



System-Security Co-design

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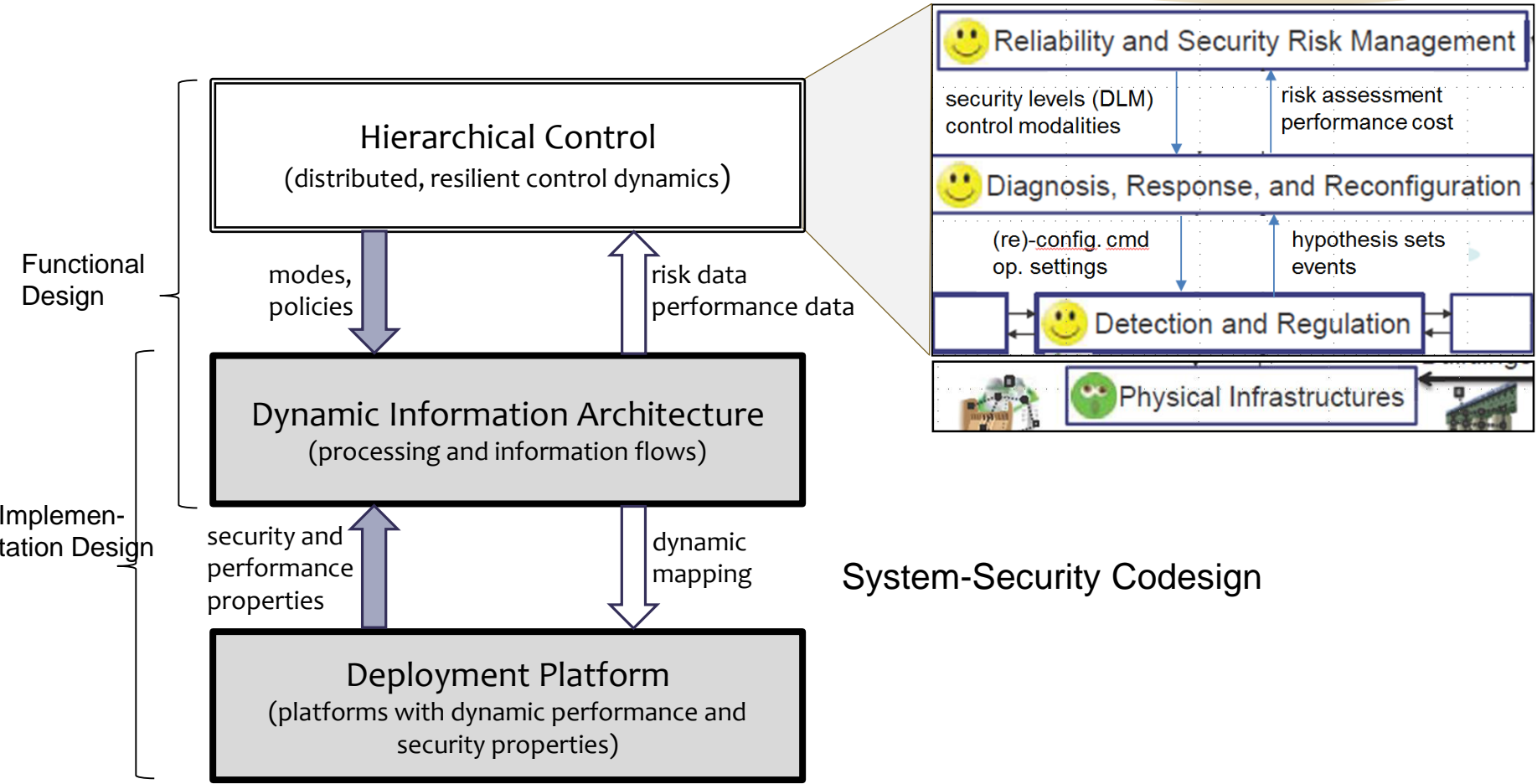
ISIS-Vanderbilt



