



A Game-Theoretic Approach for Alert Prioritization in Cyber-Physical Systems

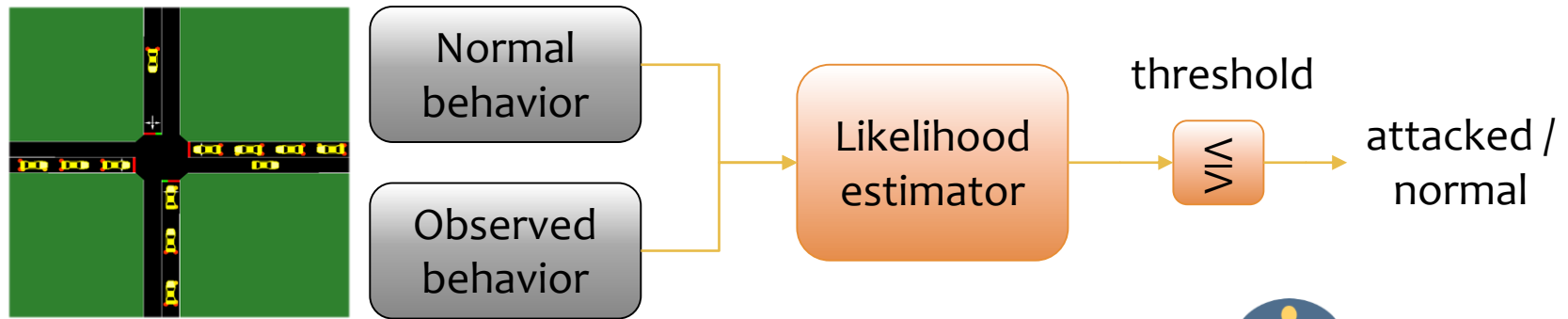
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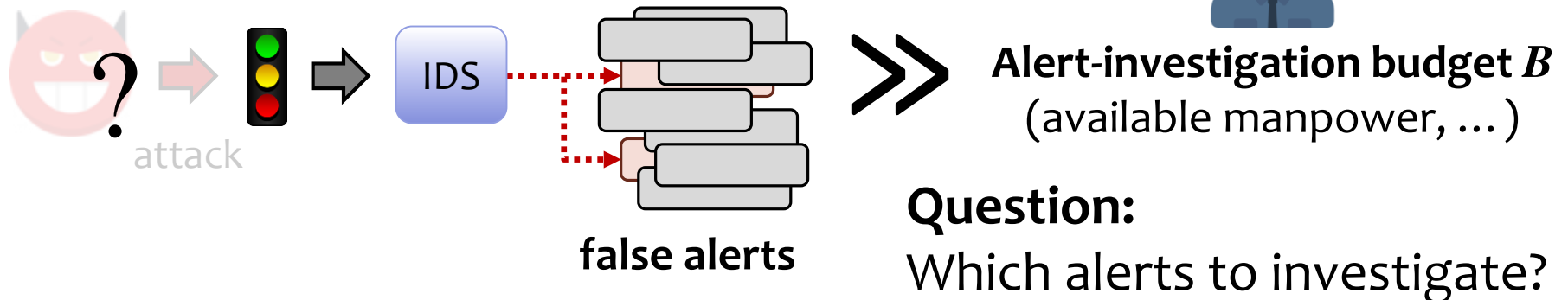


Physical Anomaly Based Intrusion Detection

- * Resilience to novel threats through anomaly-based detection
- * Previously: traffic-anomaly based attack detection



