Cyber-Physical Systems Research Challenges



George J. Pappas
Joseph Moore Professor
NSF CPS Academic Executive Board
University of Pennsylvania
pappasg@seas.upenn.edu



NSF CPS AEB Membership



George Pappas, U Penn	control
Janos Sztipanovits, Vanderbilt	embedded software
Edward Lee, UC Berkeley	education
Eric Feron, Georgia Tech	avionics
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Claire Tomlin, UC Berkeley	hybrid systems
Raj Rajkumar, CMU	real-time systems
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Cyber-physical systems*











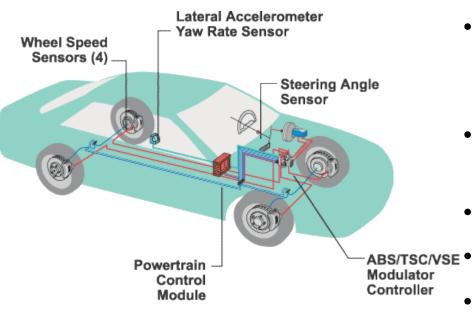




*Information systems interacting with physical systems, broadly

Automotive Systems

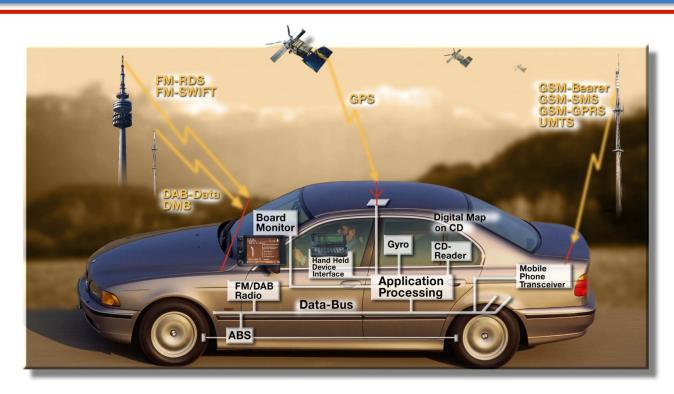




- Longitudinal dynamics : ABS (antilock brake system) and ASC (automatic stability control)
- Lateral dynamics : EDRC (engine drag reduction control) and CBC (corner braking control)
- DSC (dynamic stability control) is using all the above
- Also:
- Automatic gearboxes
- Anti-theft systems
- Multimedia systems
- Navigation systems
- etc

Automotive Systems





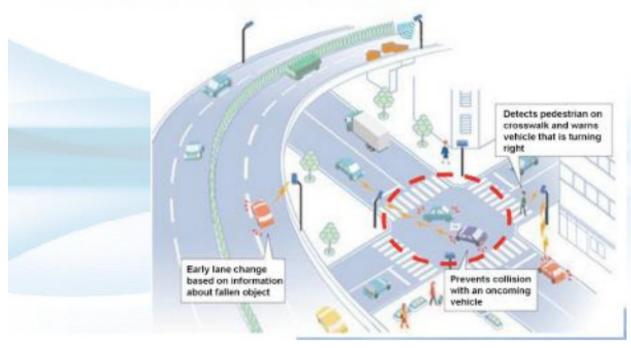
- 270 user interactive functions
- 67 embedded platforms / (5 data buses?)
- 65 MB of binary code.
- Next generation (2010):
 - \sim 1 GB of software, IP is being studied

Smart Road Infrastructure



Methods Gap - Cyber-physical systems (e.g. collision avoidance)