Objectives of Reconfiguration

- * **Change modes of operation of Detection and Regulation**
 - Diagnosis, Response and Reconfiguration forms a supervisory control mechanism – used in hierarchical control approaches (e.g. Pappas, Tabuada)
- * **Re-synthesize implementation architecture**
 - Provide interface for changing required security policies
 - Provide models of information flows required to be implemented
 - Provide models for security and performance characteristics of communication links and computing devices
 - Provide precise specification for the reconfiguration space
 - Develop methods for remapping the information architecture to the implementation architecture subject to functional, performance, timing and security constraints

Co-design Problem

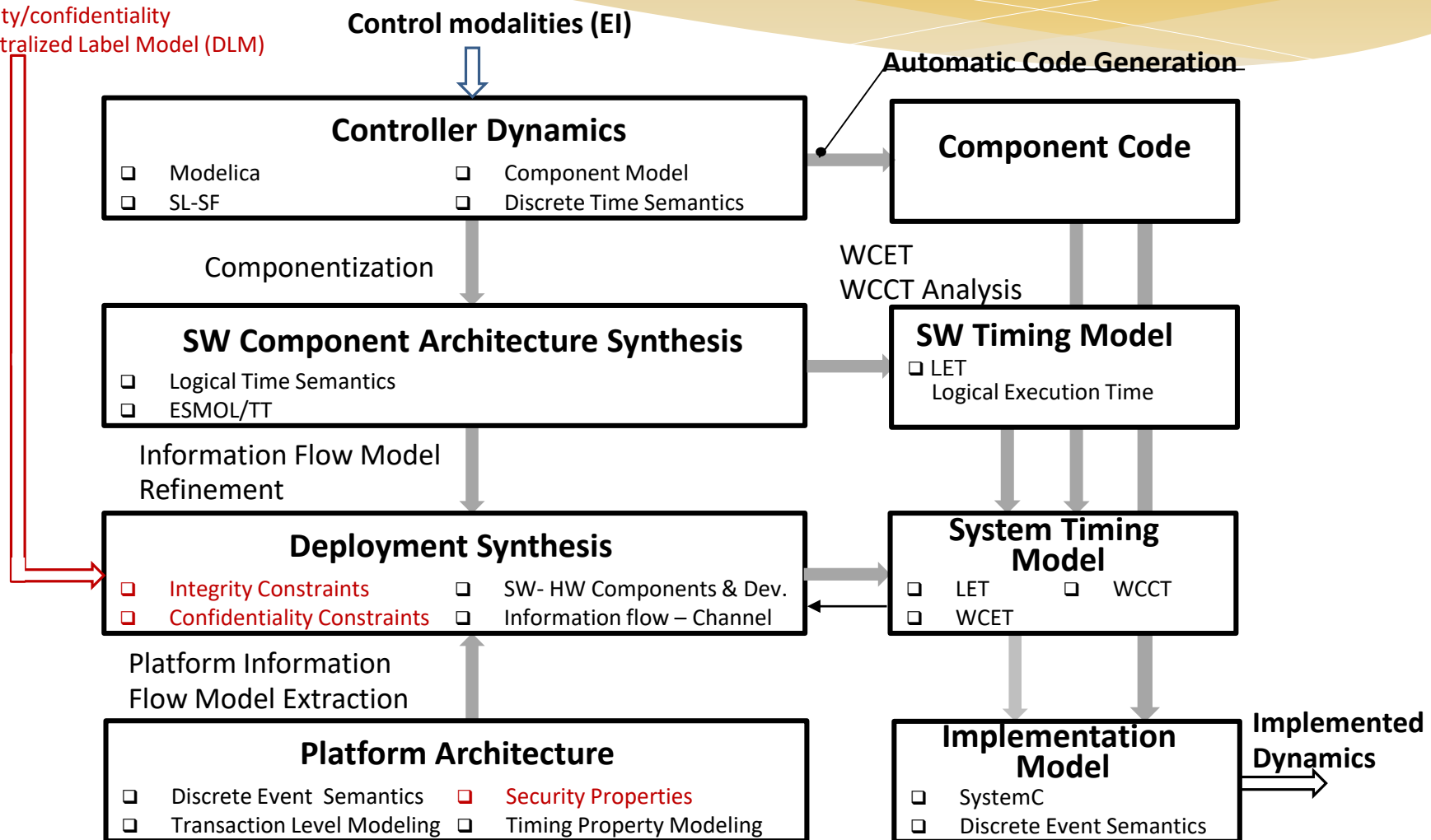


System – Security Co-design

Security Policies (EI)

integrity/confidentiality

Decentralized Label Model (DLM)



Security Concerns Addressed

- * **Integrity attacks**

- Manipulate data (value, timestamp, source identity,..)