* Problem

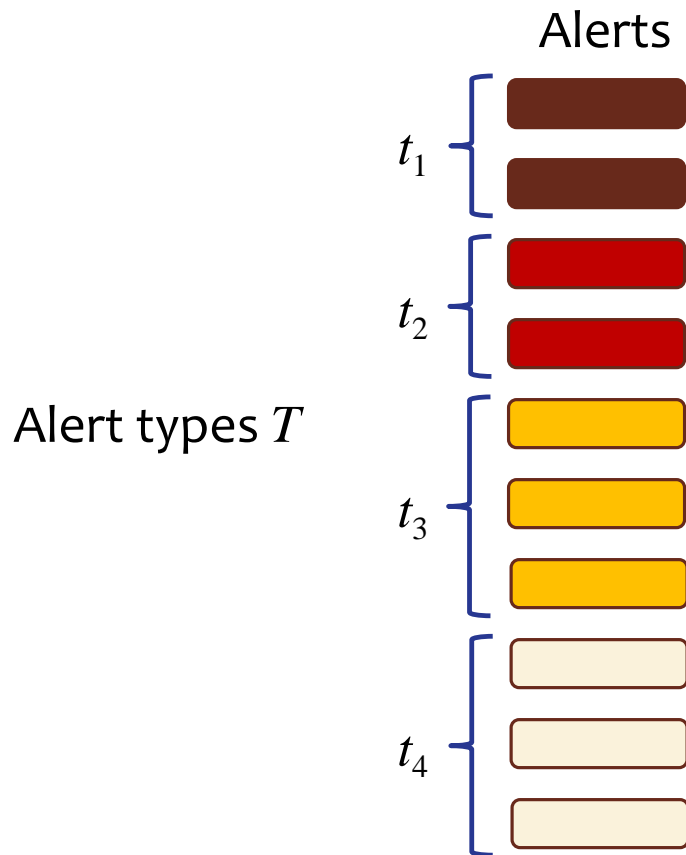


Alert Prioritization

Alerts





Alert Prioritization



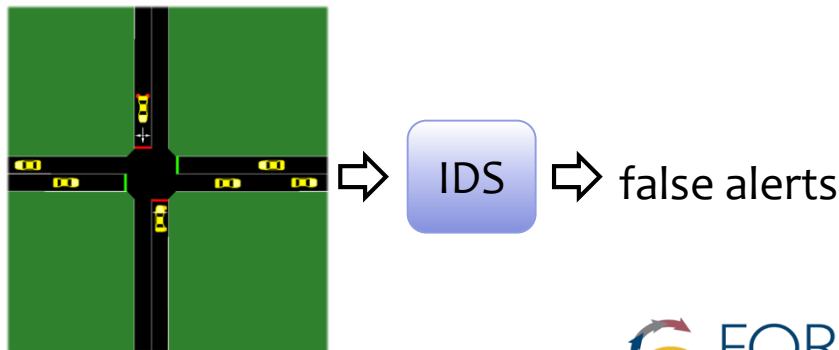
Alert Types

* Alert types T

* example:

		anomaly magnitude	
		low	high
anomalous system	sensing 	t_4	t_2
	control 	t_3	t_1

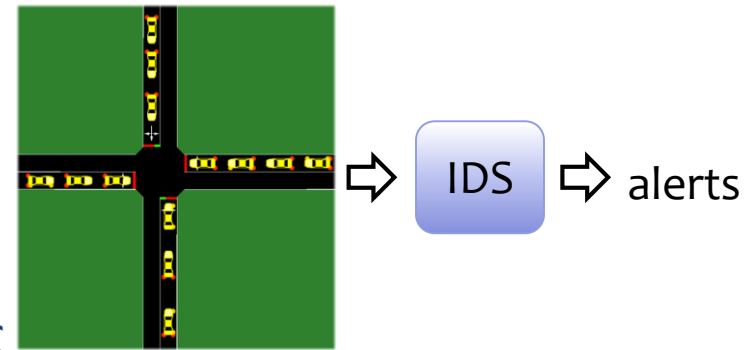
* **cumulative distribution F_t** of the number of false alerts of type t :
simulate normal behavior



