$ .



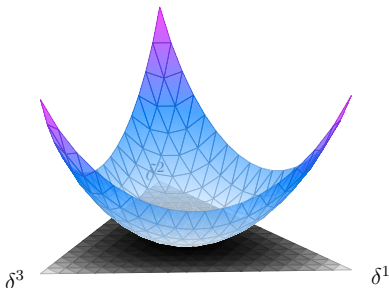
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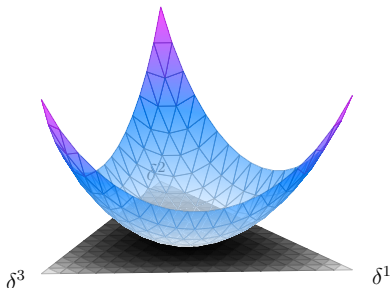
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# Convergence rates

$f$	$\eta_t$	Convergence
Weakly convex	$\frac{c}{\sqrt{t}}$ $\frac{c}{t^\alpha}, \alpha \in (0, 1)$	$\frac{1}{t} \sum_{\tau=1}^t \mathbb{E} [f(x^{(\tau)})] - f^* = O\left(\frac{1}{\sqrt{t}}\right)$ [8] $\mathbb{E} [f(x^{(t)})] - f^* = O\left(\frac{\log t}{t^{\min(\alpha, 1-\alpha)}}\right)$
Strongly convex	$\eta_t \rightarrow 0, \sum \eta_t = \infty$ $\frac{\theta}{\ell_f t^\alpha}, \alpha \in (0, 1)$	$\mathbb{E} [D_\psi(x^*, x^{(t)})] = O\left(\eta_T + e^{-\sum_T^t \eta_\tau}\right)$ $\mathbb{E} [D_\psi(x^*, x^{(t)})] = O(t^{-\alpha})$

**Figure:** Convergence rates of SMD. S. Krichene, W. Krichene, R. Dong, A. Bayen. In preparation.

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[8] A. S. Nemirovsky and D. B. Yudin. *Problem complexity and method efficiency in optimization*.

Wiley-Interscience series in discrete mathematics. Wiley, 1983.

ISBN 9780471103455

# Regret bound for SMD

Main ingredient:

## Proposition

Assume  $D_\psi$  bounded by  $D$  and  $\mathbb{E} \|\hat{g}\|^2 \leq G$ .  
SMD method with  $(\eta_t)$ .  $\forall t_2 > t_1 \geq 0$  and  $\mathcal{F}_{t_1}$ -measurable  $x$ ,

$$\sum_{\tau=t_1}^{t_2} \mathbb{E} \left[ \langle g^{(\tau)}, x^{(\tau)} - x \rangle \right] \leq \frac{\mathbb{E} [D_\psi(x, x^{(t_1)})]}{\eta_{t_1}} + D \left( \frac{1}{\eta_{t_2}} - \frac{1}{\eta_{t_1}} \right) + \frac{G}{2\ell_\psi} \sum_{\tau=t_1}^{t_2} \eta_\tau \quad (2)$$

# Distributed SMD with heterogeneous agents

- $\mathcal{X} = \mathcal{X}_1 \times \cdots \times \mathcal{X}_K$
- Agent  $k$  updates  $x^k \in \mathcal{X}_k$
- $D_{\psi^k}$  and  $\eta_t^k$  depends on  $k$

---

**Algorithm 4** DSMD Method with learning rates  $(\eta_t^k)$  and divergences  $D_{\psi^k}$

---

- 1: **for**  $t \in \mathbb{N}$  **do**
- 2:      $\hat{g}^{(t)} \in \mathcal{F}_{t+1}$
- 3:

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# Routing game with heterogeneous populations

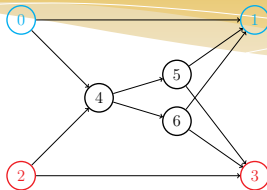


Figure: Example network

## Routing game with heterogeneous populations

Under **unbiased noisy losses**, with **heterogeneous update rules** with  $\eta_t^k = \theta_k t^{-\alpha_k}$

$$\mathbb{E} \left[ f(x^{(t)}) \right] - f^* = O \left( t^{-\min(\min_k \alpha_k, 1 - \max_k \alpha_k)} \right)$$

where  $f$  is the Rosenthal potential function

# Simulations

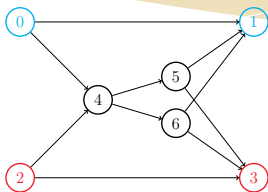


Figure: Example network

- Centered Gaussian noise on edges.
- Population 1: Hedge with  $\eta_t^1 = t^{-0.1}$
- Population 2: Hedge with  $\eta_t^2 = \frac{1}{2}t^{-0.5}$

