

# Fast Reachability As a Building Block for Verified Autonomy

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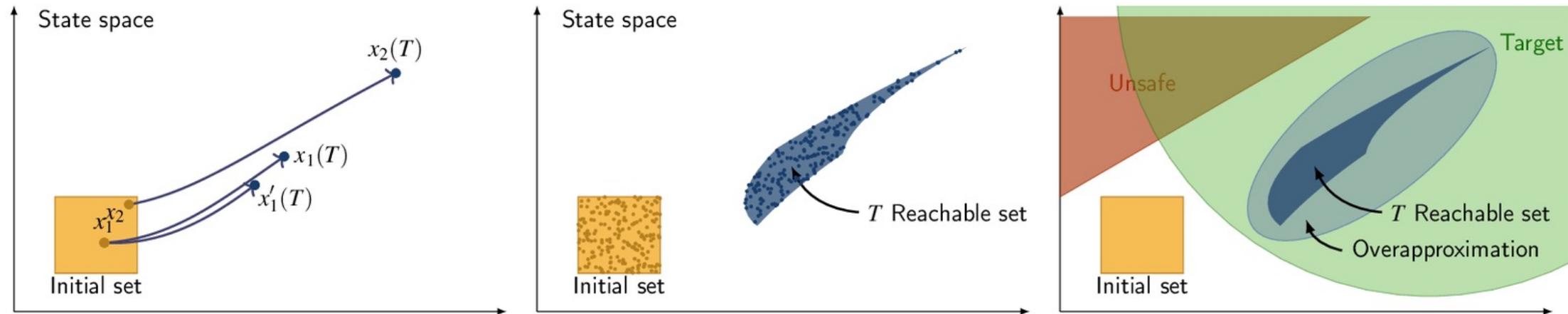


# Reachable sets of dynamical systems

**System:**  $\dot{x} = f(x, w)$

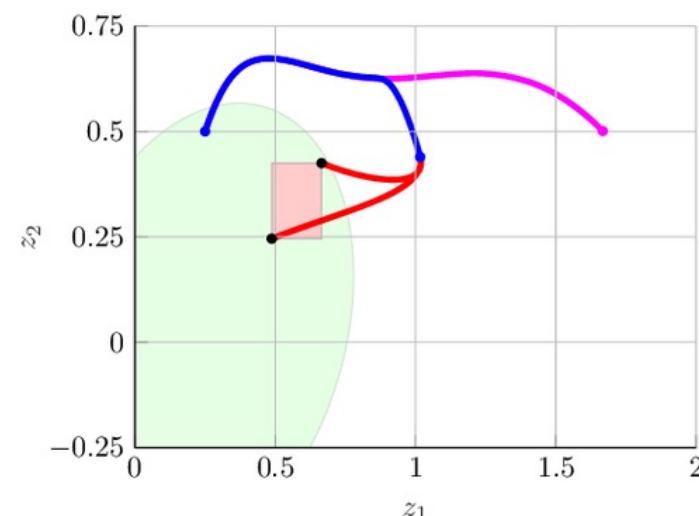
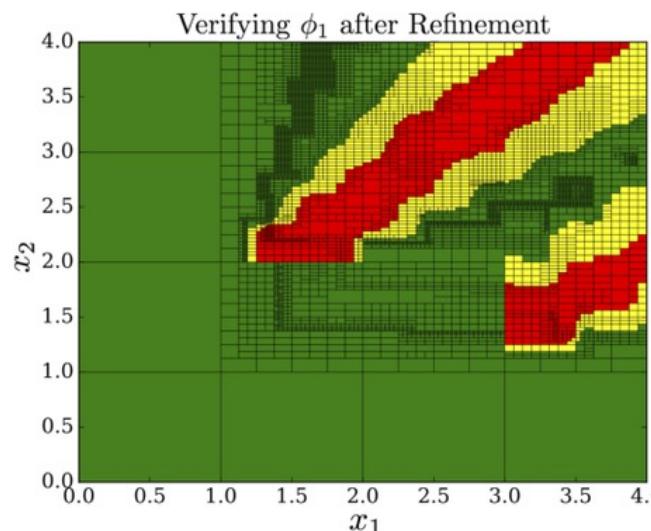
**State:**  $x \in \mathbb{R}^n$