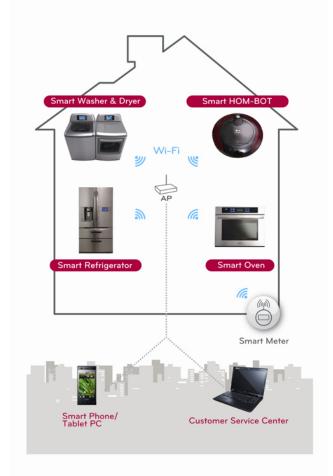
- · Heterogeneous modeling hybrid dynamics, wireless networking, dynamic agent scenarios
- · Abstractions and refinements for synthesis and analysis hierarchical systems structures
- · Component composability and consistency systems integration, rapid development
- · Verification (within and across) design and implementation



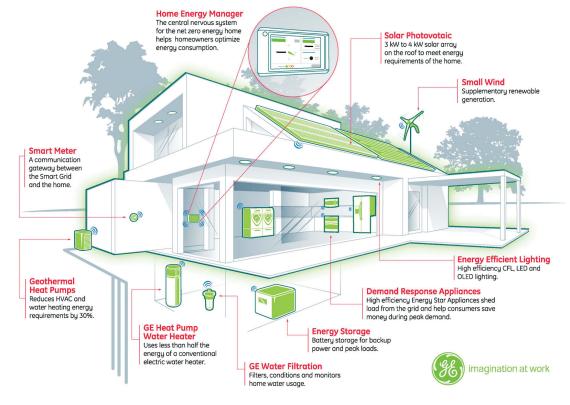
Embedded in : Home appliances







GE Targets Net Zero Energy Homes by 2015



Medical devices



Vision: Doctor-on-a-chip, artificial eye

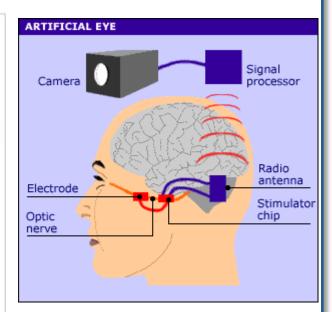
Operating room of the future & Digital hospital

Remote monitoring of elderly (body area networks)

Medical implants (artificial eyes, ears etc)







BBC news

NSF CPS AEB: What are we doing well?



CPS is impacting numerous critical sectors of society

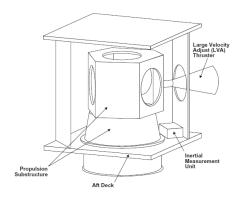
□ Energy (SmartGrid, Microgrids, Green Buildings)
 □ Health (Medical devices)
 □ Automotive/Transportation (Mobile Millenium)
 □ Aerospace/Air traffic control (NextGen)
 □ Infrastructure monitoring (bridges, lakes)
 □ DoD (META)
 □ Robotics (medical robotics, mobile sensor networks)
 □ Next generation flexible manufacturing?

Tremendous potential for broad impact

CPS problems due to software-system interactions



1998: Near Earth Asteroid Rendezvous (NEAR)



2005: Demonstration of Autonomous Rendezvous Technology (DART)