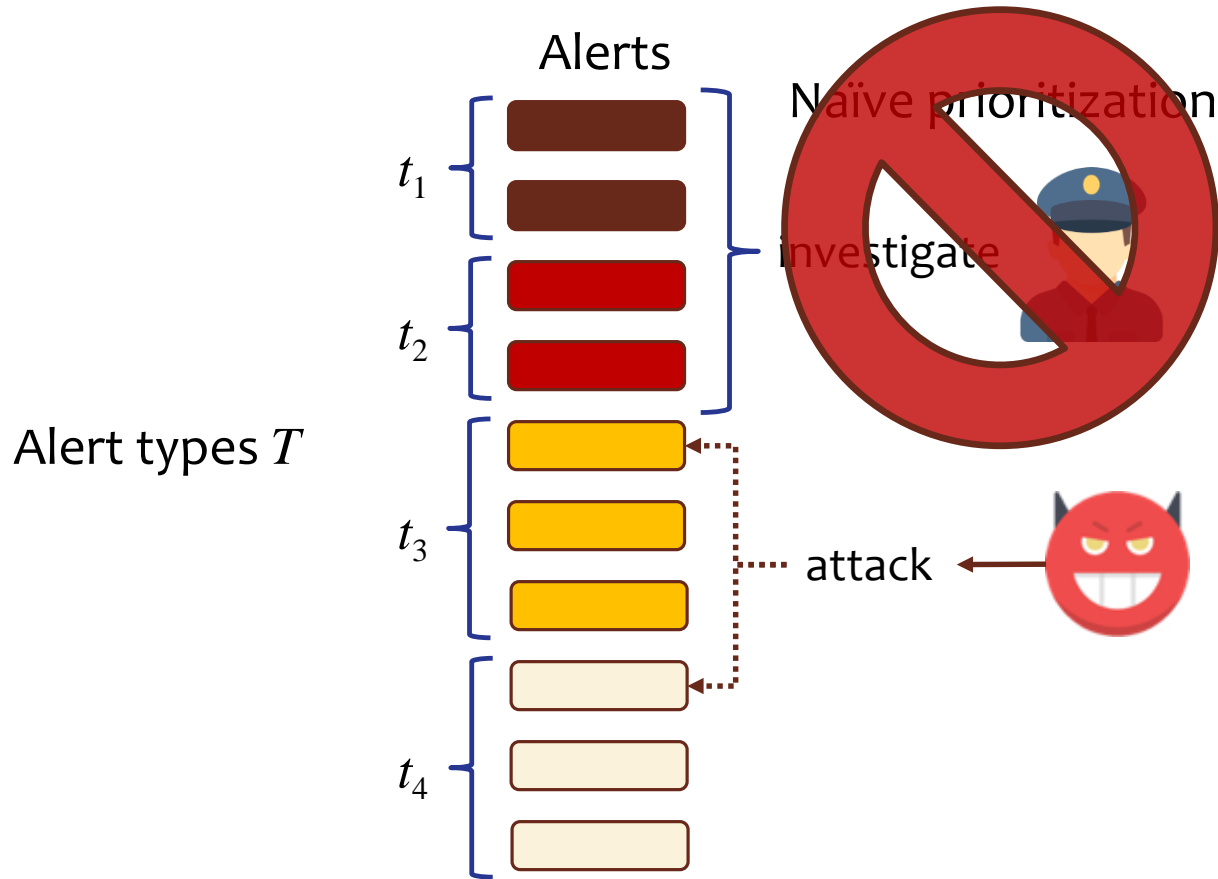
* Attacks A

- * example: attacked devices, integrity / availability
- * **impact L_a of attack a :**
our previous work (ICCPs'16)
- * **probability $R_{a,t}$ of raising an alert of type t for attack a :**

simulate attack a



Alert Prioritization Problem



Stackelberg Security Game

* Players



1. Defender

- selects alert prioritization strategy p , which is a probability distribution over orderings of T



2. Adversary

- selects an attack a from the set of possible attacks A

* Payoffs

- * let $PD(o, a)$ be probability of investigating attack a using ordering o

- * defender's expected loss = $\sum_{o \in O} p_o \cdot (1 - PD(o, a)) \cdot G_a - K_a$

- * adversary's expected payoff = $\sum_{o \in O} p_o \cdot (1 - PD(o, a)) \cdot L_a$