- * **Confidentiality attack**

- Leak critical data to unauthorized persons/systems

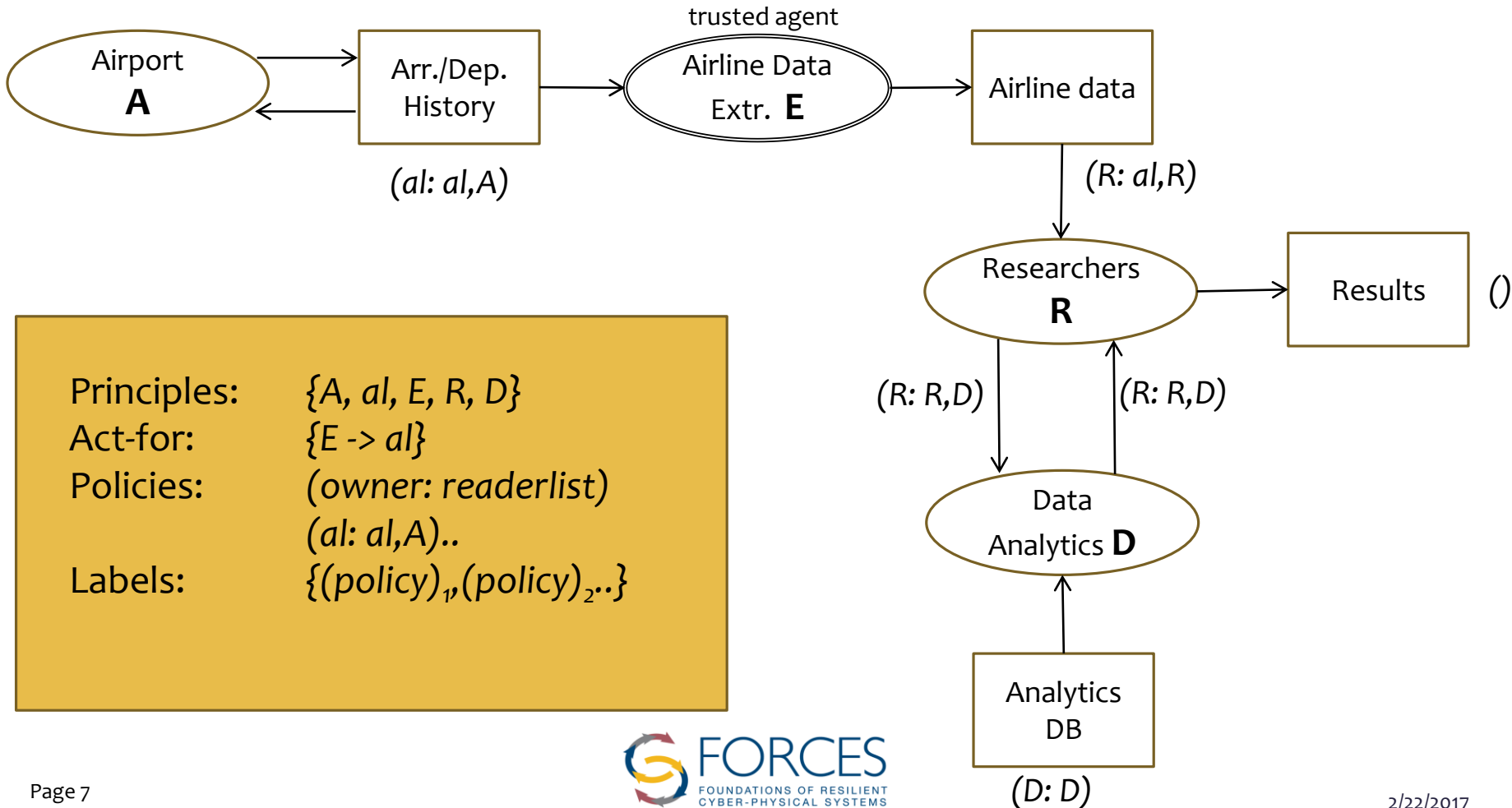
Integrity and confidentiality restrictions impose constraints on information flows.

