



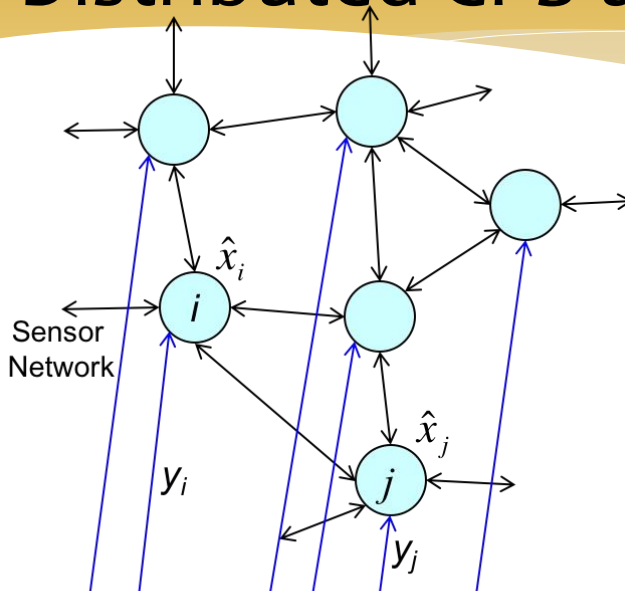
Resilience in Networked Dynamic Systems Using Trusted Nodes

Xenofon Koutsoukos

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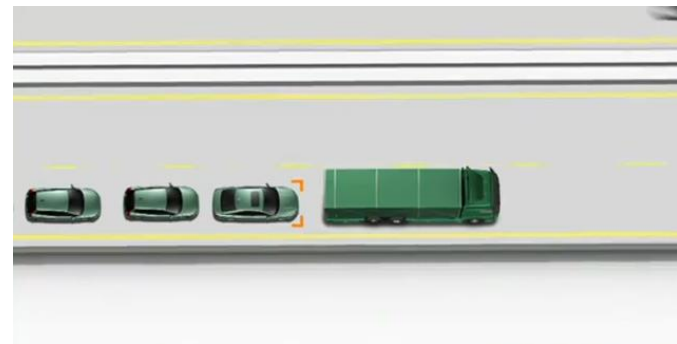
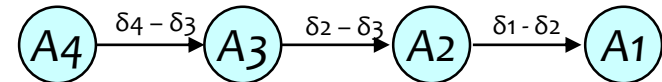
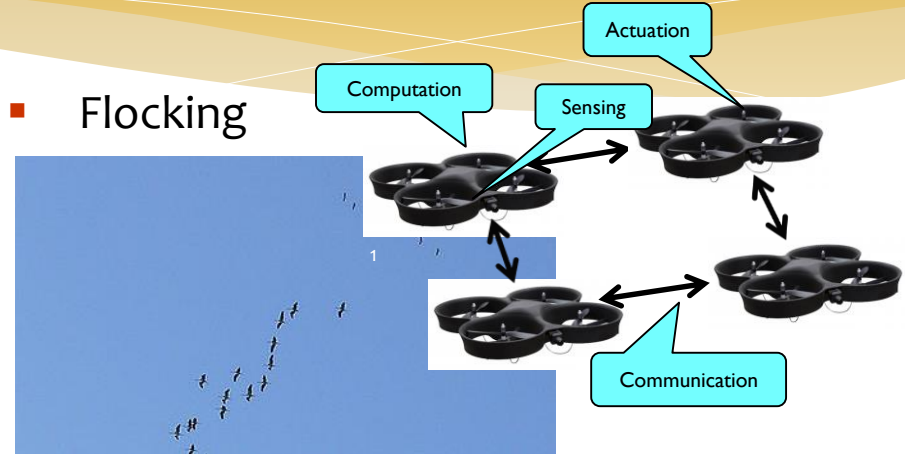
Motivation: Resilient Monitoring and Control of Distributed CPS using Consensus Algorithms



- Minimum variance estimate

$$\hat{q}_{MV} = \frac{\frac{1}{n} \hat{a}^n \frac{1}{S_i^2} y_i}{\frac{1}{n} \hat{a}^n \frac{1}{S_j^2} y_j}$$

- Flocking



- Formation control

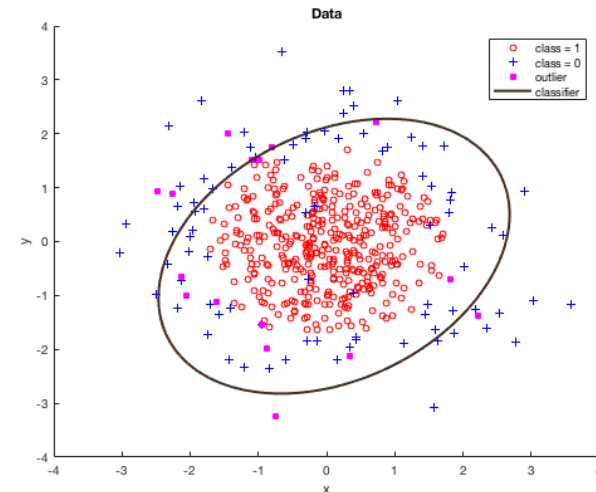
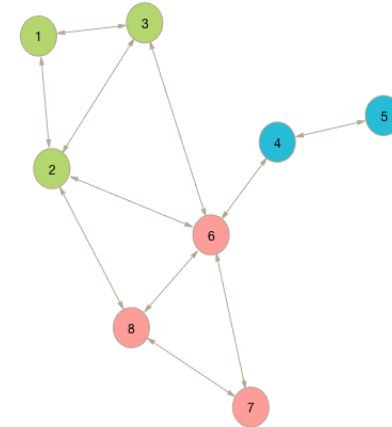
Learning by Networked Adaptive Agents