## One realization

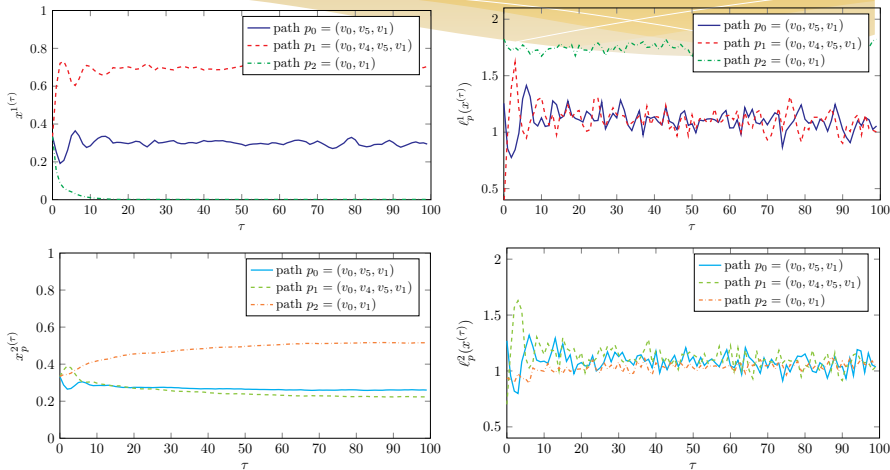


Figure: Population distributions and noisy path losses

# In Expectation

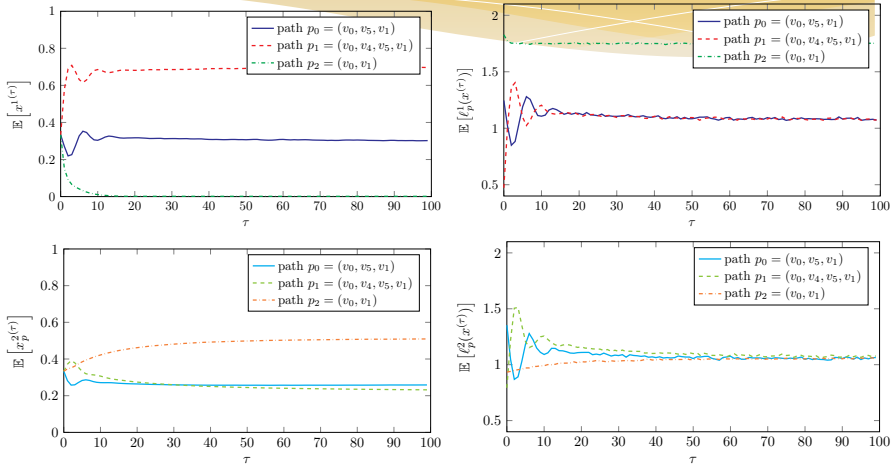


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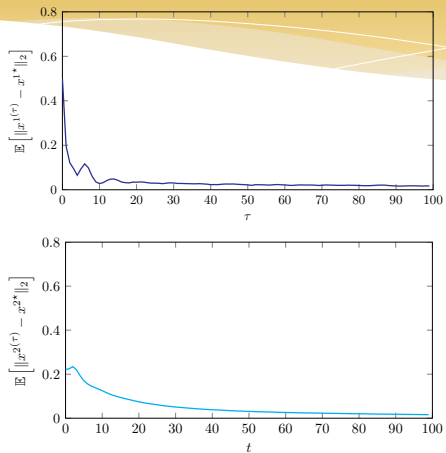


Figure: Expected distance to equilibrium

## Summary and Extensions

- How do players arrive at equilibrium?  
Any algorithm in the AREP or the DSMD class.
- How fast?  
Convergence rates for the DSMD class.
- Stability?  
Nash equilibria are stable for these dynamics [4]
- Robustness?  
Robust to unbiased perturbation, e.g. when losses are not known but estimated.

### Extensions

- Provides a model of population dynamics for optimal control problems.
- Adapt to other problems, such as network consensus.

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In *Conference on Decision and Control (CDC)*, 2014

## Summary and Extensions

- How do players arrive at equilibrium?  
Any algorithm in the AREP or the DSMD class.
- How fast?  
Convergence rates for the DSMD class.
- Stability?  
Nash equilibria are stable for these dynamics [4]
- Robustness?  
Robust to unbiased perturbation, e.g. when losses are not known but estimated.

### Extensions

- Provides a model of population dynamics for optimal control problems.
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Thank you.

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