- How to model these restrictions?
- How to integrate these restrictions with others (functional and timing) and formulate a co-design problem?

Decentralized Label Model for Information Flow Control

- * Myers, Liskov (1997): Introduced security-typed languages by labeling variables with information flow security policies
- * New semantic concepts:
 - *Principles* that represent authority entities.
 - *Labels* expressing security classes encountered in most information flow models.
 - *Policies* that are elementary security primitives used in *labels*.
 - *Labeled entities* that have attached labels, such as *values*, *slots* (*variables*, *objects*, *i/o channels*). Copies of *values* can be relabeled, *slots* cannot.
 - *Operators* that can *relabel* or *declassify* values in information flows.
- * DLM provides mechanism for static/dynamic type checking of security labels in information flows to detect policy violations.
- * Example: *Jif*, a security-typed version of Java

Simple Example

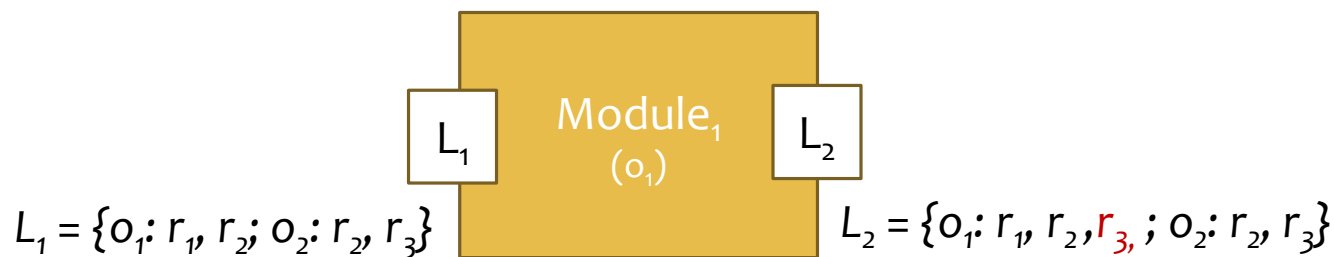


Working With Security Labels

- * Labels contain a set of policies. Each policy includes an owner and a set of readers allowed by the owner. The effective reader set for a label is the intersection of every reader set in it.

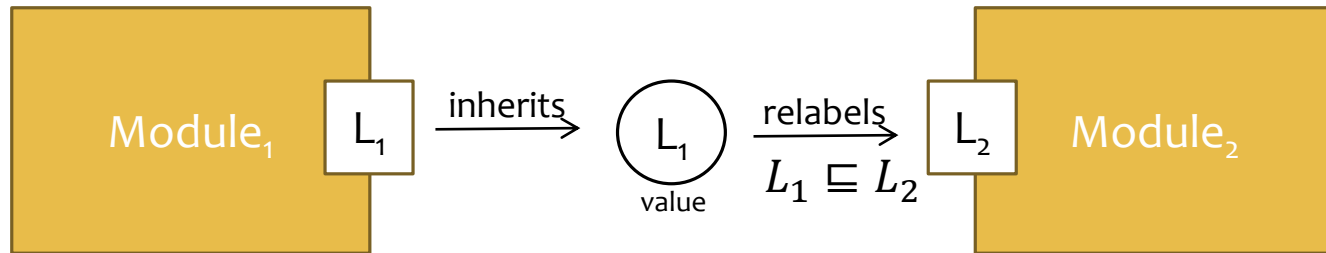
$$L = \{o_1: r_1, r_2; o_2: r_2, r_3\}$$

- * Processing blocks running under the authority of an owner can **declassify** the owner's policy by adding readers.



Propagation Rules

* Propagation rule-1:

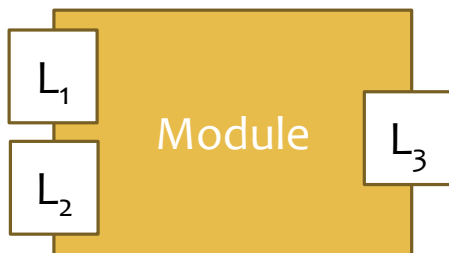


$$owners(L_1) \subseteq owners(L_2)$$

$$\forall o \in owners(L_1), readers(L_1, o) \supseteq readers(L_2, o)$$

(Labels form a security lattice.)