1996: Ariane 5



1991: Patriot Missile Software Failure

Caltech's 2005 DARPA Grand Challenge Entry, Alice





1997: Korean Air 747 in Guam



Northeast Blackout







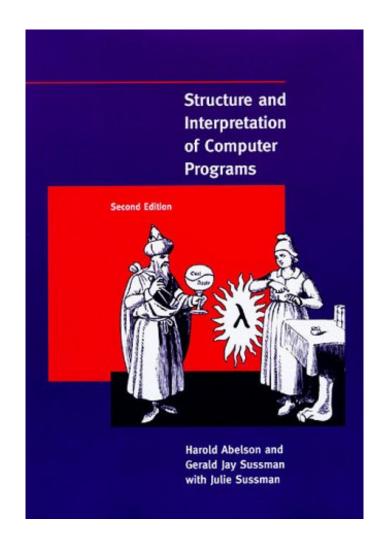
Research Challenges

Part I: Modeling

Part II: Verification

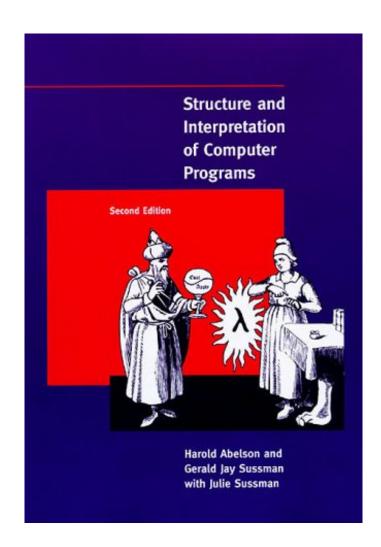
Part III: Robustness

Part IV: Security



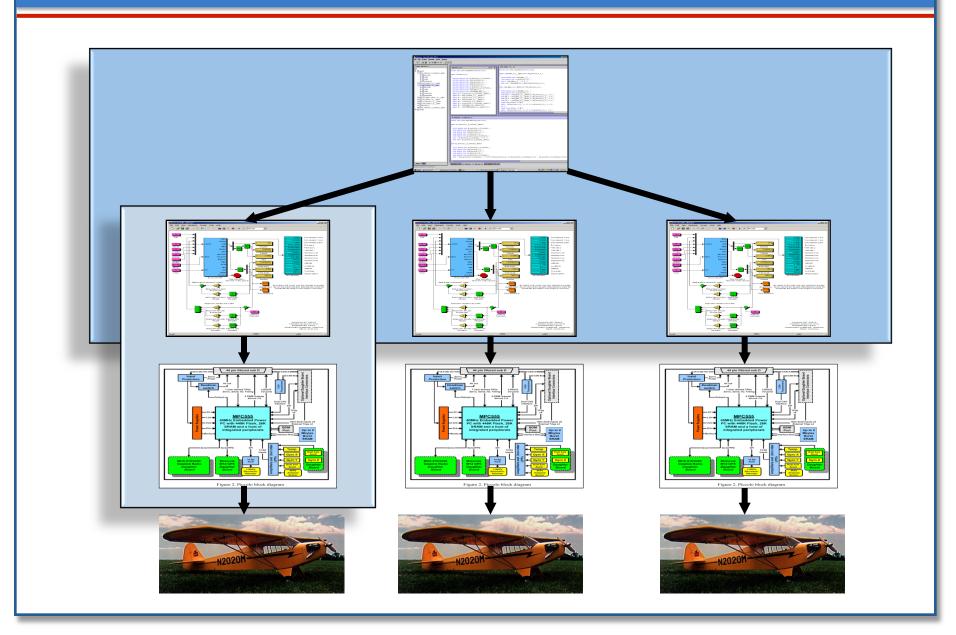


Part I: Modeling challenges



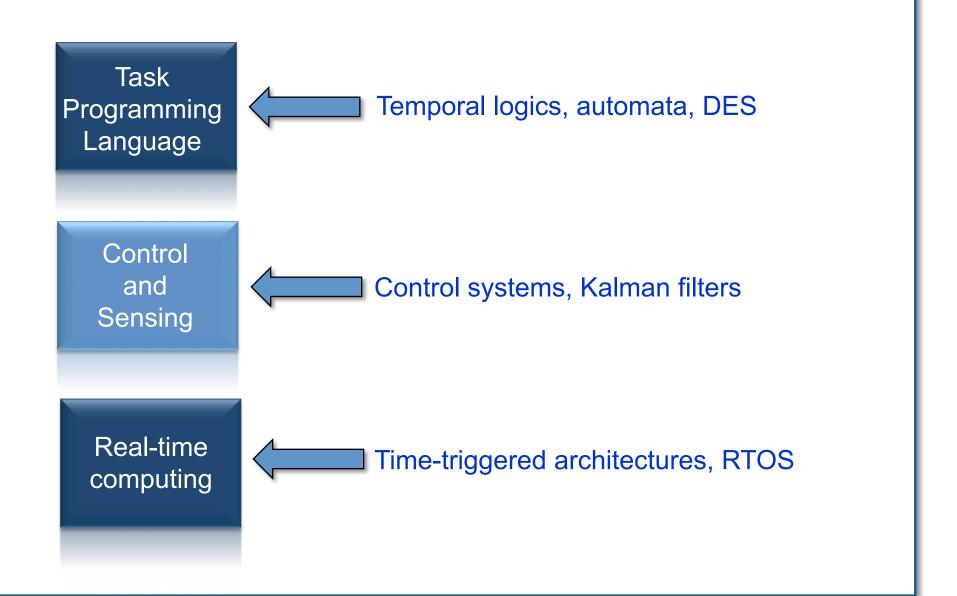
Control and computing hierarchies





Control and computing hierarchy

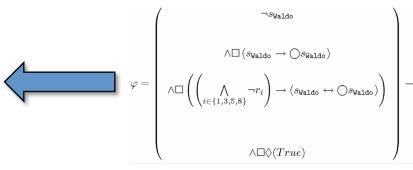


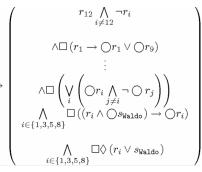


A CS view of the same hierarchy