- Distributed collaborative classification
 - Each agent has access to some data and can share information with other agents
 - The objective is to classify the data using an elliptical curve
 - Each agent employs a logistic cost function

$$J_k(w) = \rho \|w\|^2 + \mathbb{E} \left\{ \ln[1 + e^{-\gamma_k h_k^T w}] \right\}$$

- The agents solve the optimization problem using a distributed collaborative consensus-algorithm

$$\begin{aligned} \psi_{k,i} &= w_{k,i-1} - \mu_k \widehat{\nabla_{w^T} J_k}(w_{k,i-1}) \\ w_{k,i} &= \sum_{\ell \in \mathcal{N}_k} a_{\ell k} \psi_{\ell,i} \end{aligned}$$

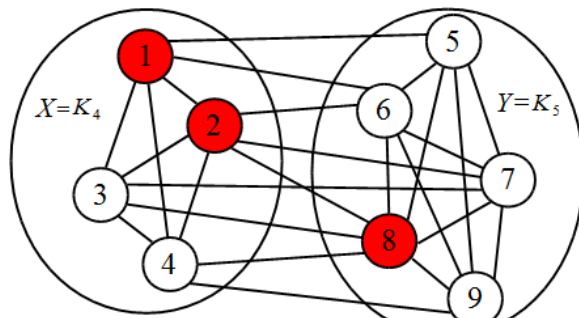


Outline

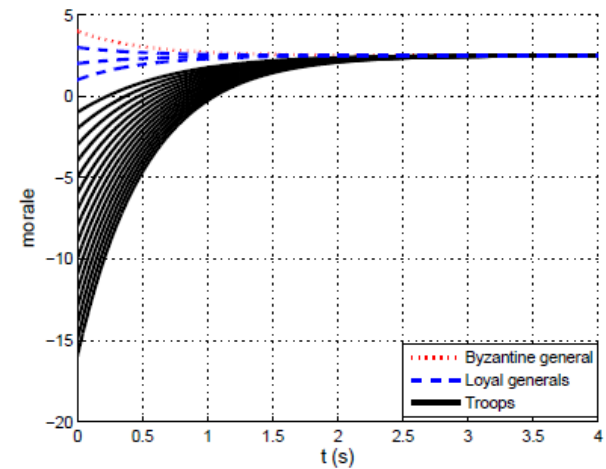
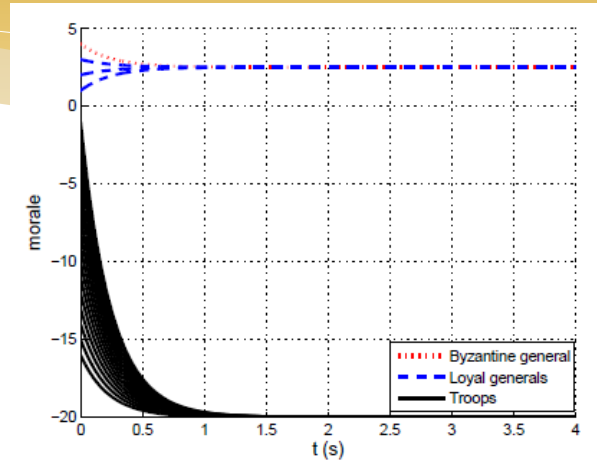
- Resilient distributed consensus in the presence of adversaries
- Resilient distributed consensus with trusted nodes
- Improving network connectivity by adding trusted nodes
- Conclusions and future directions

Resilient Distributed Consensus in the Presence of Adversaries

- * **Crash Adversary**
- * **Malicious Adversary**
 - * Must convey the same information to all neighbors
 - * Local broadcast model
- * **Byzantine Adversary**
 - * Can convey different information to different neighbors
- * **F-Total Model**
 - * At most F adversaries in the entire network
- * **F-Local Model**
 - * At most F adversaries in the neighborhood of any normal node
- * **f-Fraction Local Model**
 - * At most a fraction f of adversaries in the neighborhood of any normal node



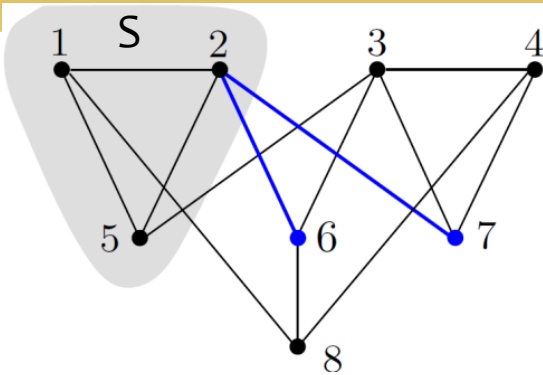
3-total, 3=local, (3/5)-fraction local



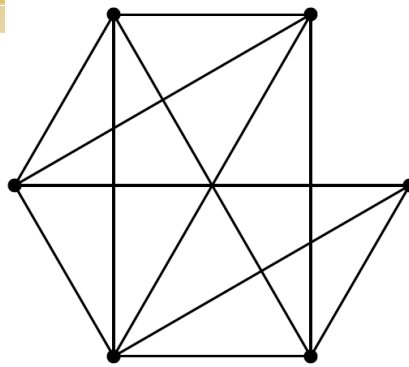
$$x_i(t+1) = w_{(i,i)}(t)x_i(t) + \sum_{j \in \mathcal{N}_i^{\text{in}}(t) \setminus \mathcal{R}_i(t)} w_{(j,i)}(t)x_{(j,i)}(t)$$

(b) ARC-P.

