Learning Dynamics, Estimation and Control In Congestion Games

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- Transportation, communication networks
- Nash equilibrium quantifies efficiency of network in steady state.



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Congestion Games

Estimation of Hedge Dynamics

Control under Hedge dynamics

Loss Function Imputation



Congestion games

Routing game

- Player k drives from source to destination node
- Chooses path from \mathcal{A}_k
- Mass of players on each edge determines cost on that edge.

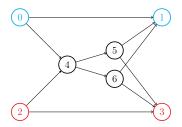


Figure : Routing game

[3]Walid Krichene, Benjamin Drighès, learning in selfish routing. In 31st International Conference on Machine Learning (ICML). JMLR, 2014

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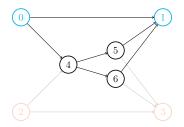


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Algorithm 1 *

- 1: for $t \in \mathbb{N}$ do 2: Play $p \sim x_{\mu}^{(t)}$
- 3: Discover $\ell_{L}^{(t)}$
- 4: Update

$$\begin{aligned} x_k^{(t+1)} &= \operatorname*{arg\,min}_{x \in \Delta^{\mathcal{A}_k}} \left\langle \ell_k^{(t)}, x_k \right\rangle + \frac{1}{\eta_k^{(t)}} D_{\mathsf{KL}}(x_k^{(t)}, x_k) \\ &\propto e^{-\eta_k^{(t)} \ell_k^{(t)}} \end{aligned}$$

5: end for

Online Learning Model



[4]Walid Krichene, Syrine Krichene, and Alexandre Bayen. Efficient Bregman projections onto the simplex. In 54th IEEE Conference on Decision and Control (CDC), Osaka, Japan, 2015

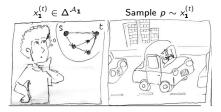
Algorithm 2 *

- 1: for $t \in \mathbb{N}$ do 2: Play $p \sim x_k^{(t)}$
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- 4: Update

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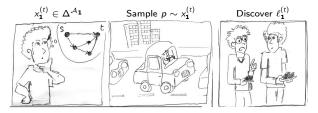
Algorithm 3 *

- 1: for $t \in \mathbb{N}$ do 2: Play $p \sim x_k^{(t)}$
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- 4: Update

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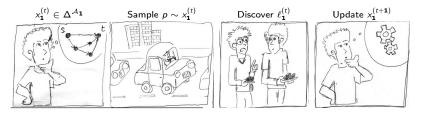
Algorithm 4 *

- 1: for $t \in \mathbb{N}$ do Play $p \sim x_k^{(t)}$ 2:
- 3: Discover $\ell_{\mu}^{(t)}$
- Update 4:

$$\begin{aligned} \mathbf{x}_{k}^{(t+1)} &= \operatorname*{arg\,min}_{x \in \Delta^{\mathcal{A}_{k}}} \left\langle \ell_{k}^{(t)}, \mathbf{x}_{k} \right\rangle + \frac{1}{\eta_{k}^{(t)}} D_{\mathbf{KL}}(\mathbf{x}_{k}^{(t)}, \mathbf{x}_{k}) \\ &\propto e^{-\eta_{k}^{(t)} \ell_{k}^{(t)}} \end{aligned}$$

5: end for

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Time remaining: 0 Turns left: 14 Logged as: Cacao -Logout Path's previous Current Previous Path cost Weight Flows Flows 0.5 Path 1.466 0.250 0.000 0.5 Path 1.263 Ó. 0.250 0.242 0.5 Path 1.374 0.250 0.000 0.5 Path 1.352 0.250 0.758 · Start Node: 3 · Destination Node: 10 · Total Flow: 1.0 Previous Cost Total Cost 1.900 -1.850 -1.850 -1.750 -1.750 -1.550 -1.550 -1.550 -1.550 -1.550 -1.450 -1.450 -1.450 -1.450 -1.450 -1.450 -1.450 -1.5 2.000 2.400 -2.200 -2.000 -1,800 0 1.573 1.000 -1 1.342 1.400 -1.200 -2 1,443 1.000 -**3** 1.312 0.800 -total_cost 1.100 1.000 -0.800 0.700 -0.600 0.500 D.400 0.300 0.200 p.100 -0.000 -



- We observe a sequence of player decisions $(\bar{x}^{(t)})$ and losses $(\bar{\ell}^{(t)})$.
- Can we fit a model of player dynamics?

