



Distributed Asynchronous Algorithms & Software Systems For Wide-Area Monitoring of Power Systems

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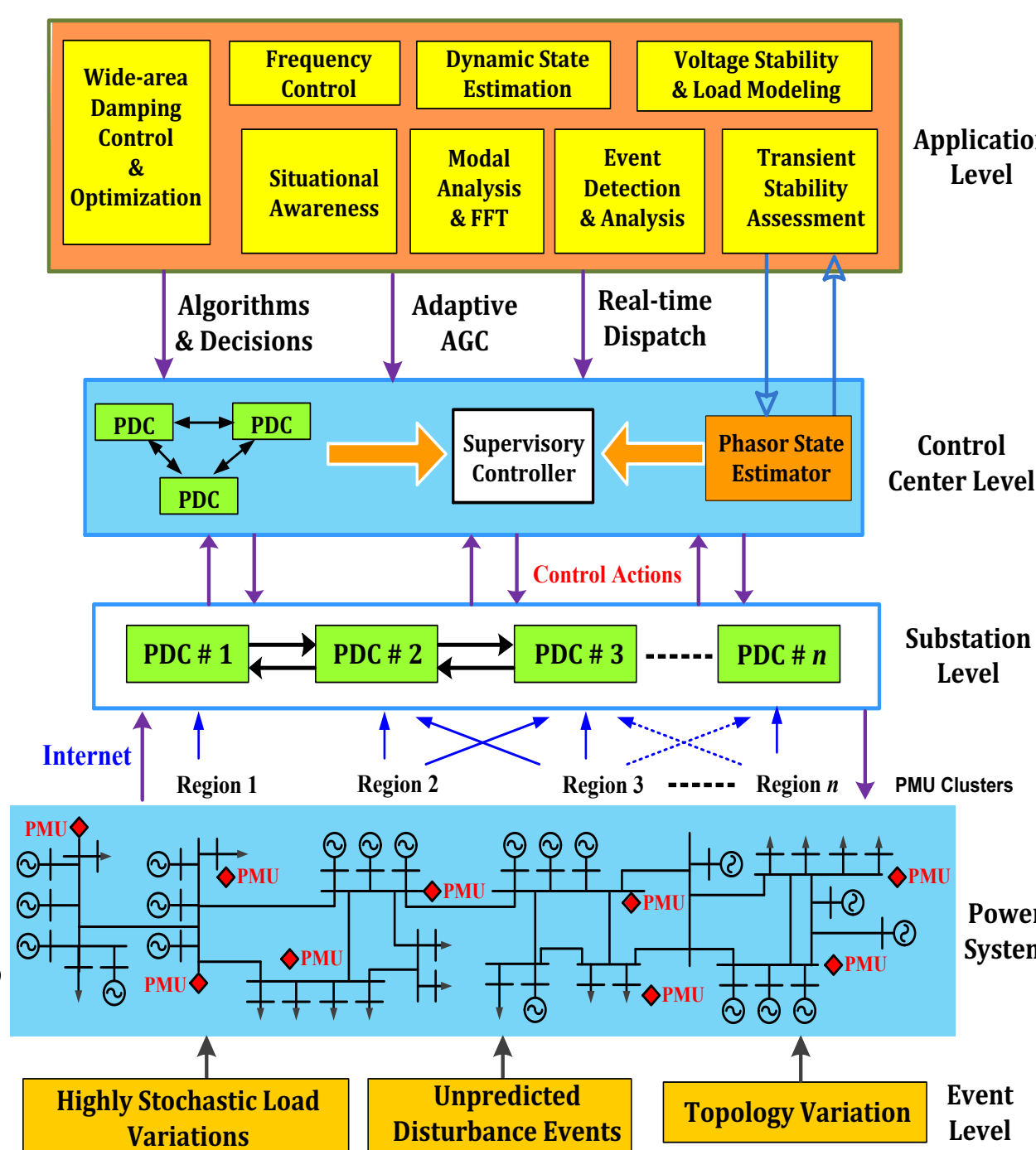
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Project Goal

To translate current state-of-art centralized processing algorithms for wide-area monitoring of large power grids using large volumes of Synchrophasor data to a completely distributed cyber-physical architecture.