* Propagation rule-2:



L_3 is the join of L_1 and L_2

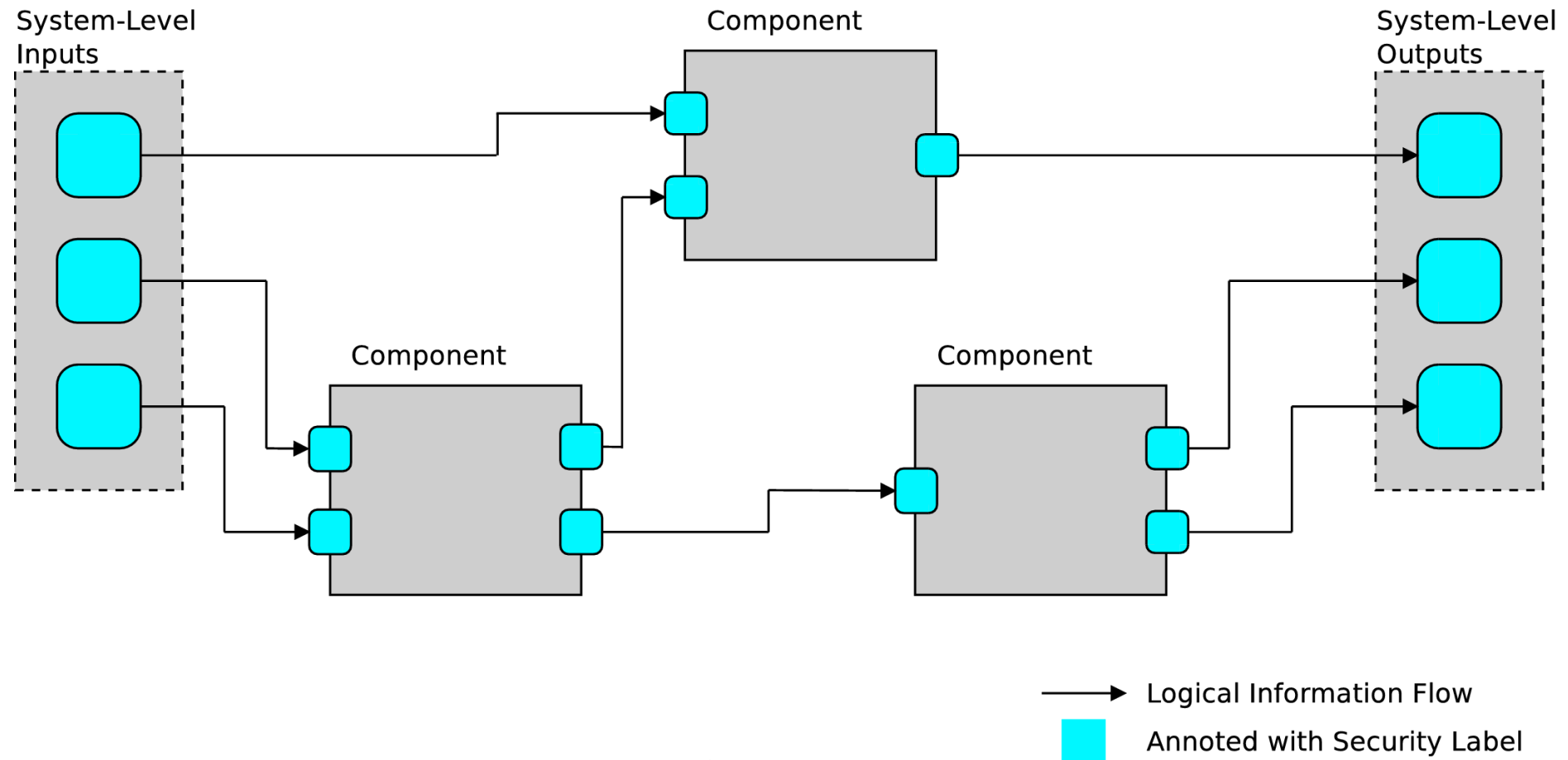
$$L_3 = L_1 \sqcup L_2$$

$$owners(L_1 \sqcup