[6]Kiet Lam, Walid Krichene, and Alexandre Bayer, Ondearning how players learn: Estimation of learning dynamics in the routing game resulting of the sector of the sector

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Mirror descent model

Estimate the learning rate in the mirror descent model

$$x^{(t+1)}(\eta) = \operatorname*{arg\,min}_{x\in\Delta^{\mathcal{A}_k}} \left\langle ar{\ell}^{(t)}, x
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angle + rac{1}{\eta} D_{\mathcal{KL}}(x,ar{x}^{(t)})$$

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Then $d(\eta) = D_{KL}(\bar{x}^{(t+1)}, x^{(t+1)}(\eta))$ is a convex function. Can minimize it to estimate $\eta_k^{(t)}$.

[6]Kiet Lam, Walid Krichene, and Alexandre Bayen On Gearning how players learn: Estimation of learning dynamics in the routing game presented at the second second

In 7th International Conference on Cyber-Physical Systems (ICCPS), 2016

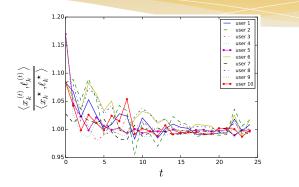


Figure : Costs of each player (normalized by the equilibrium cost)



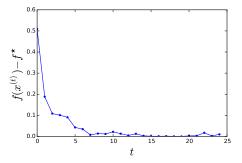


Figure : Potential function $f(x^{(t)}) - f^*$.



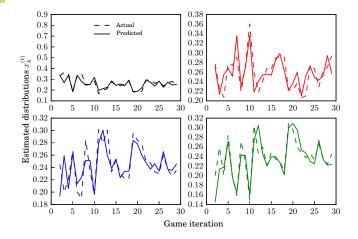


Figure : Estimated Vs. actual distribution.



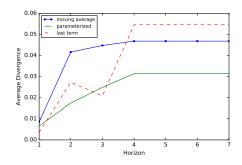


Figure : Average KL divergence between predicted distributions and actual distributions, as a function of the prediction horizon h.





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Optimal routing with learning dynamics

Assumptions

- A central authority has control over a fraction of traffic: $u^{(t)} \in \alpha_1 \Delta^{A_1} \times \cdots \times \alpha_K \Delta^{A_K}$
- ► Remaining traffic follows learning dynamics: $x^{(t)} \in (1 - \alpha_1)\Delta^{\mathcal{A}_1} \times \cdots \times (1 - \alpha_K)\Delta^{\mathcal{A}_K}$

Optimal routing under selfish learning constraints

minimize_{u(1:7),x(1:7)}
$$\sum_{t=1}^{T} J(x^{(t)}, u^{(t)})$$
subject to
$$x^{(t+1)} = u(x^{(t)} + u^{(t)}, \ell(x^{(t)} + u^{(t)}))$$

[5]Walid Krichene, Milena Suarez, and response. Transactions on Control of Networked Systems (TCNS), 2017

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Solution methods

 Greedy method: Approximate the problem with a sequence of convex problems.

minimize_{$$u(t)$$} $J(u(x^{(t-1)}, u^{(t-1)}), u^{(t)})$



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Mirror descent with the adjoint method.

Adjoint method

minimize_{*u*} J(u, x)subject to H(x, u) = 0

equivalent to

minimize J(u, X(u))

Then perform mirror descent on this function of u.



A simple example

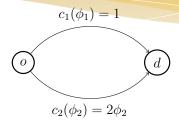


Figure : Simple Pigou network used for the numerical experiment.