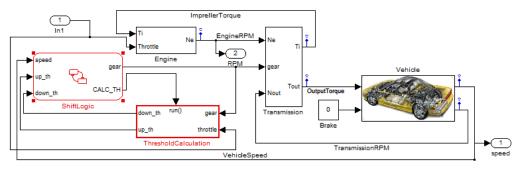






Control and Sensing





Real-time computing



```
float updateF1(pi_block* ppi, float vi;

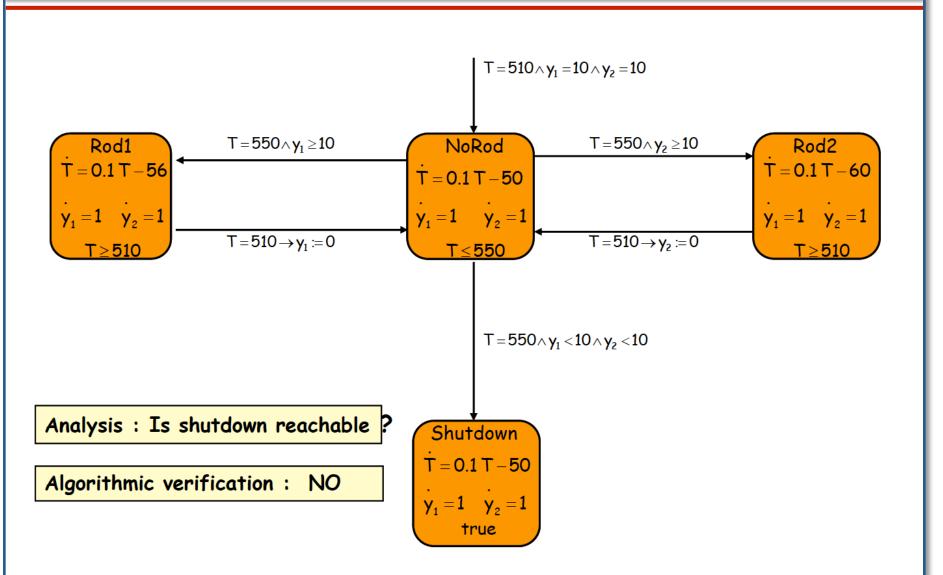
/* Integrator Block i */
void integrator() {
    double ini, in2;
    double ini, in2;
    double delict;

    curfime = getTime();
    ini = Imput();
    ini = Imput();
    delict input();
    delict input();

    delict = curfime;
    delict = curfime = prevTime;
    prevTime = curfime;
    x1 = delict();
    x2 = delict();
    y2 = delict();
    y3 = delict();
    y4 = delict();
    y4 = delict();
    y5 = delict();
    y5 = delict();
    y6 = delict();
    y7 = delict();
    y8 = delict();
    y9 = delict();
    y
```