L_2) = owners(L_1) \cup owners(L_2)$$

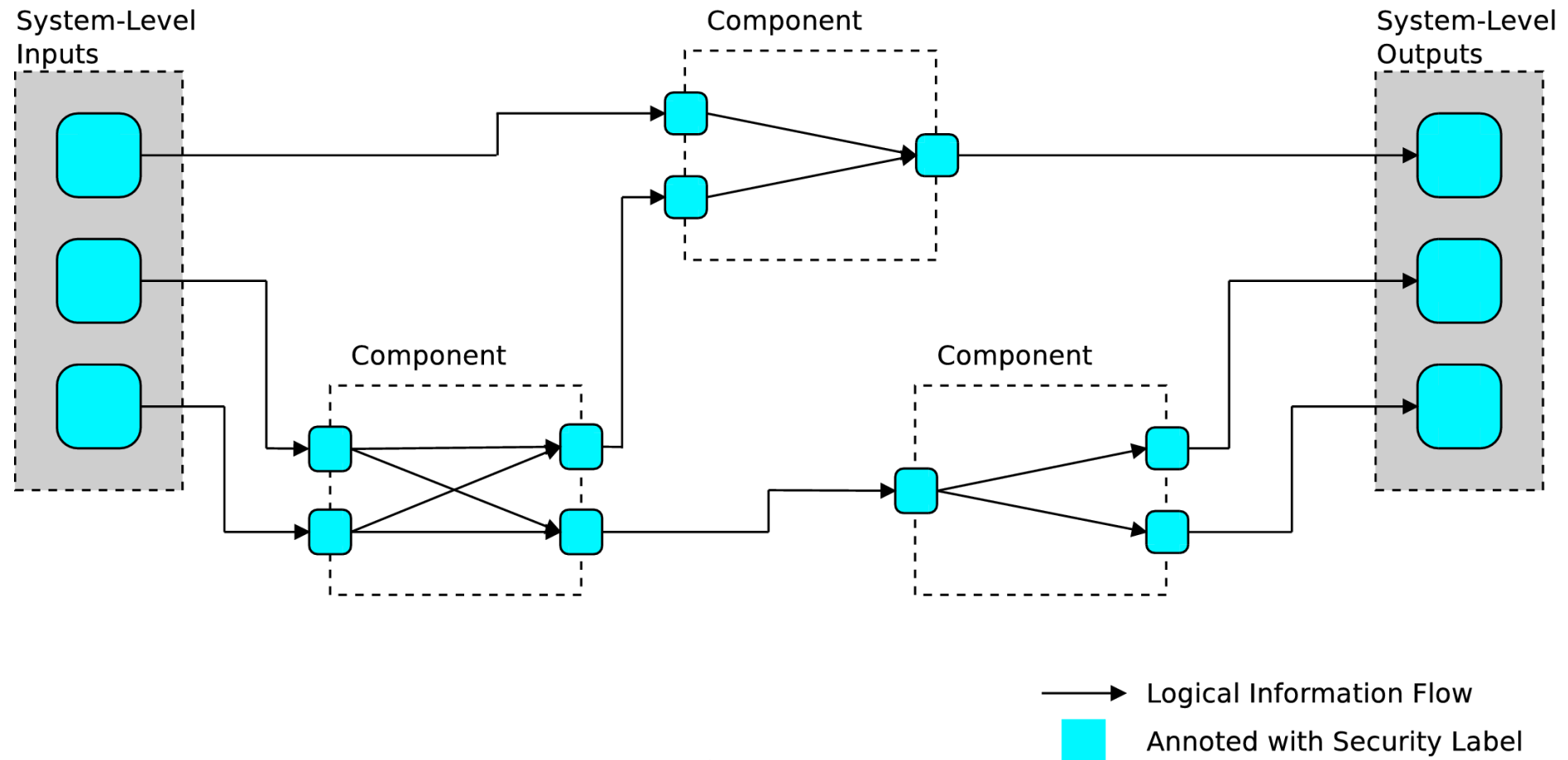
$$readers(L_1 \sqcup L_2, o) = readers(L_1, o) \cap readers(L_2, o)$$

