

Strategic Network Inspection using Resource-Constrained sUAS

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Massachusetts Institute of Technology

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How to operationalize network sensing strategies?

For a given network that faces adversarial disruptions, design and operationalize (randomized) sensing strategies subject to limitations on sensing range and resource constraints.



Malicious attacks



Randomized defense

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Approach

 Formulate a robust optimization problem over the network.



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Malicious attacks



- Formulate a robust optimization problem over the network.
 - Defender: chooses a dispatch of sUAS.



Randomized defense



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- Main contributions



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- Main contributions
 - General sensing model: heterogeneous range.



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- Formulate a robust optimization problem over the network.
 - Defender: chooses a dispatch of sUAS.
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- Main contributions
 - General sensing model: heterogeneous range.
 - Solution approach using combinatorial problems.

M. Dahan, L. Sela, S. Amin. "Randomized Network Sensing under Strategic Disruptions", Working paper





Malicious attacks



Randomized defense

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- Dispatch of sUAS
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- (Q) How to allocate a fleet of sUAS for network inspection in an adversarial environment?





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- N: Set of locations that can be visited by an sUAS.



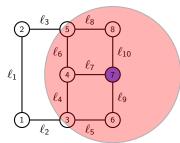
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- ▶ For every location $i \in N$, $C_i \in 2^E$ represents the subset of components that an sUAS is capable of monitoring when positioned in location *i*. For example, C_i may represent:

L. Sela, W. Abbas, X. Koutsoukos, and S. Amin. "Sensor placement for fault location identification in water networks: a minimum test cover approach", *Automatica*, 2016



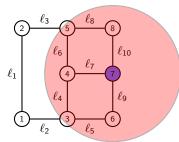
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 - The components that are within a certain distance from *i*.

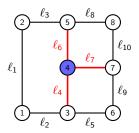


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 - The components that are within a certain distance from i.
 - The adjacent edges of node i.





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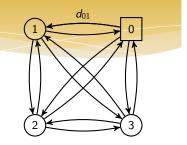
O ∈ N: Unique base node from where the sUAS are sent.



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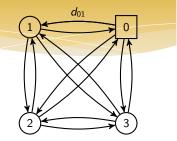
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- For every pair of locations (i, j) ∈ N², let d_{ij} denote the distance to fly from i to j.
 - The d_{ij} can take into account air space restrictions, obstacles, height difference between locations, etc.



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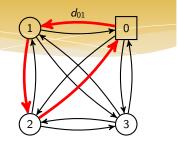
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▶ Homogeneous fuel-constrained sUAS that can fly for up to *D_{max}* miles before going back to the base node 0.



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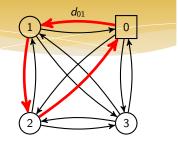
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Feasible Flight Plan

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- Set of feasible flight plans:

$$\mathcal{F} := \{(i_1, \dots, i_m) \in N^m \mid i_1 = i_m = 0 ext{ and } \sum_{k=1}^{m-1} d_{i_k i_{k+1}} \leq D_{max}, \ m \in \mathbb{N} \}.$$



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Robust optimization problem

Minimize the maximum number of failure events that remain undetected:

$$(\mathcal{P}_{\textit{insp}}) \qquad \min_{\sigma^1 \in \Delta(\mathcal{A}_1)} \max_{\mu \in \mathcal{A}_2} \mathbb{E}_{\sigma^1} \left[|\mu| - |\mathcal{C}_\eta \cap \mu|
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▶ $|\mu| - |C_{\eta} \cap \mu|$ is the total number of failures net the number of detected failures.



Auxiliary Problem



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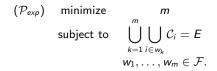
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- m^* : Optimal value of (\mathcal{P}_{exp}) .
- Can be formulated as a mixed-integer program.