* Solution concept

- * adversary's best response: $BR(p) = \operatorname{argmax}_{a \in A} \sum_{o \in O} p_o \cdot (1 - PD(o, a)) \cdot L_a$

- * **optimal prioritization strategy:** $\min_{p, a \in BR(p)} \sum_{o \in O} p_o \cdot (1 - PD(o, a)) \cdot G_a - K_a$

(details can be found in our paper
published at AAAI-17 AICS)

Computational Results

* Computing probabilities

exponential number of terms

* definition: $PD(\mathbf{o}, a) = \sum_{\hat{T} \subseteq T} \prod_{t \in \hat{T}} R_{a,t} \prod_{t \in T \setminus \hat{T}} (1 - R_{a,t}) PI(\mathbf{o}, \min\{t \mid o_t \in \hat{T}\})$

$$PI(\mathbf{o}, k) = \sum_{\substack{\mathbf{n}: \\ C_{o_k} + \sum_{i=1}^k n_i \cdot C_{o_i} \leq B}} (F_{o_k}^*(n_k) - F_{o_k}^*(n_k - 1)) \cdot \prod_{i=1}^{k-1} (F_{o_i}(n_i) - F_{o_i}(n_i - 1))$$

where C_t is investigation cost for type t

* dynamic programming algorithm:

Algorithm 1 Computing $PD(\mathbf{o}, a)$

Input: prioritization game, prioritization \mathbf{o} , attack a

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1: for  $b = 0, 1, \dots, B$  do
2:    $PD(\mathbf{o}, a, |T|, b) \leftarrow R_{a, o_{|T|}} \cdot F_{o_{|T|}}^*(\lfloor b/C_{o_{|T|}} \rfloor - 1)$ 
3: end for
4: for  $i = |T| - 1, \dots, 2, 1$  do
5:   for  $b = 0, 1, \dots, B$  do
6:      $PD(\mathbf{o}, a, i, b) \leftarrow R_{a, o_i} \cdot F_{o_i}^*(\lfloor b/C_{o_i} \rfloor - 1) + (1 - R_{a, o_i}) \sum_{j=0}^{\lfloor b/C_{o_i} \rfloor} (F_{o_i}(j) - F_{o_i}(j - 1)) \cdot PD(\mathbf{o}, a, b - j \cdot C_{o_i}, i + 1)$ 
7:   end for
8: end for
9: Return  $PD(\mathbf{o}, a) := PD(\mathbf{o}, a, 1, B)$ 

```

polynomial time

Theorem: Finding an optimal prioritization strategy is an NP-hard problem.

Column Generation Algorithm

- * Optimal strategy: linear programming based formulation

- * for each attack $a \in A$:

$$\max_p \sum_{o \in O} p_o \cdot PD(o, a)$$

subject to

$$\forall a' \in A: \sum_{o \in O} p_o \cdot D(o, a') \geq \Delta(K_{a'})$$

$o \in O$

where

$$D(o, a') = [(1 - PD(o, a))G_a - (1 - PD(o, a'))G_{a'}]$$

$$\Delta(K_{a'}) = K_a - K_{a'}$$

exponential number of possible orderings

- * Polynomial-time column-generation approach

Algorithm 2 Greedy Column Generation

Input: prioritization game, reduced cost function \bar{c}

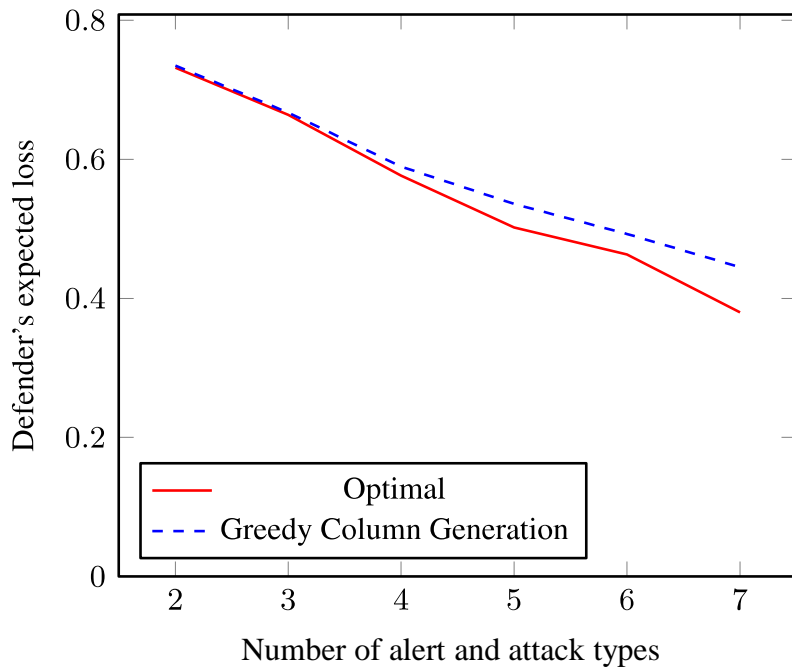
- 1: $o \leftarrow \emptyset$
 - 2: **while** $\exists t \in T \setminus o$ **do**
 - 3: $o \leftarrow o + \operatorname{argmax}_{t \in T \setminus o} \bar{c}(o + t)$
 - 4: **end while**
 - 5: Return o
-

reduced cost function:

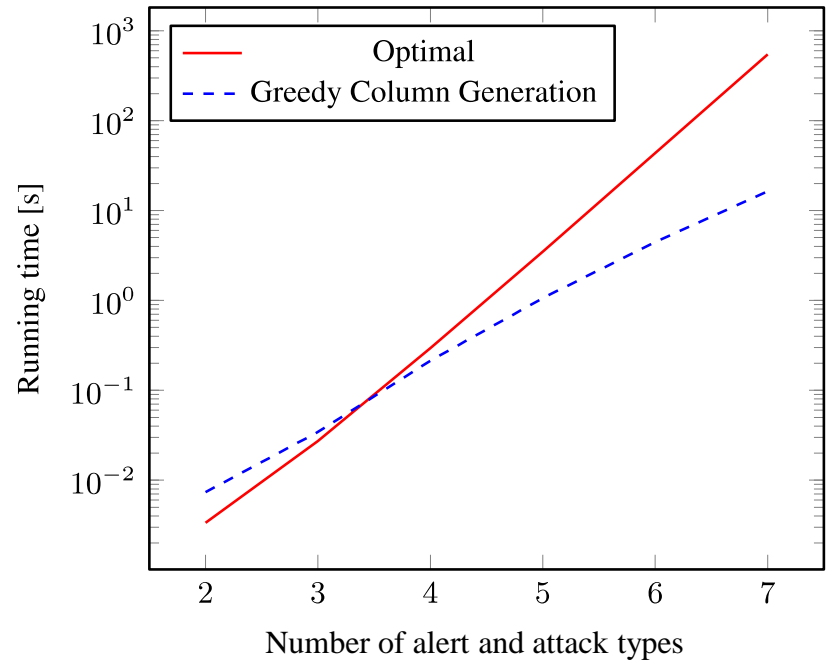
$$\bar{c}(o) = PD(o, a) + \sum_{a' \in A} y(\bar{O}, a') D(o, a')$$

Numerical Results

Expected Loss



Running Time



$K_a = 0$, $C_t = 1$, D_a and G_a are drawn at random from $[0.5, 1]$, $B = 5|T|$, each $R_{a,t}$ is either 0 (with probability $1/3$) or drawn at random from $[0, 1]$, every F_t is a Poisson whose mean is drawn at random from $[5, 15]$

Conclusion

- * Anomaly-based intrusion detection for cyber-physical systems
→ prohibitively high number of false alerts
- * Alert prioritization:
deciding which alerts to investigate with a limited budget
 - * naïve strategies are very vulnerable
- * Game-theoretic formulation and analysis
 - * finding optimal prioritization strategy is computationally hard
 - * polynomial-time dynamic-programming algorithm for computing detection probabilities
 - * polynomial-time column-generation approach for alert prioritization

Thank you for your attention!
Questions?