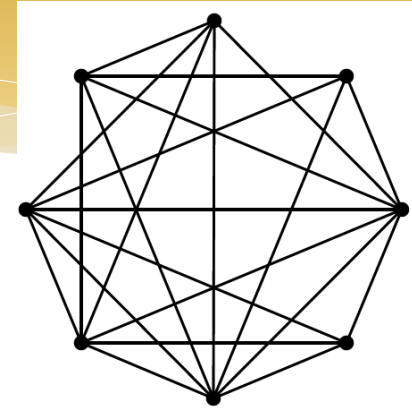
Robust Network Topologies



2-robust graph: Node 2 has two neighbors outside of S



(2,2)-robust

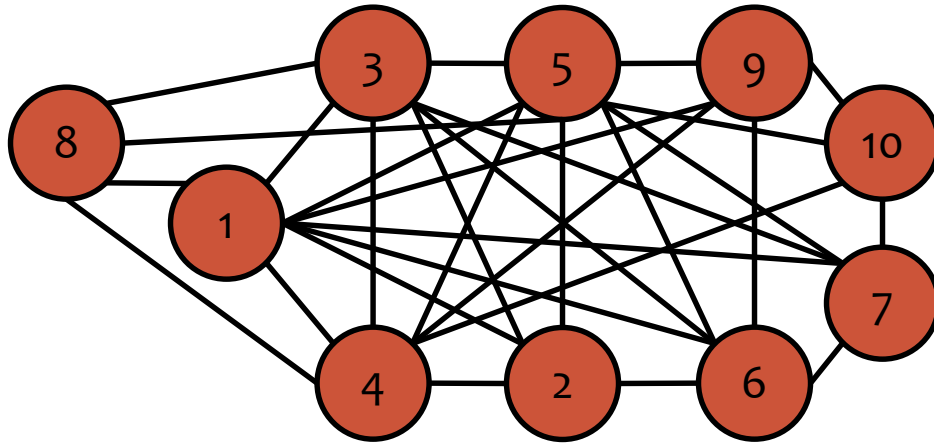


(3,3)-robust

- **Graph robustness:** New graph theoretic property to capture local redundancy
- Characterize a minimum number of nodes that are sufficiently influenced from outside their set
- **r -robustness:** For every pair of nonempty disjoint sets, at least one set has a node that has at least r neighbors outside the set
- **(r,s) -robustness:** For every pair of nonempty disjoint sets, there are at least s nodes with at least r neighbors outside their respective sets

Threat	Scope	Necessary	Sufficient
Crash & Malicious	F -Total	$(F+1, F+1)$ -robust	$(F+1, F+1)$ -robust
Crash & Malicious	F -Local	$(F+1, F+1)$ -robust	$(2F+1)$ -robust
Crash & Malicious	f -Fraction local	f -fraction robust	p -fraction robust, where $2f < p \leq 1$
Byzantine	F -Total & F -Local	Normal Network is $(F+1)$ -robust	Normal Network is $(F+1)$ -robust
Byzantine	f -Fraction local	Normal Network is f -robust	Normal Network is p -robust where $p > f$

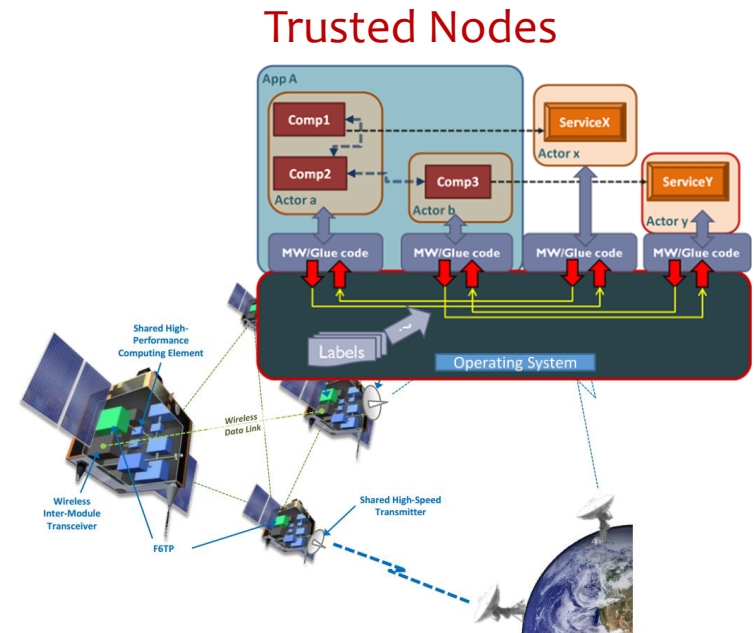
Construction of Robust Networks



Preferential-attachment model

- * Initial graph: K_5
- * K_5 is (3,2)-robust
- * Num edges / round: 4
- * End up with (3,2)-robust graph
- * Achieve resilient consensus in the presence of 1 adversary

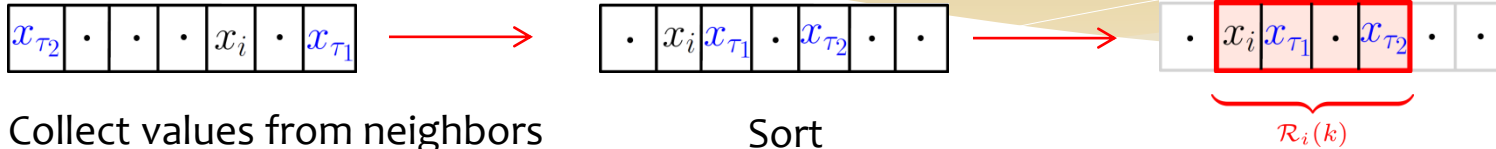
- * Resilience requires high degree of redundancy (high connectivity)
- * Redundancy increases the attack surface
- * How can improve resilience without adding redundancy?