Intellectual Merits:

1. Distributed oscillation monitoring
2. Distributed voltage monitoring
3. Distributed middleware
4. Fault-tolerance
5. Experimental verification using Exo-GENI network
6. Real-time testing of QoS and cyber-security



Distributed Oscillation Monitoring

Problem statement: Compute power flow oscillation frequencies (eigenvalues), mode shapes (eigenvectors), damping, and residue from PMU measurements *using distributed algorithms* implemented via DRCP and DLAP.

Power System Dynamic Model

$$\begin{bmatrix} \Delta \dot{\delta} \\ M \Delta \dot{\omega} \\ \Delta \dot{E} \end{bmatrix} = \begin{bmatrix} 0 & I & 0 \\ -L(G) & -D & -P \\ K & 0 & J \end{bmatrix} \begin{bmatrix} \Delta \delta \\ \Delta \omega \\ \Delta E \end{bmatrix} + \begin{bmatrix} 0 \\ \text{col}_{i=1}^{(1)n}(\gamma_i) \\ \text{col}_{i=1}^{(1)n}(\rho_i) \end{bmatrix} + \begin{bmatrix} 0 & 0 \\ 0 & I \\ I & 0 \end{bmatrix} \begin{bmatrix} \Delta P_m \\ \Delta E_F \end{bmatrix}$$

due to load

$$y = \text{col}_{i \in \mathcal{S}}(\Delta V_i, \Delta \theta_i).$$

$$\Delta \theta_i(t) = \frac{b_0 + b_1 z^{-1} + b_2 z^{-2} + L + b_{2n} z^{-2n}}{1 + a_1 z^{-1} + a_2 z^{-2} + L + a_{2n} z^{-2n}}$$

Distributed LS problem

$$\text{minimize}_{\mathbf{a}_i, \mathbf{K}_i, \mathbf{a}_N, \mathbf{z}} \sum_{i=1}^N \frac{1}{2} \|\hat{\mathbf{H}}_i \mathbf{a}_i - \hat{\mathbf{c}}_i\|_2^2$$