Hybrid model of nuclear reactor

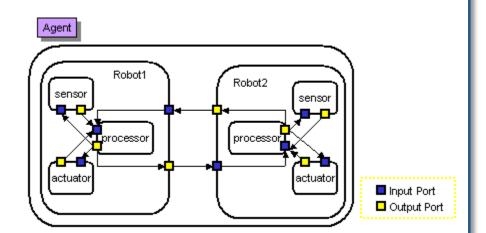


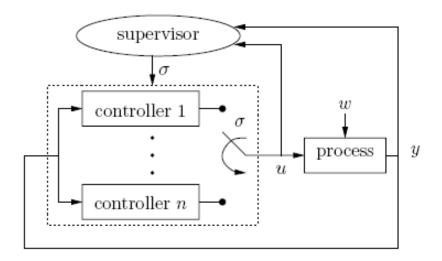


A zoo of hybrid systems



- Hybrid Automata
- Hybrid Input-Output Automata
- Hybrid Petri Nets
- Simulink/Stateflow MATLAB models
- Supervisory control systems
- Switched systems
- Nonsmooth systems
- Piece-wise affine systems (PWA)
- Mixed Logical Dynamical
- Linear complementarity models





Major switching types



Time-triggered: Switching depends on time only

Switching and dynamics are decoupled

Switching times are known a priori

Switched systems more appropriate

Event-triggered: Switching also depends on state

Switching and dynamics are coupled

Switching times not known a priori

Hybrid automata more appropriate

Similarly there is a large variety of concurrency/synchronization types

CPS AEB: Modeling challenges for CPS

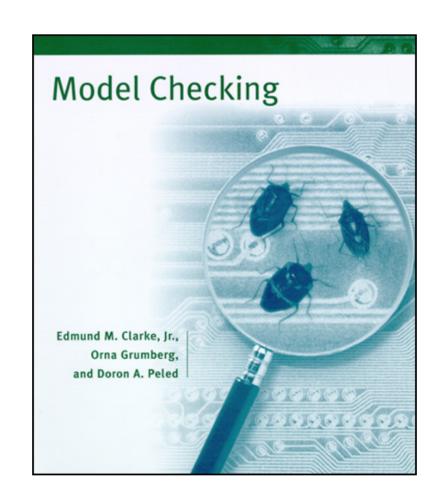


A richer, systems view of computer science is needed. Ingredients include:

- Enriching CS models with relevant physical/resource properties
 - Physical, model-based computing
 - Resource aware (time/energy) computing
- ☐ Formal composition of **multiple** physics, models of computing, languages
 - Composition of heterogeneous components
- ☐ Impact of cyber components on physical components and vice versa
 - Physically-aware computing



Part II: Composable verification for CPS



Traditional verification view



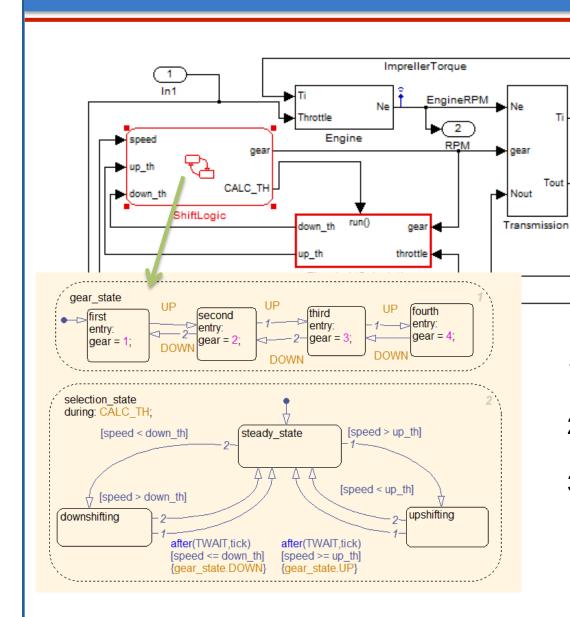
Safety verification: Is Reach(S) \cap S_F empty?

