

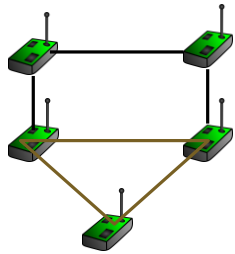


Resilient Consensus Through Trusted Nodes

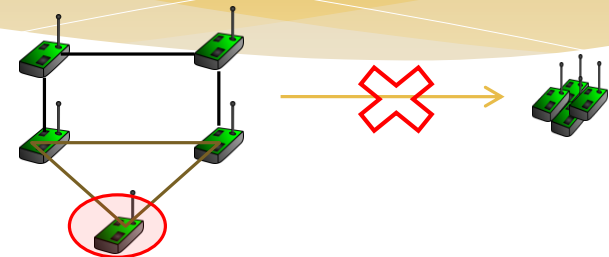
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Motivation



Consensus: agents agree upon a common value
Applications: optimization, sensor networks, estimation, etc.



Resilient consensus: consensus in the presence of adversaries

Existing Resilient Consensus Protocols

Non-local information

e.g., node knows the overall network topology

High connectivity

e.g.,
- $(2F+1)$ -connected
- (r,s) -robust

Conditions hard to verify

e.g., (r,s) -robustness

Dependence on parameter F
(no. of possible attacks)

e.g., F -Total, F -local etc.

Can these issues be addressed?

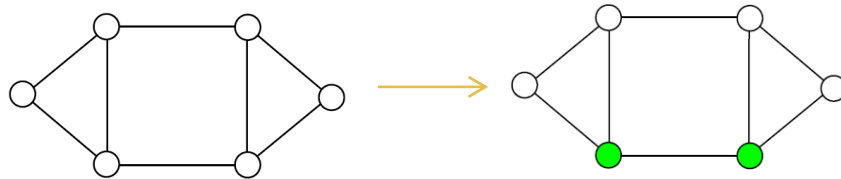
Motivation

Resilient networks
(various approaches to resilience)

**Connectivity of
underlying graph**

**Nodes' capabilities
(security levels)**

By **how much** can we improve the **resilience** of a network by adding **few trusted nodes**?



some nodes are **more secured** against adversarial attacks

Trusted Nodes

- Highly secure.
- Cannot be compromised by attacks.

Resilient Consensus with Trusted Nodes

Design an update rule such that all normal nodes, a subset of which consists of **trusted nodes**, achieve consensus even in the presence of **any** number of adversaries.

i.e., design an update rule that achieves consensus.

Consensus:

- Agreement
(as $k \rightarrow \infty$, $x_i(k) = x_j(k)$)
- Safety
($\forall k$, $x_i(k) \in [x_{\min}(0) \quad x_{\max}(0)]$)

Trusted nodes:

For any no. of adversarial attacks, find **necessary and sufficient** conditions on the

- number
- location
- connectivity

of trusted nodes to achieve consensus.

Resilient Consensus Protocol (RCP-T)

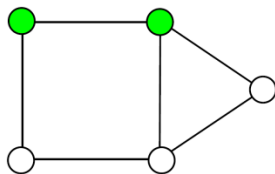
Under RCP-T, consensus is always achieved in the presence of *arbitrary number of adversaries* if and only if there exists a set of trusted nodes that form a **connected dominating set**.

Under RCP-T

- Any number of attacks can be handled
- Sparse networks can be made resilient

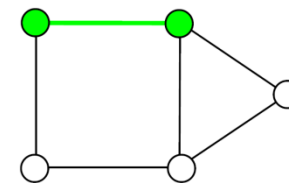
Dominating Set:

$$D \subseteq V, \quad \text{s.t.} \quad \bigcup_{v_i \in D} \mathcal{N}[v_i] = V$$



Connected Dominating Set:

Nodes in the dominating set induce a **connected** subgraph

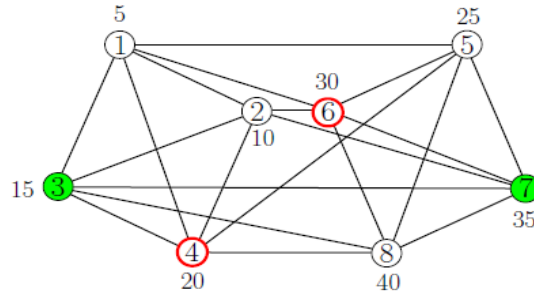


Example

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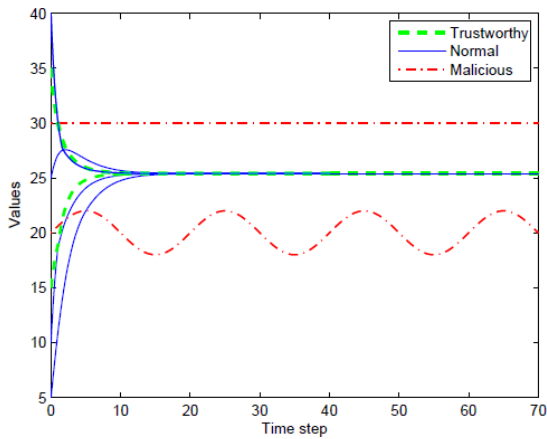
Attacked nodes = {4, 6}

Trusted nodes = {3, 7}

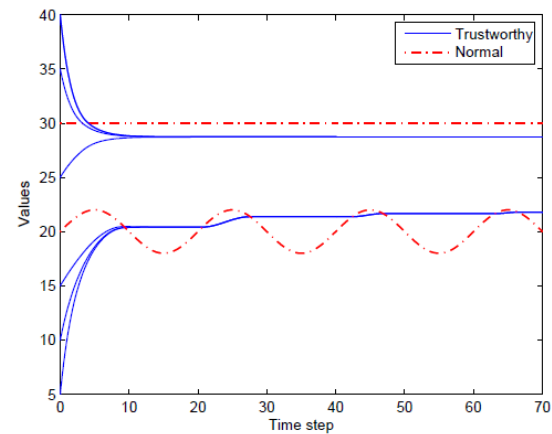


(2,2)-robust graph

RCP-T



W-MSR



- RCP-T achieves consensus even with two attacks
- W-MSR algorithm can handle a single attack, but not two adversaries

Conclusions

Trusted nodes and Network Robustness

γ_c = Connected domination number

Resilience of networks if
No. of trusted nodes $< \gamma_c$

Sometimes adding as many as $(\gamma_c - 1)$ trusted nodes does not improve the resilience.

$\gamma_c \sim (r, s)$ – robust graphs

‘Trusted nodes’ ?

How can we **generalize** the notion of ‘trusted nodes’ for resilience in networks?

Graph domination

How can we utilize the concept of **domination in graphs** for resource distribution in networks.