**Disturbance:**  $w \in \mathcal{W}$



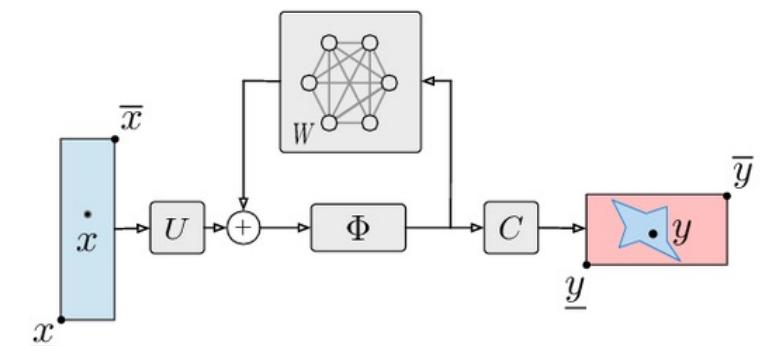
- ▶ Reachable sets characterize possible system evolution
- ▶ Overapproximations of reachable sets are appropriate for verification and safety

# Need for fast reachability methods



For **formal methods**:  
Reachability from each  
region of a finite abstraction

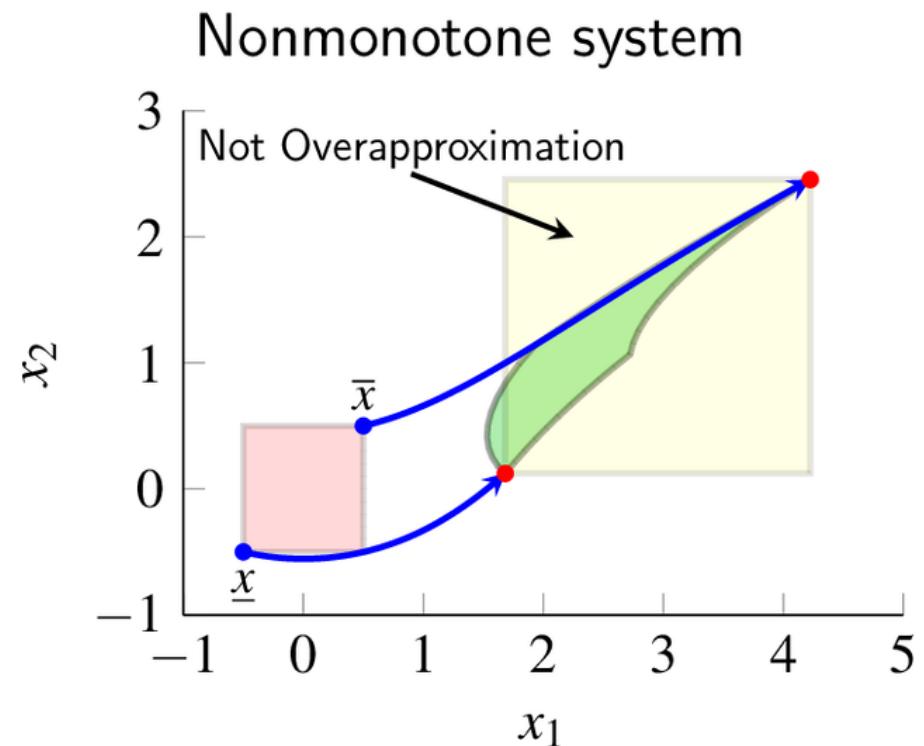
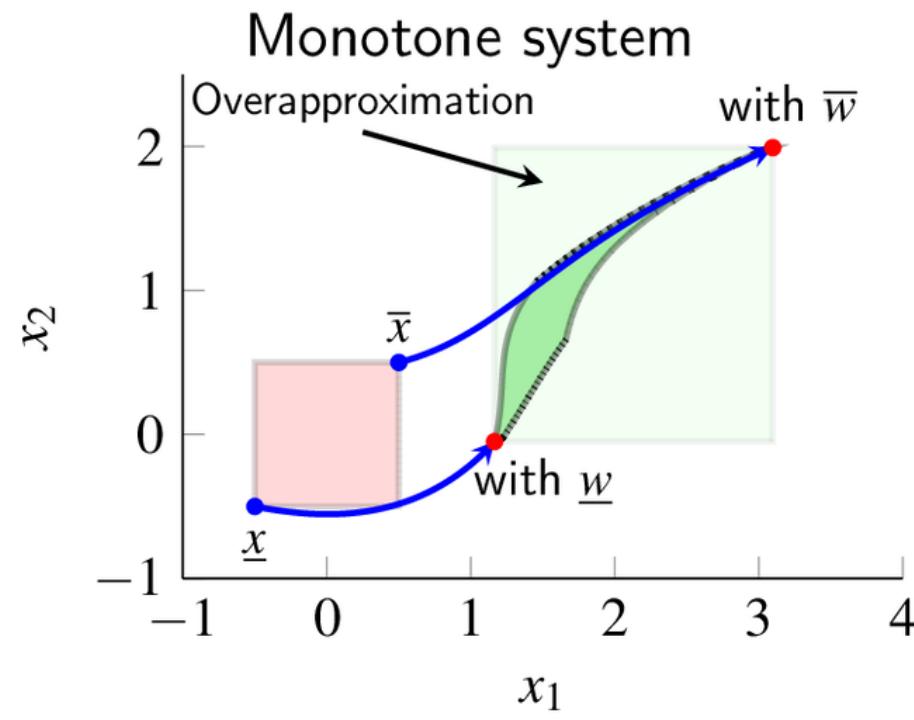
For **safe autonomy**:  
Reachability in the control  
loop for runtime assurances