Karsai et al., DARPA F6 program

Resilient Consensus Protocol with Trusted Nodes (RCP-T)

- If node i is connected to at least one trusted node



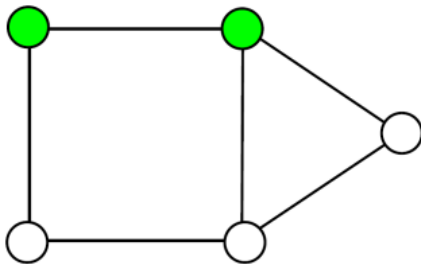
Collect values from neighbors

Sort

(x_{τ_1}, x_{τ_2} are trusted nodes' values)

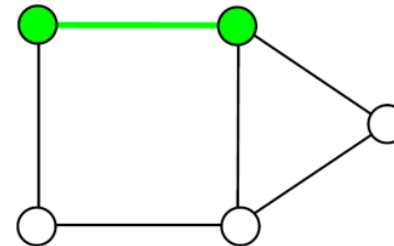
- Under RCP-T, consensus is always achieved in the presence of *arbitrary number of adversaries* if and only if there exists a set of trusted nodes that form a **connected dominating set**

Dominating Set

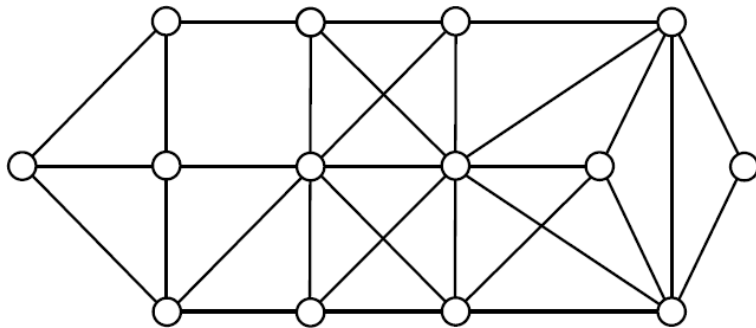


Connected Dominating Set

Nodes in the dominating set induce a connected subgraph



Trusted Nodes and Network Robustness



(2,2)-robust \longleftrightarrow Resilient against a single attack (with no trusted nodes)

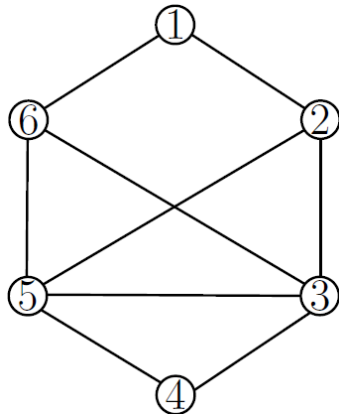
$d = 4$ \longleftrightarrow Resilient against any no. of attacks (with 4 trusted nodes)

Can we improve resilience if the number of trusted nodes $< d$?

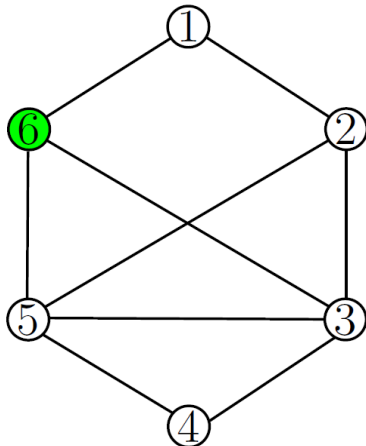
- * The **connected domination number** d is the number of vertices in the minimum connected dominating set
- * If the number of trusted nodes is at least d , the network can be made resilient against any number of adversaries

(r,s) -Robustness with Trusted Nodes

(r,s) -robustness with trusted nodes: For every pair of nonempty disjoint sets, there are at least s nodes with at least r neighbors or have trusted neighbors outside their respective sets

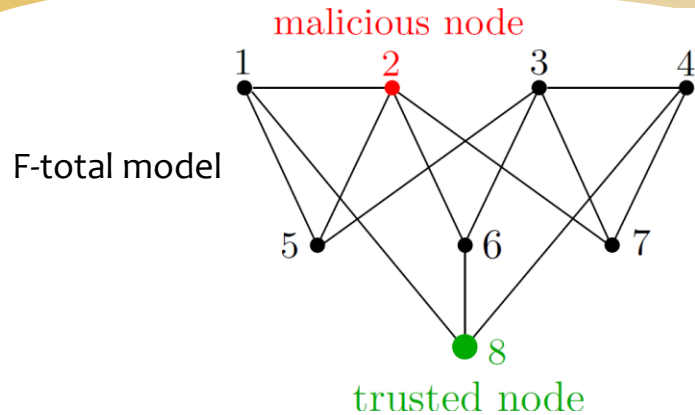


- The graph is **not (2,2)-robust**.
- For instance, consider $S_1 = \{1,2\}$; $S_2 = \{3,4,5,6\}$
- Node 1 has only neighbor outside S_1

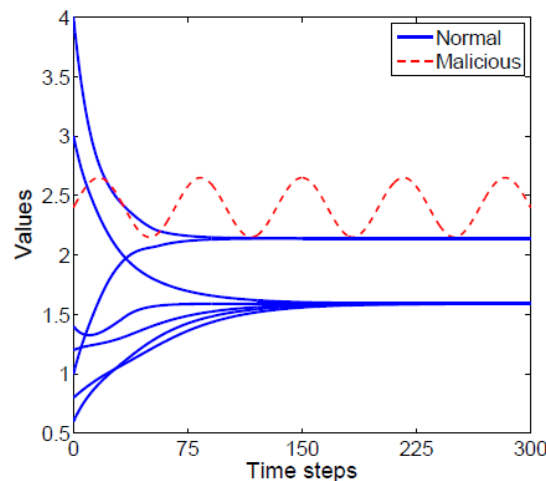


- The graph is **(2,2)-robust with 6 as a trusted node**.
- For instance, consider $S_1 = \{1,2\}$; $S_2 = \{3,4,5,6\}$
- Node 1 has a trusted neighbor outside S_1

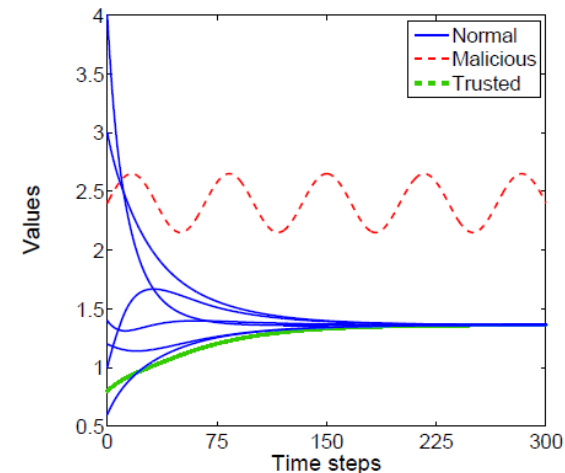
Resilient Consensus with Trusted Nodes: Example



- G is $(2,2)$ -robust with $T = \{8\}$.