Model checking: Does S satisfy temporal logic formula ϕ ?

Controller Synthesis: Does S||C satisfy temporal logic formula?

Automatic trasmission verification





1. the vehicle speed v exceeds 120km/h

Vehicle

OutputTorque

Brake

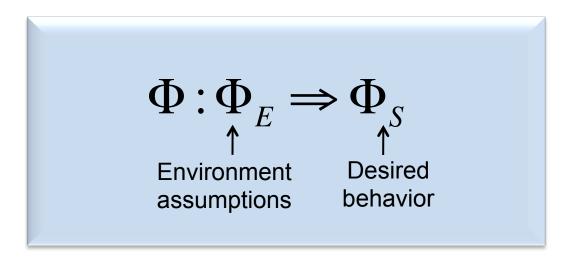
TransmissionRPM

- 2. the engine speed ω exceeds 4500RPM
- 3. all states are reached in the switching logic

Assume-guarantee reasoning



Consider component specifications of the form



If the component assumptions are met, then desired behavior should be guaranteed Assumptions can model other components, physical or computational

How to analyze the overall design?



Task Programming Language

Control and Sensing

Real-time computing

Technical Challenges

Composable interfaces between control and computing components

CPS AEB: Methods & Tools Challenges



□ From verification of C and P systems to verification of CPS systems
 □ Compositional verification and design

 Requires composable interfaces
 Assume-guarantee reasoning (contracts)

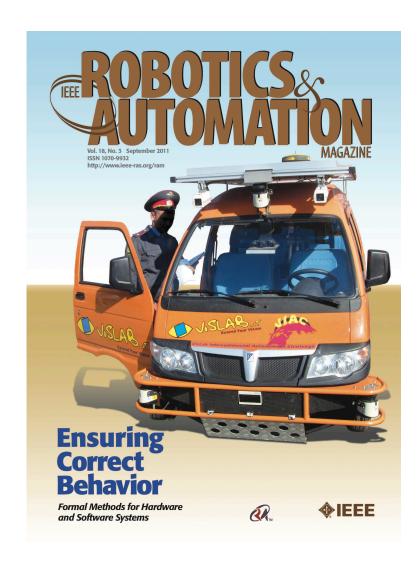
 □ Tradeoffs between different system views

 Control performance versus scheduling flexibility
 Understanding the price of compositionality

 □ From methods for flat CPS to methods for distributed CPS systems
 □ From verification methods to synthesis and code-generation