For **NN verification**:  
High dimensional  
reachability

# Mixed monotonicity for interval reachability estimates

*Reachability analysis for monotone systems.* For a monotone system,  
Reachable set  $\subseteq [\text{lower trajectory}, \text{upper trajectory}]$ .



Goal of **mixed monotonicity**: Embed nonmonotone system in a monotone system

# Reachability from embedding system

System:  $\dot{x} = f(x, w)$ , disturbance input  $w \in \mathcal{W} = [\underline{w}, \bar{w}] = \{w : \underline{w} \leq w \leq \bar{w}\}$



2n dimensional embedding system:  $\begin{bmatrix} \dot{x} \\ \dot{\hat{x}} \end{bmatrix} = e(x, \hat{x}) := \begin{bmatrix} d(x, \underline{w}, \hat{x}, \bar{w}) \\ d(\hat{x}, \bar{w}, x, \underline{w}) \end{bmatrix}$

- ▶ *d* is a *decomposition function* constructed from the dynamics *f*.
- ▶ **MM is fast:** A single trajectory of the deterministic embedding bounds reachable sets of the original mixed monotone system
- ▶ **MM is scalable:** If  $(\underline{x}_{\text{eq}}, \bar{x}_{\text{eq}})$  is an equilibrium for embedding system, then hyperrectangle  $[\underline{x}_{\text{eq}}, \bar{x}_{\text{eq}}]$  is robustly forward invariant