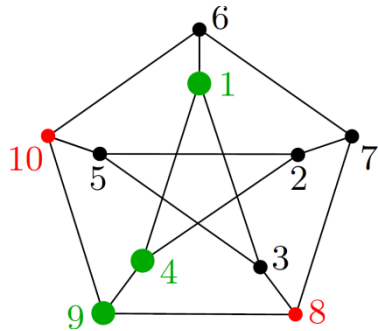
WMSR - algorithm
If there is no trusted node,
consensus cannot be achieved.



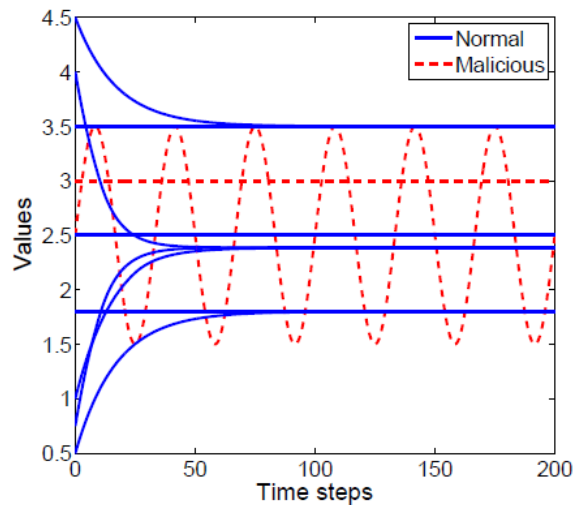
RCP-T - algorithm
Consensus is achieved with one
trusted node.

Resilient Consensus with Trusted Nodes: Example

F-local model

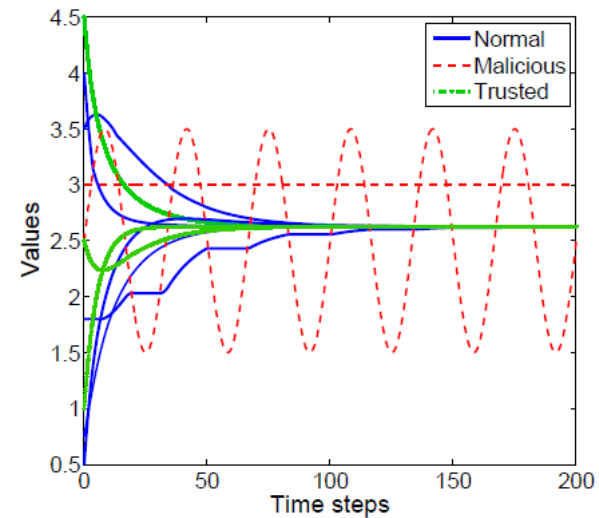


- G is 3-robust with $T = \{1, 4, 9\}$.



WMSR - algorithm

If there is no trusted node, consensus cannot be achieved.



RCP-T - algorithm

Consensus is achieved with three trusted nodes.

Resilient Consensus with Trusted Nodes

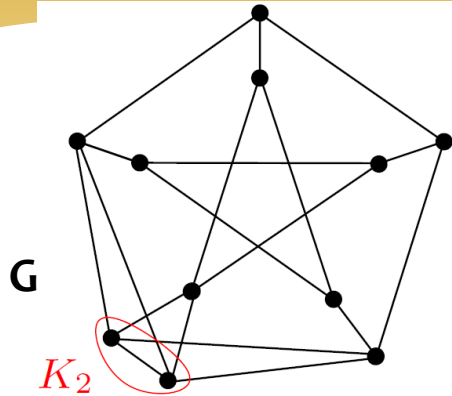
- The conditions for resilient consensus based on (r,s) -robustness can be reformulated using the notion of **(r,s) -robustness with trusted nodes**
- For instance

Theorem :

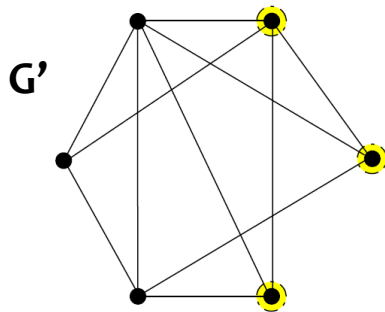
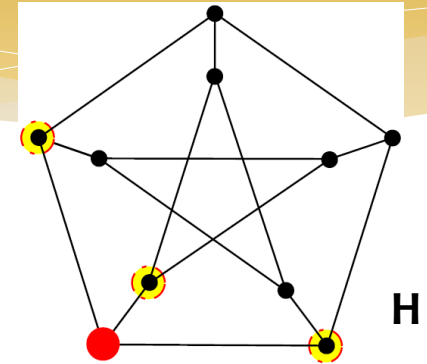
Under the F -total malicious model, resilient asymptotic consensus is achieved using RCP-T algorithm if and only if the network topology is **$(F+1,F+1)$ -robust with trusted nodes.**

- A graph that is $(F+1,F+1)$ -robust with trusted nodes could be much sparser than the one that is $(F+1,F+1)$ -robust without trusted nodes