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Detection Guarantees

Theorem

Given an optimal solution of (\mathcal{P}_{exp}) , we can construct a randomized strategy $\tilde{\sigma}^1$ such that:

M. Dahan, A. Weinert, and S. Amin. "Network Exploration and Inspection Using Distance-Constrained sUAS", *Submitted*, 2016



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b₁: number of available sUAS

b₂: maximum attack size

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▶ *m*^{*}: Optimal value of (*P*_{exp})

M. Dahan, A. Weinert, and S. Amin. "Network Exploration and Inspection Using Distance-Constrained sUAS", *Submitted*, <u>2016</u>



Detection Guarantees

Theorem

Given an optimal solution of (\mathcal{P}_{exp}) , we can construct a randomized strategy $\tilde{\sigma}^1$ such that:

1. The expected number of undetections in the worst case is upper bounded by:

$$\max_{\mu \in \mathcal{A}_2} \mathbb{E}_{\widetilde{\sigma}^1} \left[|\mu| - |\mathcal{C}_\eta \cap \mu|
ight] \leq b_2 \left(1 - rac{b_1}{m^*}
ight).$$

2. The detection rate, defined as the ratio between the number of detections and the total number of failure events, in the worst case, is lower bounded in expectation by:

$$\min_{\mu \in \mathcal{A}_2} \mathbb{E}_{\widetilde{\sigma}^1} \left[\frac{|\mathcal{C}_{\eta} \cap \mu|}{|\mu|} \right] \geq \frac{b_1}{m^*}$$

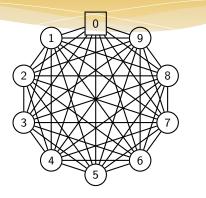
- b1: number of available sUAS
 b2: maximum attack size
- *m*^{*}: Optimal value of (*P_{exp}*)

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Case Study: Complete Network

- Fully connected network.
- 10 locations uniformly placed on a circle of radius 1 mile.
- The sUAS can travel for 4 miles.
- Vulnerable components are the network edges that can be monitored from its end nodes.

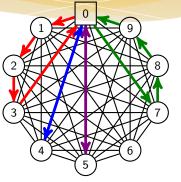


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Case Study: Complete Network

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• Optimal solution of (\mathcal{P}_{exp}) : $w_1^* = (0, 1, 2, 3, 0)$, $w_2^* = (0, 4, 0)$, $w_3^* = (0, 5, 0)$ and $w_4^* = (0, 7, 8, 9, 0)$.





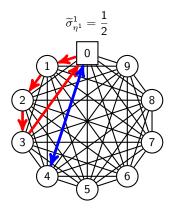
• If the operator has 2 sUAS, then $\tilde{\sigma}^1$ is illustrated as follows:

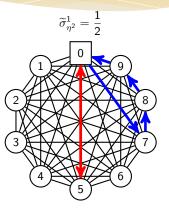


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Inspection

• If the operator has 2 sUAS, then $\tilde{\sigma}^1$ is illustrated as follows:

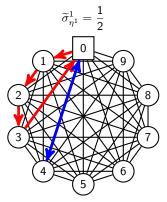


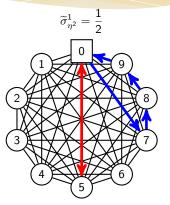




Inspection

• If the operator has 2 sUAS, then $\tilde{\sigma}^1$ is illustrated as follows:





▶ At least 50 % of the failures will be detected.



Conclusion

- Summary
 - Resource allocation problem for network inspection using fuel-constrained sUAS.
 - Flexible model that can take into account constraints imposed by the sUAS platform and the environment.
 - Mixed-integer programming formulation for the network exploration problem.
 - Extension to the inspection problem, and performance guarantee on the detection score in worst-case scenarios.



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 - Flexible model that can take into account constraints imposed by the sUAS platform and the environment.
 - Mixed-integer programming formulation for the network exploration problem.
 - Extension to the inspection problem, and performance guarantee on the detection score in worst-case scenarios.
- Future Work
 - Include heterogeneity in the vulnerability or importance of components.
 - Account for imperfect (and noisy) information on network state in designing exploration/inspection strategies.





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- 2. MIT Thurber Fellowship

Thank you!

Questions: mdahan@mit.edu



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