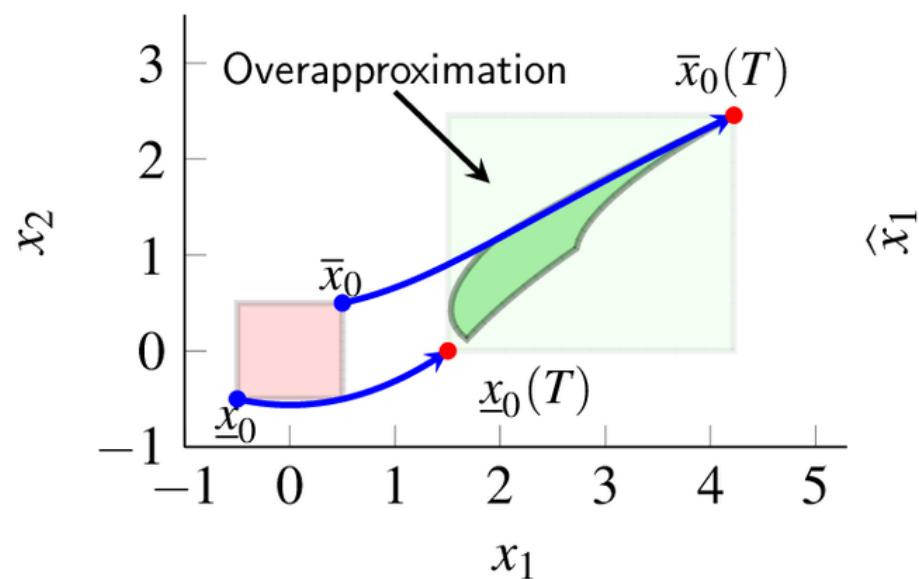
# A simple example

## Mixed Monotone System:

$$\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \end{bmatrix} = \begin{bmatrix} x_2^2 + 2 \\ x_1 \end{bmatrix}$$

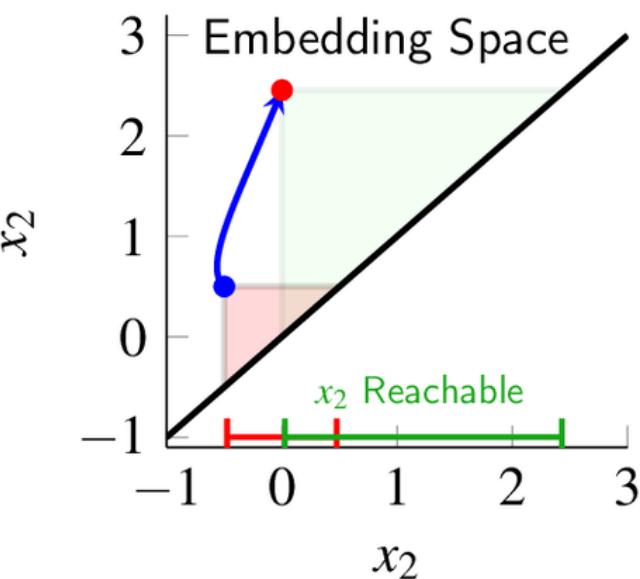
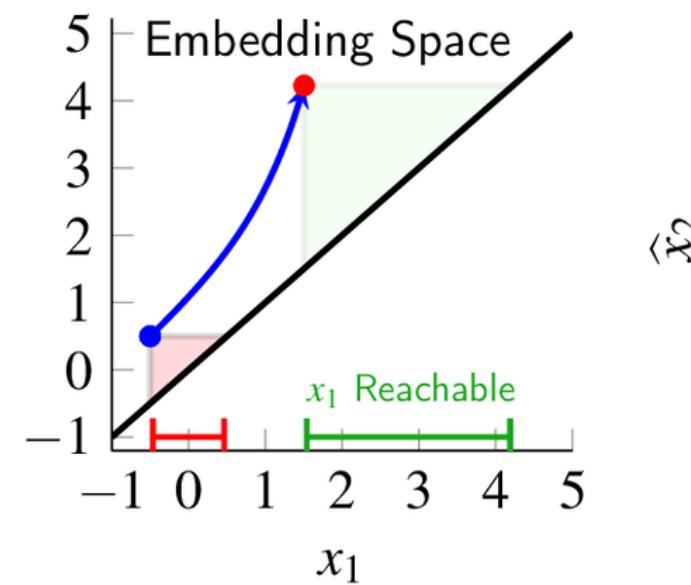
$$[\underline{x}, \bar{x}] = [(-0.5, -0.5), (0.5, 0.5)]$$