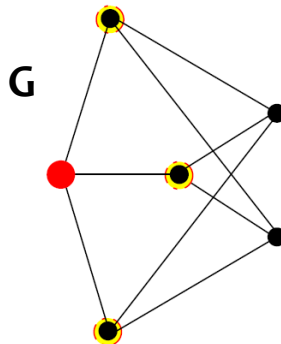
Construction of Robust Graphs with Trusted Nodes



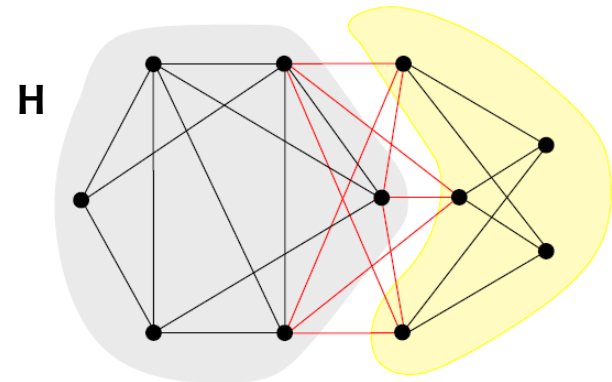
- G is 2-robust
- Clique K_2 is replaced by a trusted node
- H is still 2-robust



- **G' is 3-robust.**
- Nodes in subset η are highlighted

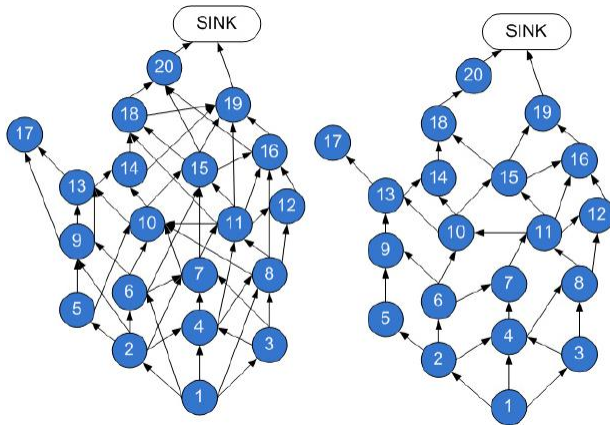


- **G is 3-robust with red trusted node.**
- Neighbors of trusted node are highlighted



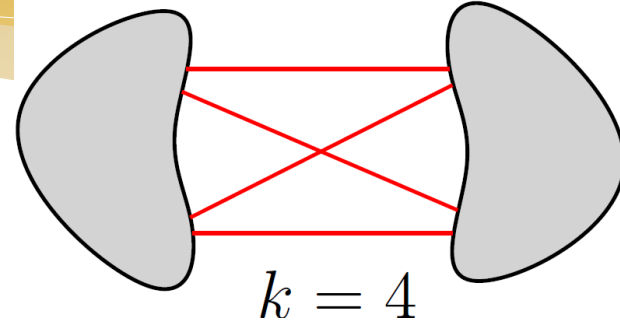
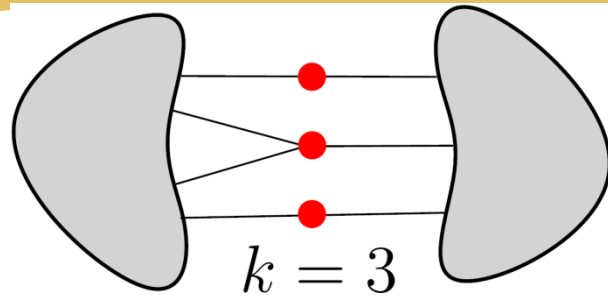
- **H is also 3-robust.**
- New edges added are shown in red.

Improving Network Connectivity Using Trusted Nodes and Edges



- Connectivity is primary attribute of every network
- Many important network properties depend on vertex (edge) connectivity
- How can we efficiently place trusted nodes in a network to increase vertex (edge) connectivity?

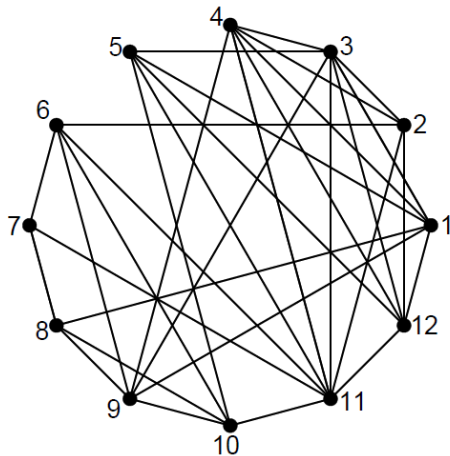
Improving Network Connectivity



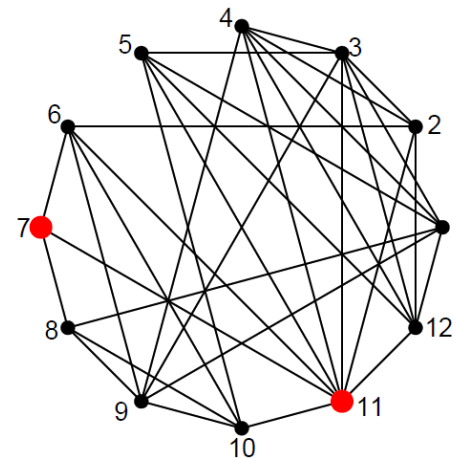
- * **k-vertex Connectivity:** A graph remains connected even if any set of $(k-1)$ vertices are removed
- * **k-edge Connectivity:** A graph remains connected even if any set of $(k-1)$ edges are removed
- * **Connectivity augmentation:** Determine the *smallest set of edges* which must be added to a given graph to make it k -edge connected or k -vertex connected
- * Connectivity augmentation may be difficult due to practical and economical reasons and increases the attack surface
- * **Improving network connectivity using trusted nodes:** Deploy a small subset of trusted nodes

Network Connectivity with Trusted Nodes

Network connectivity can be measured by the number of *non-trusted nodes* that need to be removed to make the graph disconnected.