subject to $\mathbf{a}_i - \mathbf{z} = 0$, for $i = 1, \dots, N$

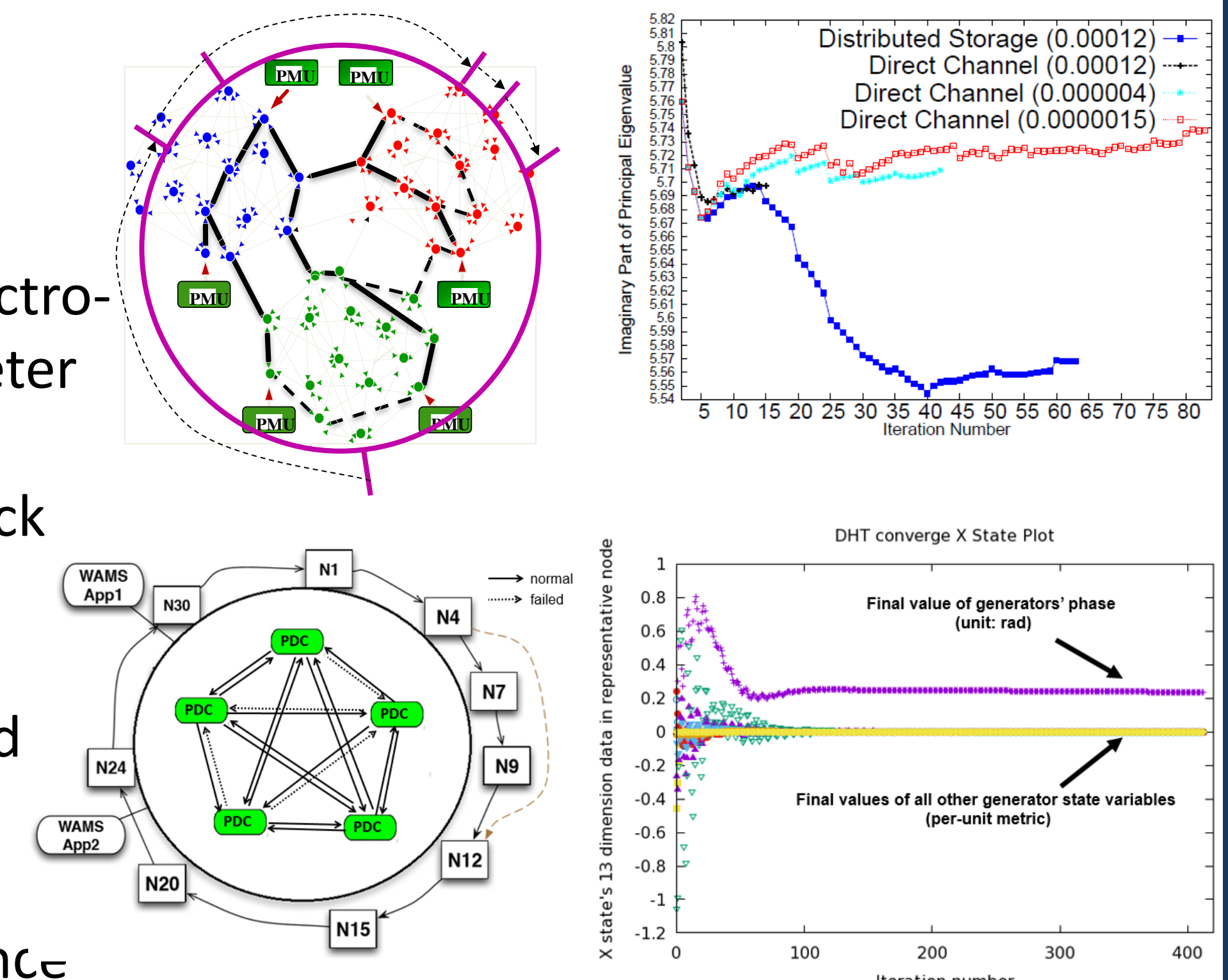
Sparse distributed optimization:

1. ADMM, Stochastic Gradient
2. Round-Robin ADMM for fault tolerant optimization

Distributed Middleware

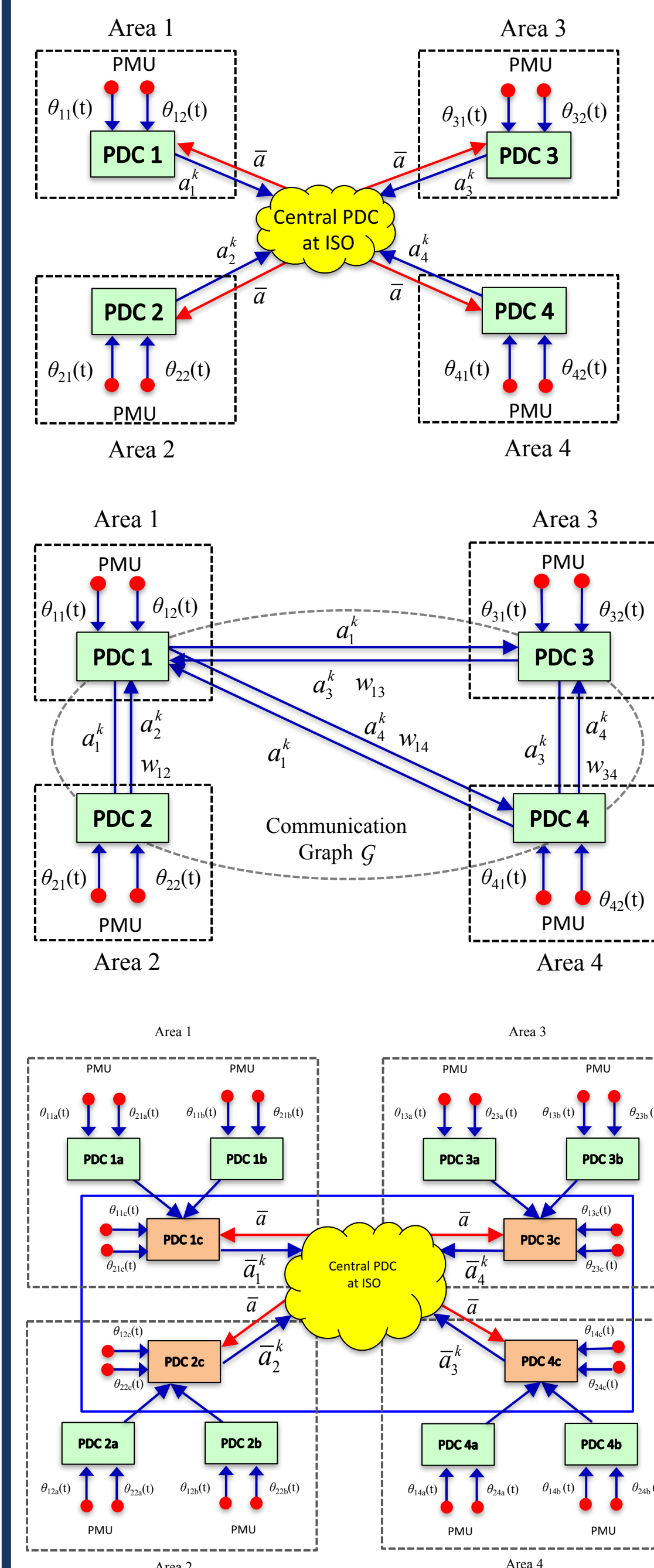
Integrated two classes of important applications into Resilient Real Time Distributed System (R2TDS):

- Single server measurement for electro-mechanical oscillation mode parameter estimate
- Fully distributed wide-area feedback control
- Alternative access points of R2TDS guarantee continual data storage and retrieval
- The decreasing of convergence accuracy and speed or even divergence can be avoided by using R2TDS



Technical Approach

Proposed Distributed Cyber-Physical Architecture for PMU-PDC Communication:



Dynamic Rate Control Problem (DRCP):

- Find optimal PMU data exporting rates, and frequency of information exchange between local PDCs and inter-regional PDCs to minimize computation error between centralized and distributed estimation

Dynamic Link Assignment Problem (DLAP):

- Find optimal communication topologies in real-time connecting local and inter-regional PDCs to maximize computational speed for the overall global estimation/monitoring/control problem.