$$T = 1$$



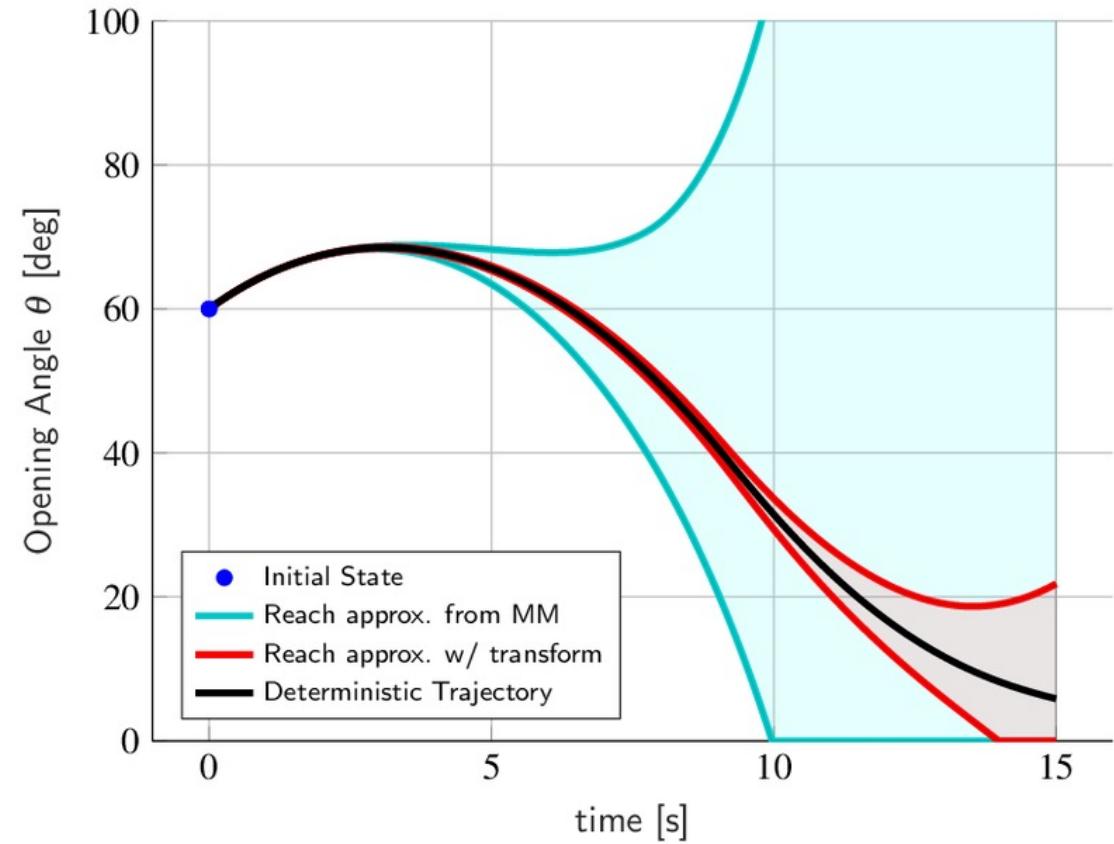
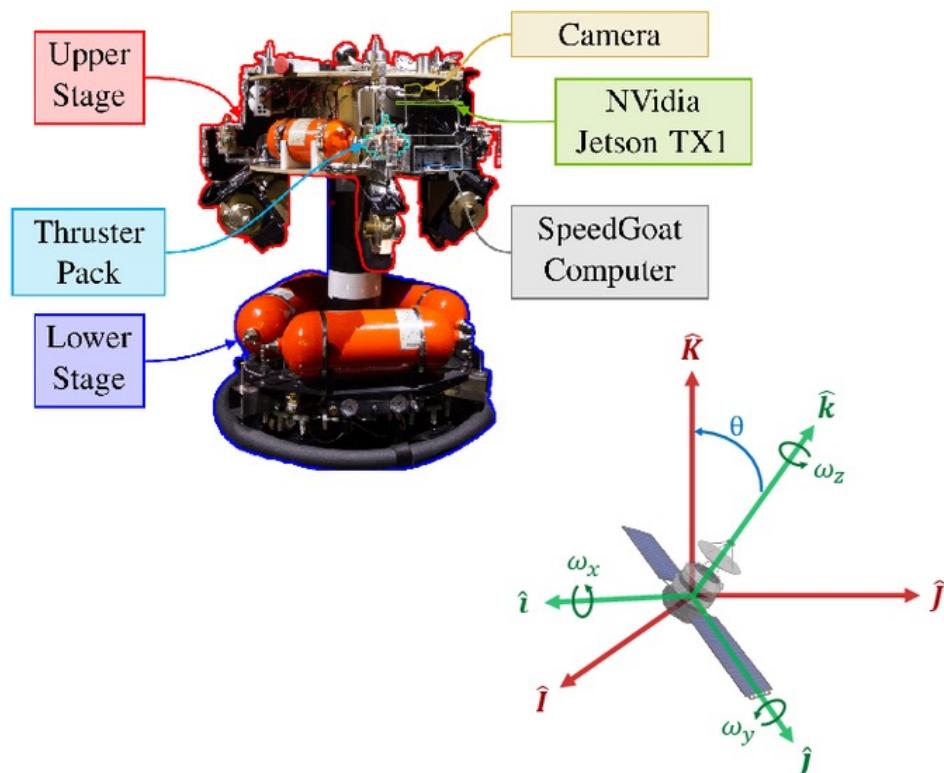
## Decomposition Function:

$$d_1(x, \hat{x}) = \begin{cases} x_2^2 + 2 & \text{if } x_2 \geq 0 \text{ and } x_2 \geq -\hat{x}_2, \\ \hat{x}_2^2 + 2 & \text{if } \hat{x}_2 \leq 0 \text{ and } x_2 < -\hat{x}_2, \\ 2 & \text{if } x_2 < 0 \text{ and } \hat{x}_2 > 0. \end{cases}$$
$$d_2(x, \hat{x}) = x_1$$



# Fast reachability for runtime assurance (RTA) from mixed monotonicity

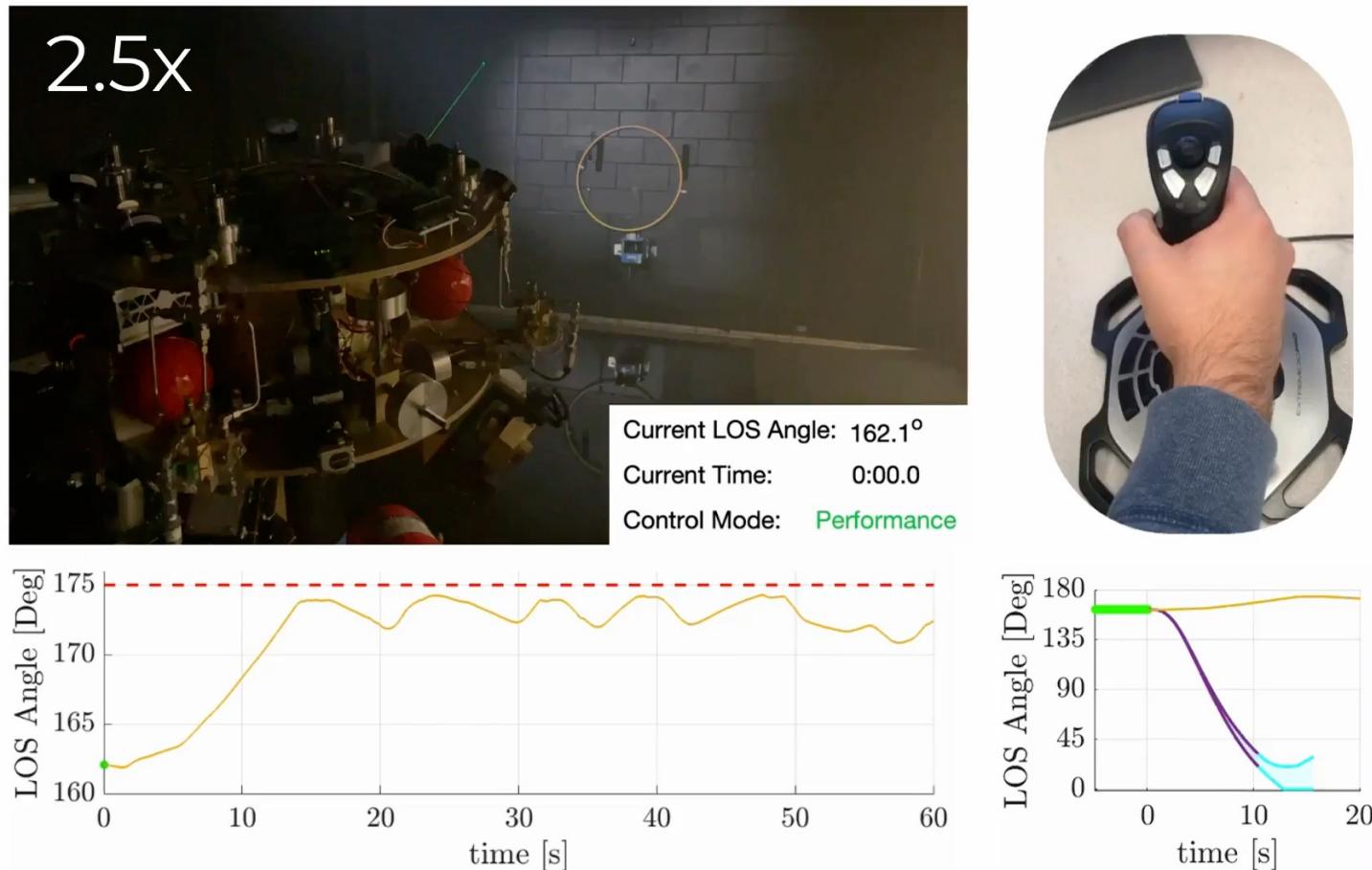
- ▶ Torque-controlled nonlinear spacecraft system in free rotational motion
- ▶ Safety objective: Line-of-sight angle constraint,  $\theta \leq \theta_{max}$