- The graph is 3-vertex connected.
- At least 3 nodes need to be removed to disconnect the graph.



- By making nodes *7* and *11* as *trusted*, we need to remove at least 6 of the remaining nodes to disconnect the graph.
- In other words, with nodes 7 and 11 as trusted, the graph behaves like a 6-vertex connected.

Menger's Theorem: Independent Paths

The minimum number of nodes whose removal disconnects two nodes, say u and v , is equal to the maximum number of pairwise node-independent paths from u to v .

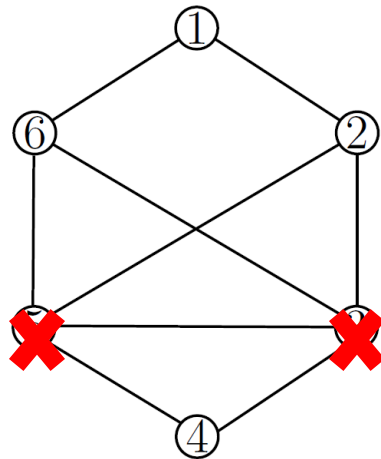
In other words,

Node
connectivity

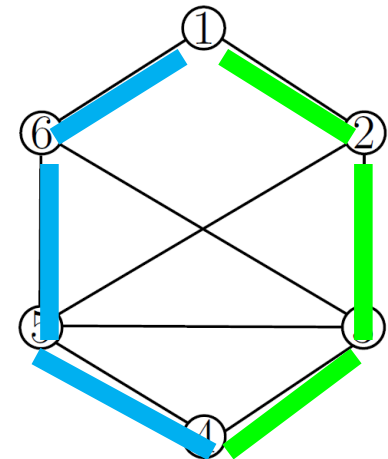


No. of node independent paths
between any two nodes

- Removal of two nodes can disconnect the graph.
- In particular, nodes 1 and 4 can be disconnected by 2 node removals.



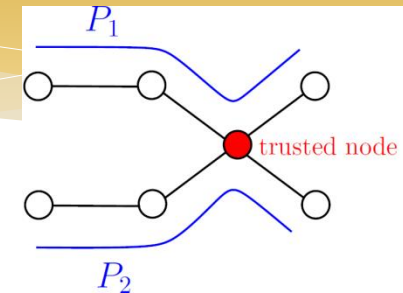
- There are two node-independent paths between any two nodes.
- The ones between nodes 1 and 4 are shown.



Independent Paths with Trusted Nodes

Definition (Node-independent paths with trusted nodes):

Two paths are node-independent with trusted nodes if common nodes between them are only the trusted nodes.



P_1 and P_2 are node-independent paths with trusted nodes.

Definition (Node trusted path):

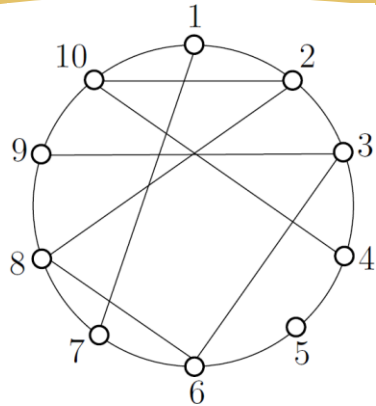
A path with all trusted nodes is a node trusted path.

Theorem:

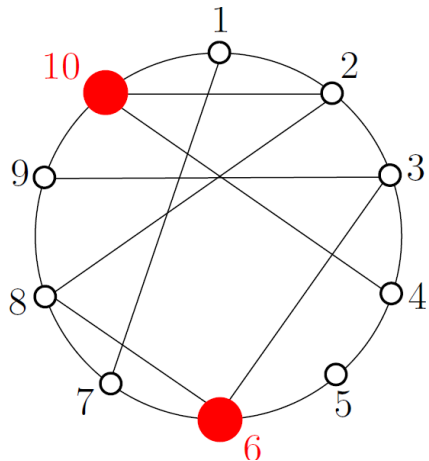
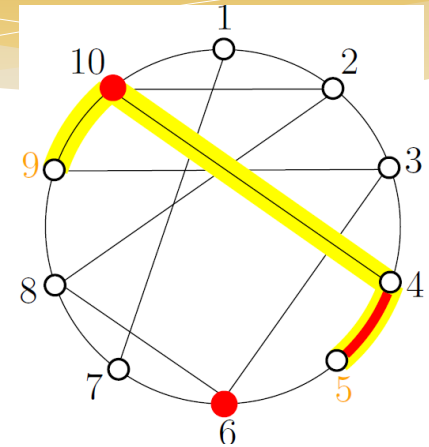
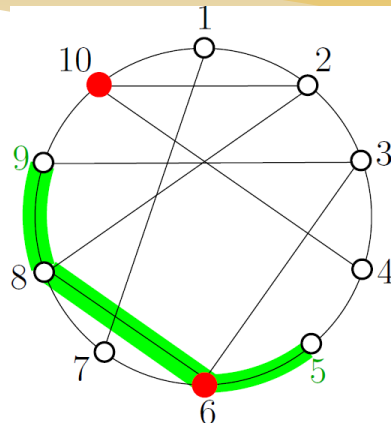
For a graph $G(V,E)$ and a set of trusted nodes T_v , following statements are equivalent.

1. G is k -vertex connected with T_v .
2. For any two distinct, non-adjacent vertices u and v , either there exists a node-trusted path between u and v , or there exists at least k paths between u and v that are vertex-independent with T_v .