- Social optimum: $\left(\frac{3}{4}, \frac{1}{4}\right)$
- Nash equilibrium $(\frac{1}{2}, \frac{1}{2})$
- Control over $\alpha = \frac{1}{2}$ of traffic



A simple example

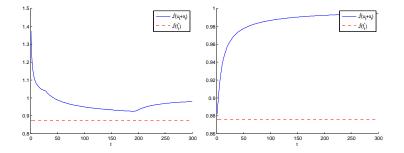


Figure : Social cost $J^{(t)}$ over time induced by adjoint solution (left) and the greedy solution (right). The dashed line shows the social optimal allocation.



A simple example

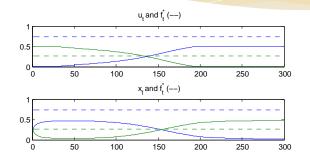


Figure : Adjoint controlled flows (top), selfish flows (bottom). The green lines correspond to the top path, and the blue lines to the bottom path. The dashed lines show the social optimal flows $x^{SO} = (\frac{3}{4}, \frac{1}{4})$.



Application to the L.A. highway network

- Simplified model of the L.A. highway network.
- Cost functions uses the B.P.R. function, calibrated using the work of [8].

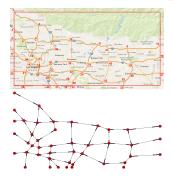


Figure : Los Angeles highway network.

^[8]J. Thai, R. Hariss, and A. Bayen. A miltiple approach to latency inference and control in traffic equilibria from sparse data. In American Control Conference (ACC), 2015, pages 689–695, July 2015

Application to the L.A. highway network

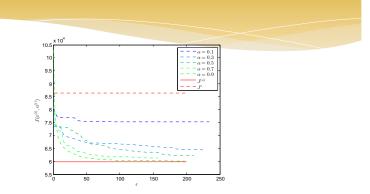


Figure : Average delay without control (dashed), with full control (solid), and different values of $\alpha.$





Congestion Games

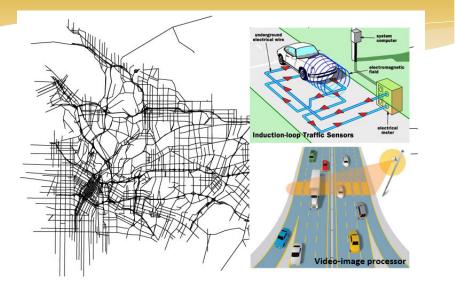
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Motivation





Framework

 $VI(\mathcal{K}, F)$ consists in finding $\mathbf{x} \in \mathcal{K}$, *i.e.*, $A\mathbf{x} = \mathbf{b}$, $\mathbf{x} \ge 0$, such that $F(\mathbf{x})^T (\mathbf{by} - \mathbf{x}) \ge 0, \forall \mathbf{by} \in \mathcal{K}$

Examples of parametric VI's $VI(\mathcal{K}(\mathbf{p}), F(\cdot, \mathbf{p}))$:

▶ Routing game with fixed latency functions and variable demand:

 $\ell(\mathbf{x})^{\mathsf{T}}(\mathsf{by} - \mathbf{x}) \ge 0 \quad \forall \mathsf{ feasible path flow by for demand } \mathbf{d}(\mathbf{p})$ (2)

Parametric convex optimization:

min $f(\mathbf{x}, \mathbf{p})$ s.t. $\mathbf{x} \in \mathcal{K}(\mathbf{p}) \iff VI(\mathcal{K}(\mathbf{p}), \nabla f(\cdot, \mathbf{p}))$ (3)

• Controller fitting, consumer behavior etc. [2, 1]

[2]A. Keshavarz, Y. Wang, and S. Boyd. Imputing a convex opjective function. In *IEEE International Symposium on Intelligent Control (ISIC)*, 2011
[1]Dimitris Bertsimas, Vishal Gupta, an equilibrium using inverse optimization. *Math. Program.*, pages 595–633, 2015 (1)

Problem statement

Inputs:

- parametric polyhedron $\{\mathcal{K}(\mathbf{p})\}_{\mathbf{p}}$
- ▶ parametric observation process $g(\cdot, \mathbf{p}) : \mathbb{R}^n \to \mathbb{R}^q$
- ▶ *N* observations of equilibria $\mathbf{z}^{(j)} := g(\mathbf{x}^{(j)}, \mathbf{p}^{(j)}) + \mathbf{w}^{(j)}, j \in [N]$

Objective:

Impute parametric map $F(\cdot, \mathbf{p})$ and decision vectors $\mathbf{x}^{(j)}$ such that

- (a) $\mathbf{x}^{(j)}$ is an approximate solution to $VI(\mathcal{K}(\mathbf{p}^{(j)}), F(\mathbf{p}^{(j)}))$.
- (b) $\mathbf{x}^{(j)}$ agrees with the observations $\mathbf{z}^{(j)}$.

Method:

The idea is to find $F(\cdot, \mathbf{p})$ and $\mathbf{x}^{(j)}$ that minimize both

- (a) sum of sub-optimality gaps r_{eq} of the VI's
- (b) observation residual $r_{obs} := \sum_{j=1}^{N} \phi\left(g(\mathbf{x}^{(j)}, \mathbf{p}^{(j)}), \mathbf{z}^{(j)}\right)$



Formalization

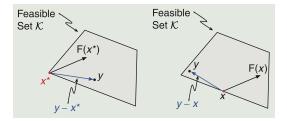
Define the optimality gap for $VI(\mathcal{K}, F)$:

$$r_{VI}(\mathbf{x}) = \max_{by \in \mathcal{K}} F(\mathbf{x})^{\mathsf{T}}(\mathbf{x} - by)$$
(4)

Note that

►
$$r_{VI}(\mathbf{x}) \ge 0, \forall \mathbf{x} \in \mathcal{K}$$

► $r_{VI}(\mathbf{x}) \le \epsilon \iff \min_{b\mathbf{y} \in \mathcal{K}} F(\mathbf{x})^T (by - \mathbf{x}) \ge -\epsilon$





Residual of the primal-dual system

Define residual function for $\mathbf{x} \in \mathcal{K}$ and by $\in \mathbb{R}^n$

$$r_{\text{PD}}(\mathbf{x}, \mathbf{by}) = F(\mathbf{x})^T \mathbf{x} - \mathbf{b}^T \mathbf{by}$$

Theorem [1]

The following holds for any $\epsilon \geq 0$ and $\mathbf{x} \in \mathcal{K}$

$$r_{VI}(\mathbf{x}) \leq \epsilon \iff \exists by \in \mathbb{R}^n : \mathbf{A}^T by \leq F(\mathbf{x}), r_{PD}(\mathbf{x}, by) \leq \epsilon$$
 (6)

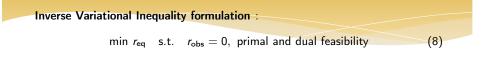
Set
$$r_{eq} = \sum_{j} r_{PD}(\mathbf{x}^{(j)}, by^{(j)}, \mathbf{p}^{(j)}) = \sum_{j} F(\mathbf{x}^{(j)}, \mathbf{p}^{(j)})^{\mathsf{T}} \mathbf{x}^{(j)} - \mathbf{b}^{\mathsf{T}}(\mathbf{p}^{(j)}) by^{(j)}$$

and solve:

$$\begin{array}{ll} \min_{\substack{F,\mathbf{x},\ by}} & w_{eq} r_{eq} + w_{obs} r_{obs} \\ \text{s.t.} & \mathbf{x}^{(j)} \in \mathcal{K}(\mathbf{p}^{(j)}), \quad \mathbf{A}(\mathbf{p}^{(j)})^{\mathsf{T}} \text{ by }^{(j)} \leq F(\mathbf{x}^{(j)}, \mathbf{p}^{(j)}), \quad \forall j \end{array}$$
(7)