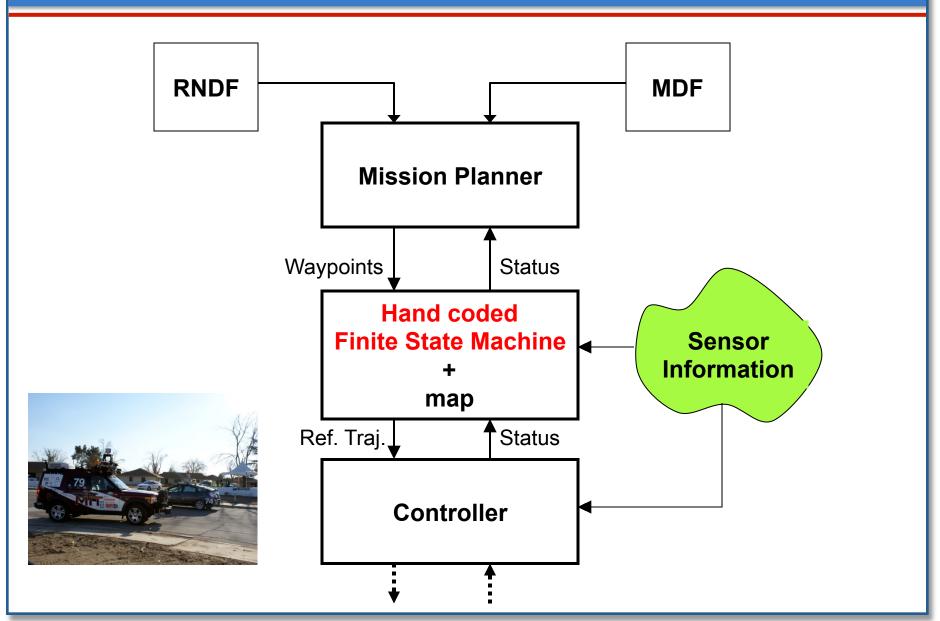


Part III: Robustness



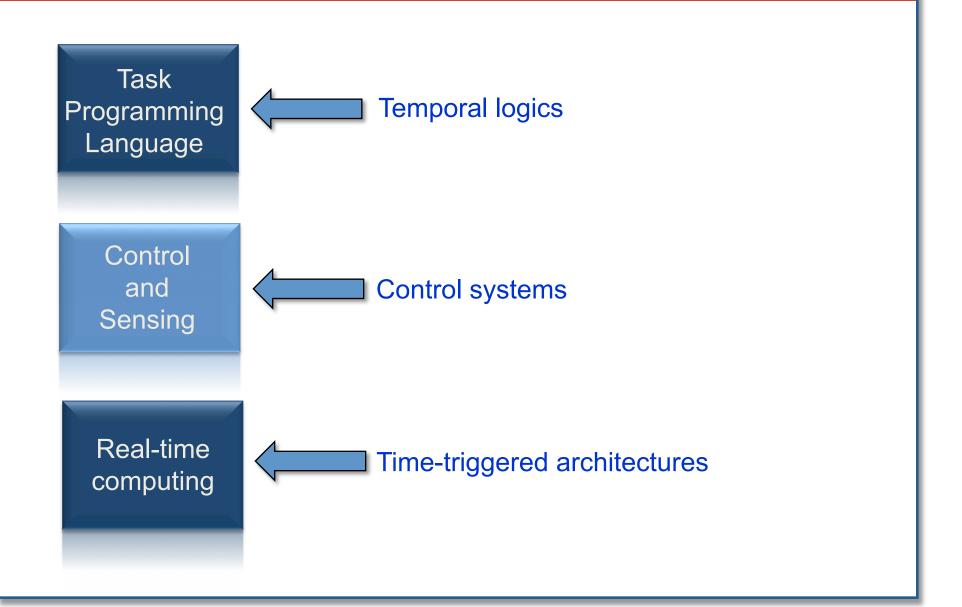
DARPA Urban Challenge





Control and computing hierarchy





From exact to approximate



- Exact relationships useful for binary answers/relations
- Exact results are fragile with respect to uncertainty
- When interacting with the physical world, we need approximations

Labeled Markov processes (Desharnais et. al., TCS 2004)