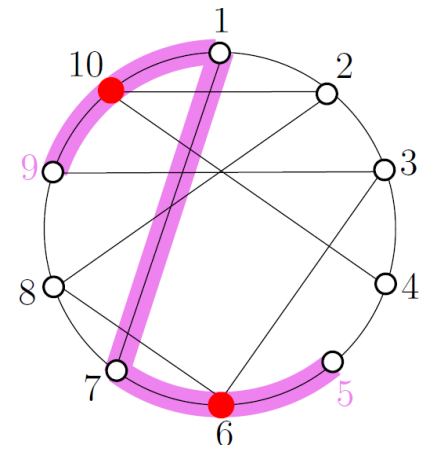
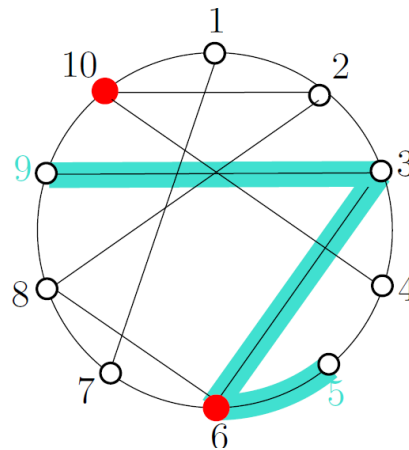
Vertex-Connectivity with Trusted Nodes: Example



A graph is 2-vertex connected.



A graph is 4-vertex connected with the **red** nodes as the **trusted nodes**, i.e., between any two nodes there always exist four node-independent paths with trusted nodes.



Four node-independent paths with (red) trusted nodes between nodes 5 and 9

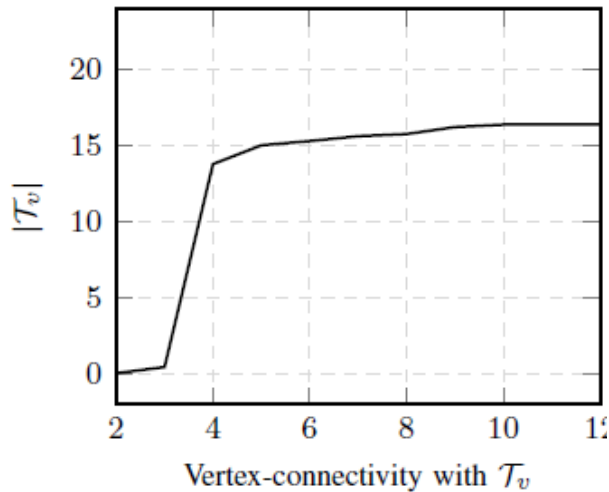
Placement of Trusted Nodes

Theorem:

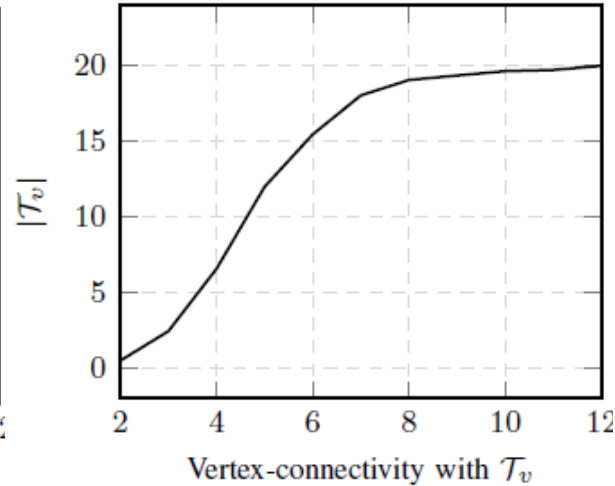
Given a graph $G(V, E)$, a desired vertex connectivity k , and the number of trusted nodes T , determining if there exists a set of trusted nodes T_v of cardinality T such that G is k -vertex connected with T_v is NP-hard.

1. Heuristics based on Connected Dominating Set (CDS)
 - * If trusted nodes form a CDS, then, between any two nodes, there is always a path consisting of only trusted nodes.
 - * Start with a set of trusted nodes forming a CDS, and then successively reduce the set of trusted nodes as long as the desired connectivity is obtained.

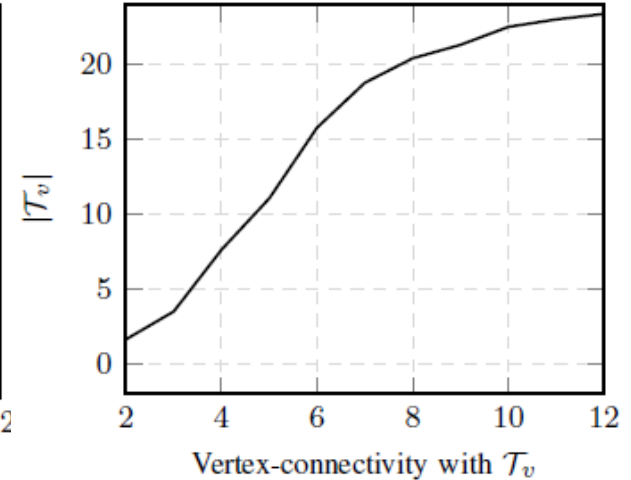
Number of Trusted Nodes T_v as a Function of the Vertex Connectivity



Preferential attachment networks



Erdos-Renyi random networks



Random geometric networks

(Each network has a total of 100 nodes. Details of the networks are in Abbas et al., 2017)

Conclusions and Future Directions

- * Resilient Consensus Protocols in the Presence of Adversaries
 - * Exploit local redundancy to ensure asymptotic consensus
 - * Characterize robust network topologies
- * Resilient Consensus Protocols with Trusted Nodes
 - * Increase resilience by exploiting trust instead of redundancy
- * Improving Network Connectivity with Trusted Nodes
 - * Improve reliability, resilience, and other properties based on connectivity
- * Can trusted nodes be used to improve resilience of other properties in networked CPS?
 - * Participatory sensing
 - * Learning by networked agents