M. Abate, M. Mote, M. Dor, C. Klett, S. Phillips, K. Lang, P. Tsiotras, E. Feron, S. Coogan, in submission.

# Fast reachability for runtime assurance (RTA) from mixed monotonicity

- RTA mechanism overrides human input when reachable set could become unsafe



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# Fast reachability for formal methods verification

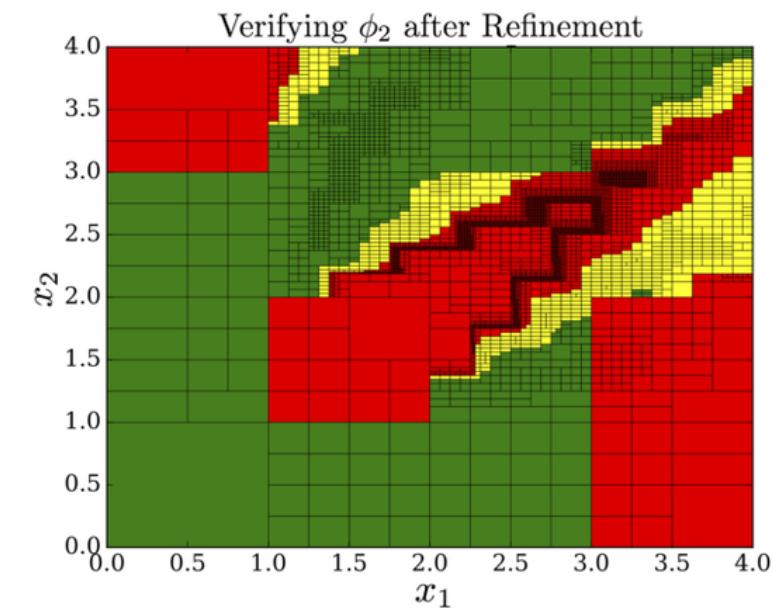
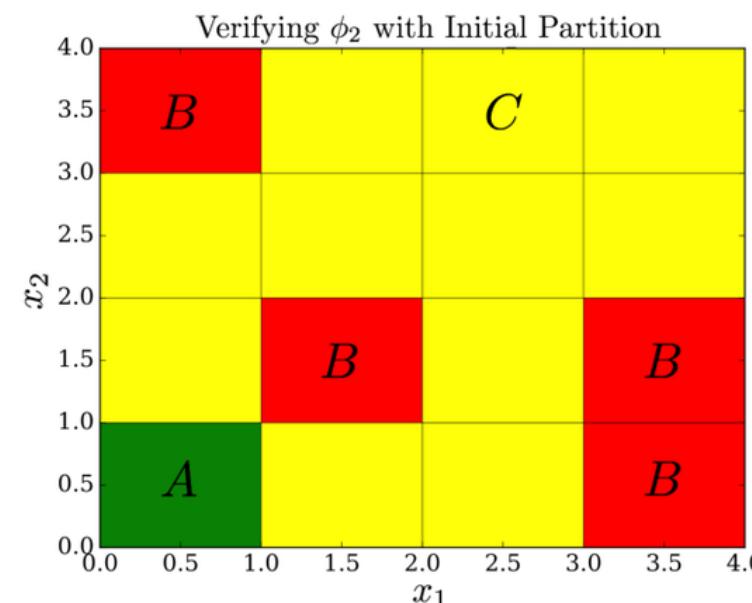
Less than 90% chance of: reaching a  $B$  state if it eventually always remains in  $A$  state,  
and always stays outside of  $B$  state if it reaches a  $C$  state