DLM in Model-Based Design

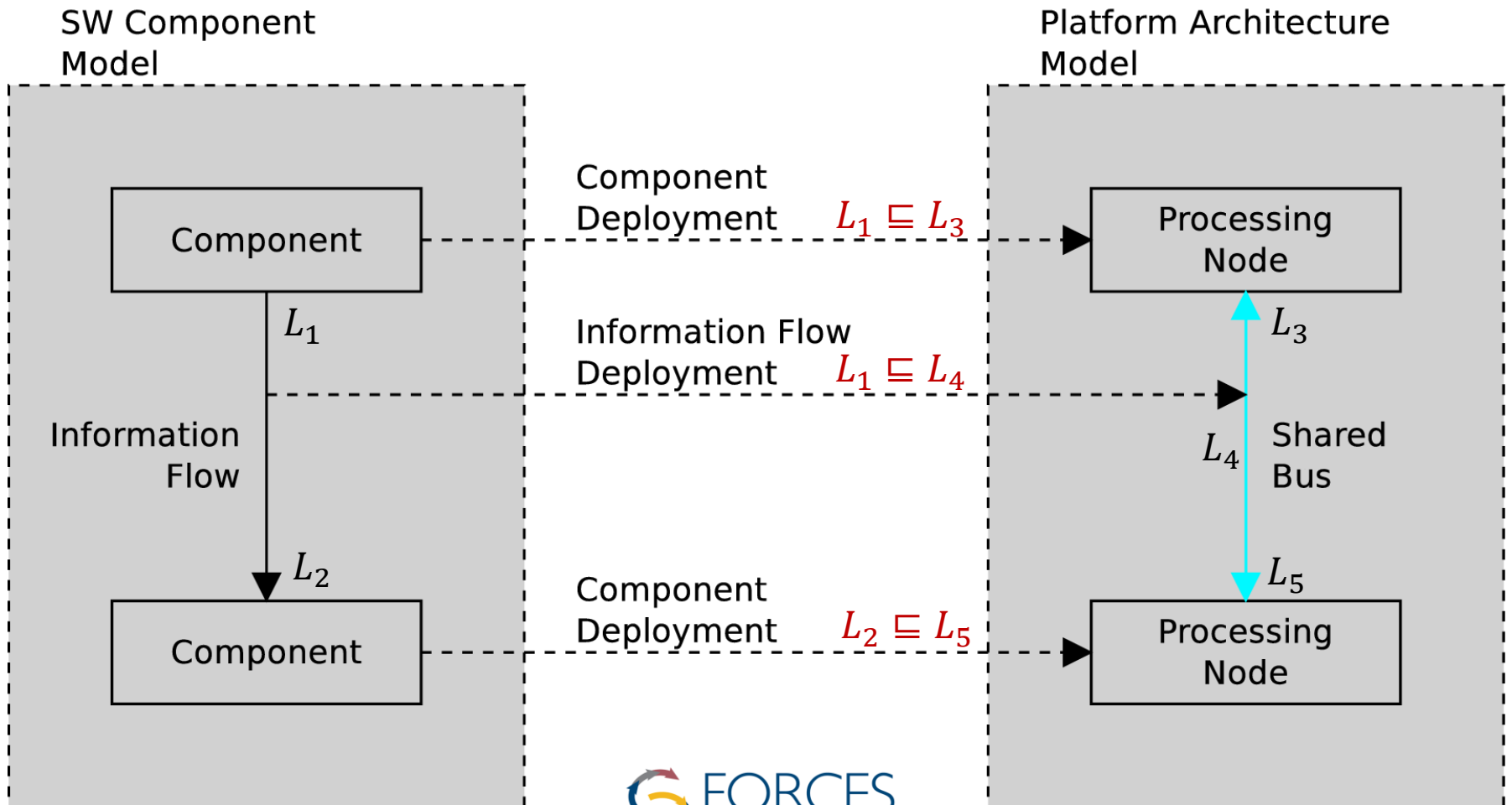
Information Flow Over SW Component Model



Information Flow Over SW Component Model



Information Flow Over Hardware Buses



Workflow for Designing Secure Distributed Embedded Systems