[1]Dimitris Bertsimas, Vishal Gupta, and partis Paschatidis. Data-driven estimation in equilibrium using inverse optimization. Math. Program., pages 595–633, 2015 (5)

Previous works



- assumes complete and noiseless observations
- not robust to measurement errors

Bilevel programming formulation:

$$\min r_{\rm obs} \quad \text{s.t.} \quad r_{\rm eq} = 0 \tag{9}$$

- difficult to solve due to bilevel structure
- KKT system as constraints pose numerical difficulties (ref)

[1]Dimitris Bertsimas, Vishal Gupta, and the prise paschatidis. Data-driven estimation in equilibrium using inverse optimization. Math. Program., pages 595–633, 2015

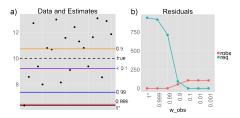
Pareto optimization

Solve the Weighted Sum Program (WSP) :

min $w_{eq}r_{eq} + w_{obs}r_{obs}$ s.t. primal and dual feasible

Algorithm 1 Weighted-sum() Weighted sum method

- 1: Normalize objectives $r_{\rm eq}$ and $r_{\rm obs}$ for consistent comparisons.
- 2: Solve with $w_{obs} + w_{eq} = 1$ and $w_{obs} \in \{0.001, 0.01, 0.1, 0.9, 0.99, 0.999\}$
- 3: Check the values of $r_{\rm eq}$ and $r_{\rm obs}$.
- Robust to noise or outliers with appropriate choice of ϕ in r_{obs}
- smoothing with penalization r_{PD} instead of constraint $r_{PD} = 0$



[7]R.T. Marler and J.S. Arora. Survey and the presence of the

(10)

Asymptotic behavior

Let r_{eq}^{\star} be the optimal obj. value of the inverse VI {min r_{eq} s.t. $r_{obs} = 0$ }

Let r_{obs}^{\star} be the optimal obj. value of the BP {min r_{obs} s.t. $r_{eq} = 0$ }

Let the WSP {min $w_{eq}r_{eq} + w_{obs}r_{obs}$ s.t. feasibility} with $w_{eq} + w_{obs} = 1$.

Theorem [9]

Any optimal solution $u^{\star} \in \mathcal{S}(w_{\mathsf{eq}}, w_{\mathsf{obs}})$ to the WSP is such that

$$r_{
m obs}(\mathbf{u}^{\star}) \leq r_{
m eq}^{\star}(w_{
m obs}^{-1}-1) \approx r_{
m eq}^{\star}(1-w_{
m obs}) \text{ as } w_{
m obs} \stackrel{\leq}{\longrightarrow} 1$$
 (11)

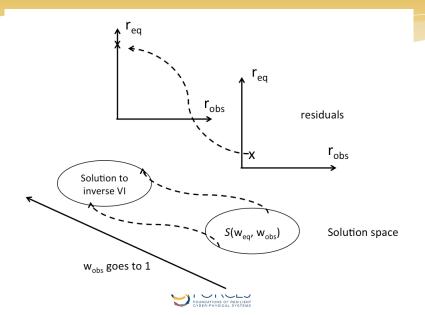
$$r_{\rm eq}(\mathbf{u}^{\star}) \xrightarrow{\leq} r_{\rm eq}^{\star}$$
 (uniformly) (12)

Given compactness and any sequence $w_{obs}^{(n)} \longrightarrow 1$, there exists a sequence $\{\mathbf{u}^{(n)}\}_n \in \{\mathcal{S}(w_{eq}^{(n)}, w_{obs}^{(n)})\}_n$ and a sub-sequence of it converging to a solution to the inverse VI.

[9]Jerome Thai and Alexandre Bayen. function: a robust approach. Journal of Mathematical Analysis and Applications, 2016

Asymptotic behavior

Opens a new dimension for which the 'edge' are the inverse VI and the BP.



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- [1] Dimitris Bertsimas, Vishal Gupta, and Ioannis Paschalidis. Data-driven estimation in equilibrium using inverse optimization. *Math. Program.*, pages 595–633, 2015.
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Thank you!