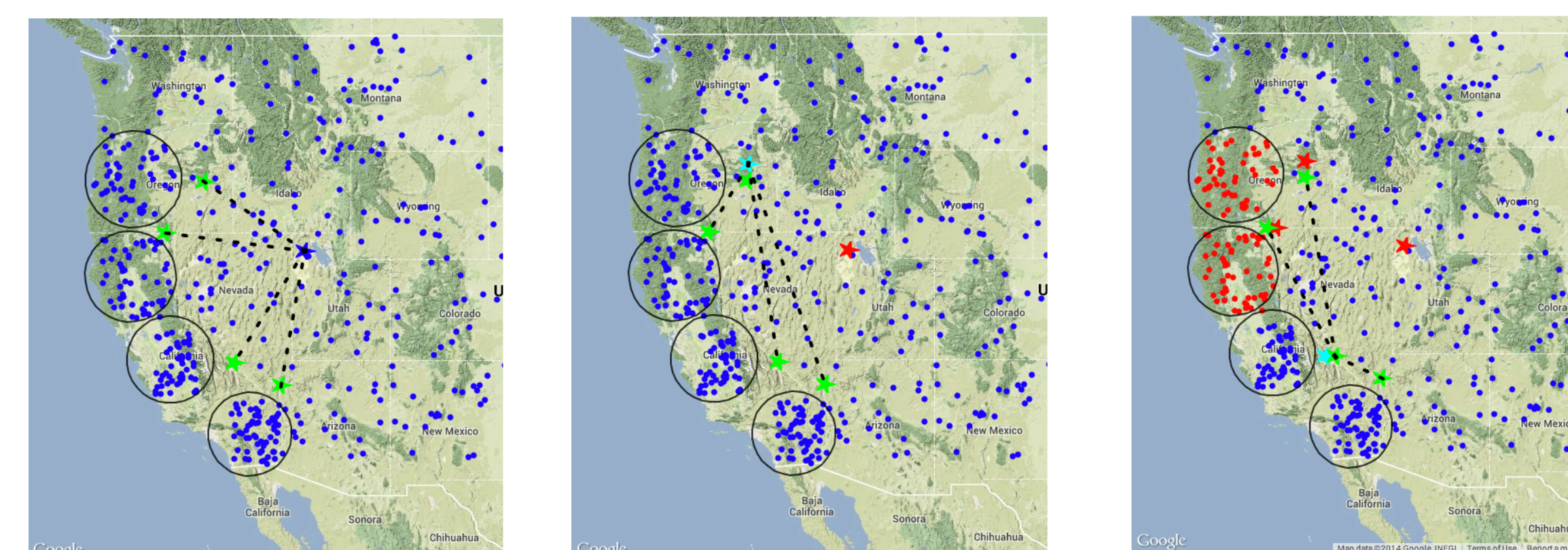
New Algorithms and Results using ADMM

- 1: Update both primal and dual estimation variables at every local control center:

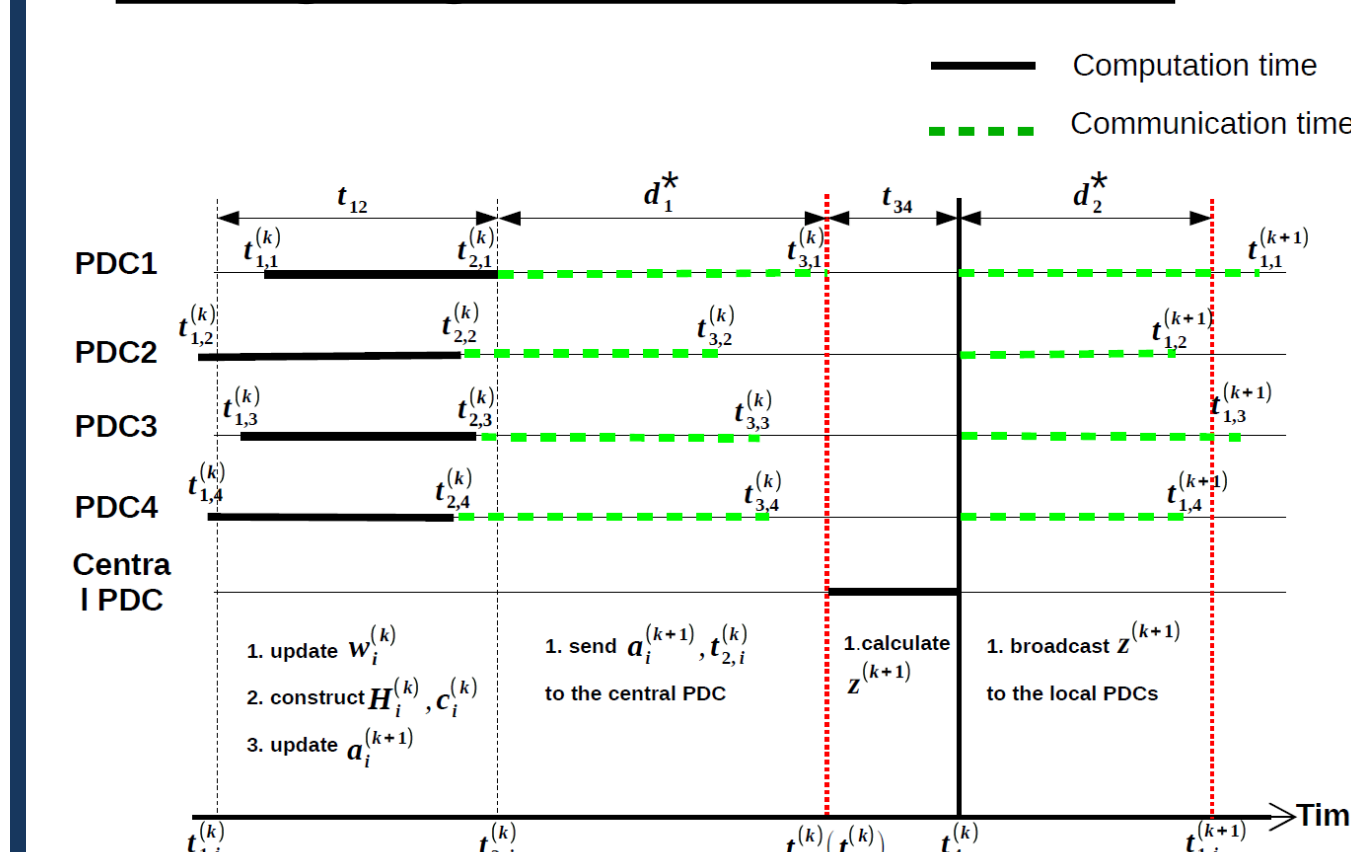
$$\beta_i^{(k+1)} = ((H_i^{(k)})^T H_i^{(k)} + \rho I)^{-1} ((H_i^{(k)})^T c_i^{(k)} - w_i^{(k)} + \rho \bar{\beta}^{(k)})$$