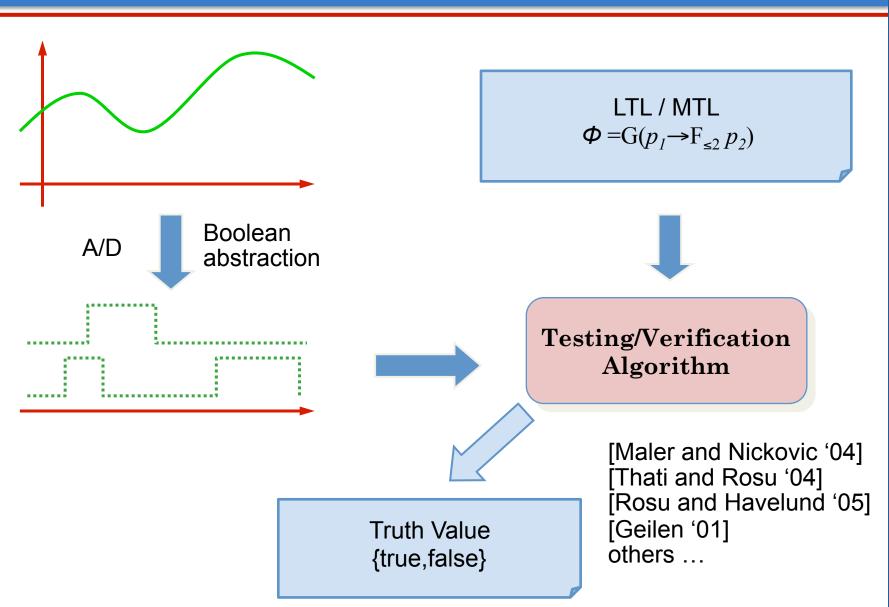
Quantitative transition systems (de Alfaro et. al., ICALP 2004)

Quantitative generalizations of languages (Henzinger, DLT 2007)

- Approximate system relationships
 - Enable larger system "compression"
 - Quantify error/complexity tradeoffs
 - Provide measures of robustness
 - Potentially introduce different algorithms

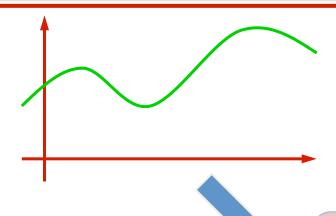
Boolean testing & verification of CPS





Robust verification for CPS



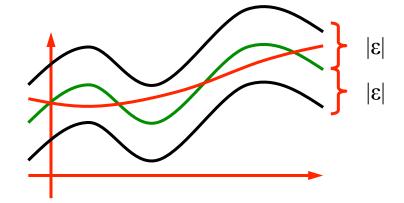


LTL / MTL $\boldsymbol{\Phi} = G(p_1 \rightarrow F_{\leq 2} p_2)$

Robust Testing/ Verification Algorithm



Robustness Estimate ϵ



CPS AEB: Robust methods for CPS



- ☐ Computing needs to transform from exact to approximate.
 - From boolean to robust computing
 - From qualitative to quantitatve
- ☐ Robust computing interacting with uncertain world (stochastic)
 - Uncertainty due to P or C components, networking
 - From deterministic to stochastic verification
- ☐ Robustness verification of CPS systems
 - Closer to sensitivity analysis
- □ Verification of numerical software



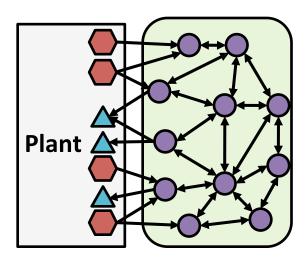
Part IV : CPS Security Beyond Cybersecurity



Monitoring for faulty and malicious behavior



- Many CS efforts focus on building security walls
- What happens after a fault or after an attack?
- What if certain wireless nodes become faulty or malicious?



Monitoring for faulty and malicious behavior



- Security of control networks in industrial control systems is a major issue [NIST Technical Report, 2008]
 - Data Historian: Maintain and analyze logs of plant and network behavior
 - Intrusion Detection System: Detect and identify any abnormal activities
- Is it possible to design an Intrusion Detection System (IDS) to determine if any nodes are not following protocol?
- Can IDS scheme avoid listening all nodes? Under what conditions? Which nodes?

CPS AEB: Secure CPS



- ☐ CPS Beyond cybersecurity
 - From pre-attack defense to post-attack methods
 - Detection, identification, reconfiguration, graceful degradation
- ☐ Theory of resilient and trustworthy CPS
 - Security/trust metrics
- ☐ Privacy for CPS
 - Crucial in many contexts (i.e. medical CPS)
 - Interesting ideas in database theory (differential privacy)