$$w_i^{(k+1)} = w_i^{(k)} + \rho(\beta_i^{(k+1)} - \bar{\beta}^{(k+1)})$$

- 2: Gather the values of $\beta_i^{(k+1)}$ at the central ISO
- 3: Compute the average of $\beta_i^{(k+1)}$ at the central ISO
- 4: Broadcast the average to local control centers and iterate to Step 1



Timing diagram of message arrival



Mitigating Asynchrony in Wide-Area Communication

If a message doesn't arrive at ISO by a delay threshold d_1^*

- Strategy 1: Skip missing data

$$z^{(k+1)} = \frac{1}{S_1^{(k)}} \sum_{i \in S_1^{(k)}} (a_i^{(k+1)} + \frac{1}{\rho} w_i^{(k)})$$

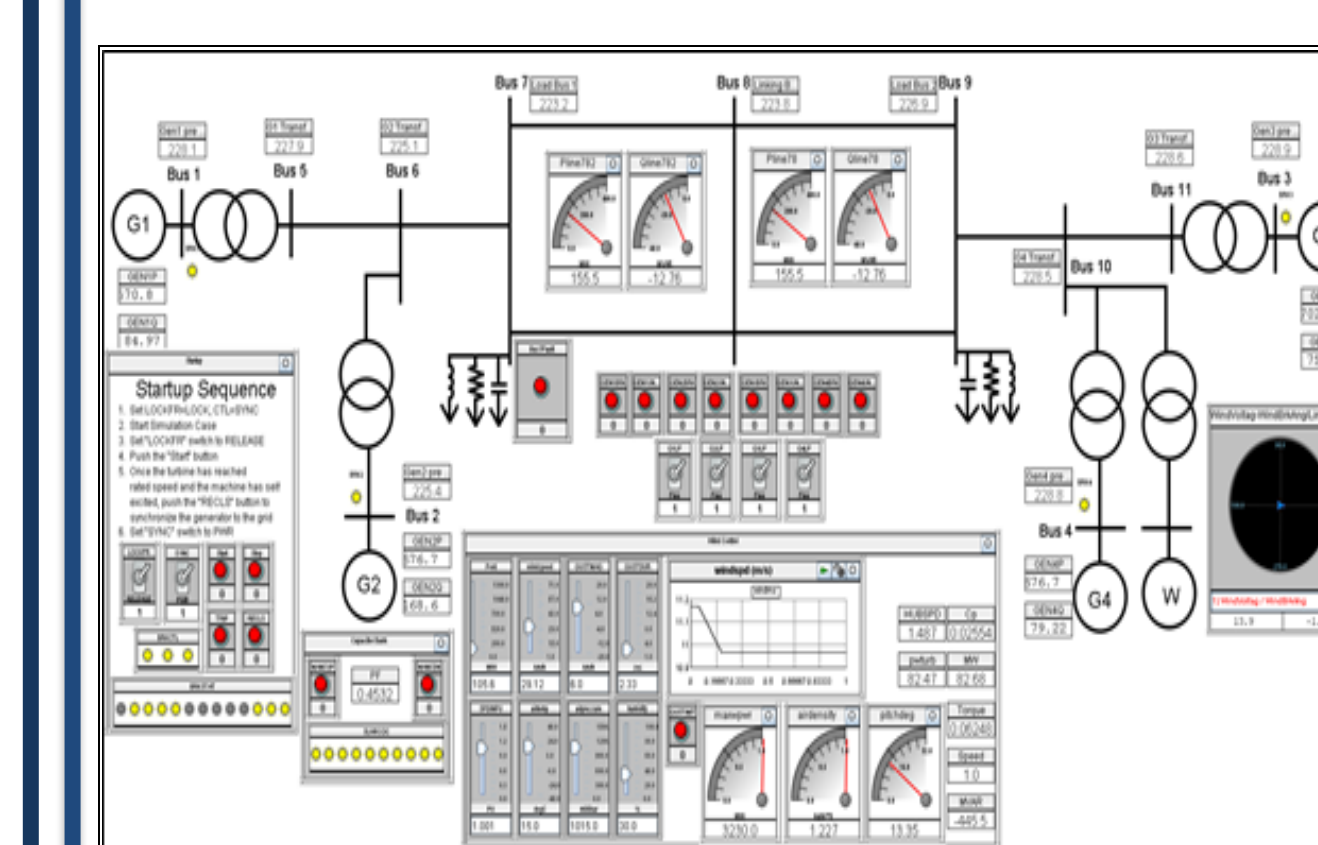
- Strategy 2: Use zero-order hold

$$z^{(k+1)} = \frac{1}{N} \left(\sum_{i \in S_1^{(k)}} (a_i^{(k+1)} + \frac{1}{\rho} w_i^{(k)}) + \sum_{i \in S_2^{(k)}} (a_i^{(k)} + \frac{1}{\rho} w_i^{(k-1)}) \right)$$

Fault-Tolerance & Cyber-Security

- Design application specific fault-tolerance mechanisms to meet real-time needs of the DRCP and DLAP monitoring algorithms
- Crash failures
- Byzantine failures – *Graph designs* to prevent arbitrary byzantine faults
- Leverage the redundancy of sensors and the correlation among sensor data to reduce the cost of fault-tolerance
- Protecting a small subset of PMU data may be necessary and sufficient to detect false data injection attacks
- Leverage application characteristics to design approximate or safe algorithms that can tolerate asynchrony and message loss

Experimental Testbed



- Participated in US Ignite Application Summits 2013-2017, and in Smart America Challenge 2014 Initiative of NIST and US White House
- Federated ExoGENI-WAMS: Multi-vendor PMU-based hardware-in-loop simulation testbed at NCSU to showcase wide-area oscillation monitoring and control

Broader Impacts

- Smart America Challenge 2014, US Ignite 2013-2017 demos
- Undergraduate, K-12 and minority education via Science House and FREEDM ERC programs at NC State
- Undergraduate summer internship at Information Trust Institute at UIUC
- Industry collaborations with power utilities and vendors such as SCE