$$\mathcal{P}_{\leq 0.90}[(\Diamond \Box A \rightarrow \Diamond B) \wedge (\Diamond C \rightarrow \Box \neg B)]$$

$$x_1^+ = x_1 + (-ax_1 + x_2) \cdot \Delta T + w_1$$

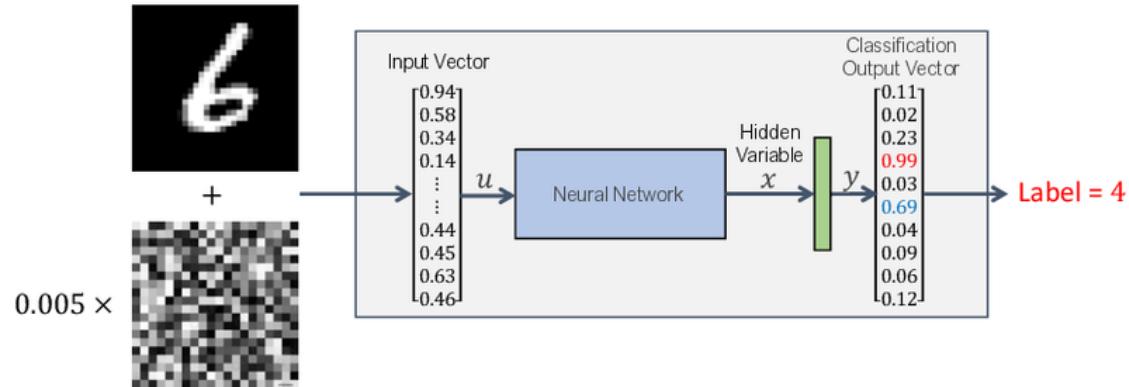
$$x_2^+ = x_2 + \left( \frac{(x_1)^2}{(x_1)^2 + 1} - bx_2 \right) \cdot \Delta T + w_2$$

$w_1, w_2$  truncated Gaussian



# Scalable mixed monotonicity for robustness analysis/training in NNs

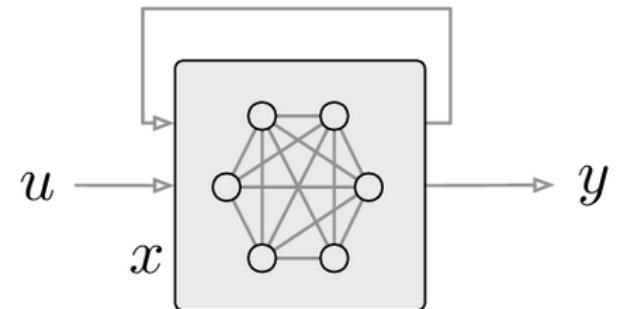
A well known problem: Neural networks are brittle



Increased interest in *Implicit* Neural Networks defined by fixed point equation:

$$x = \Phi(Ax + Bu)$$

$$xy = Cx$$

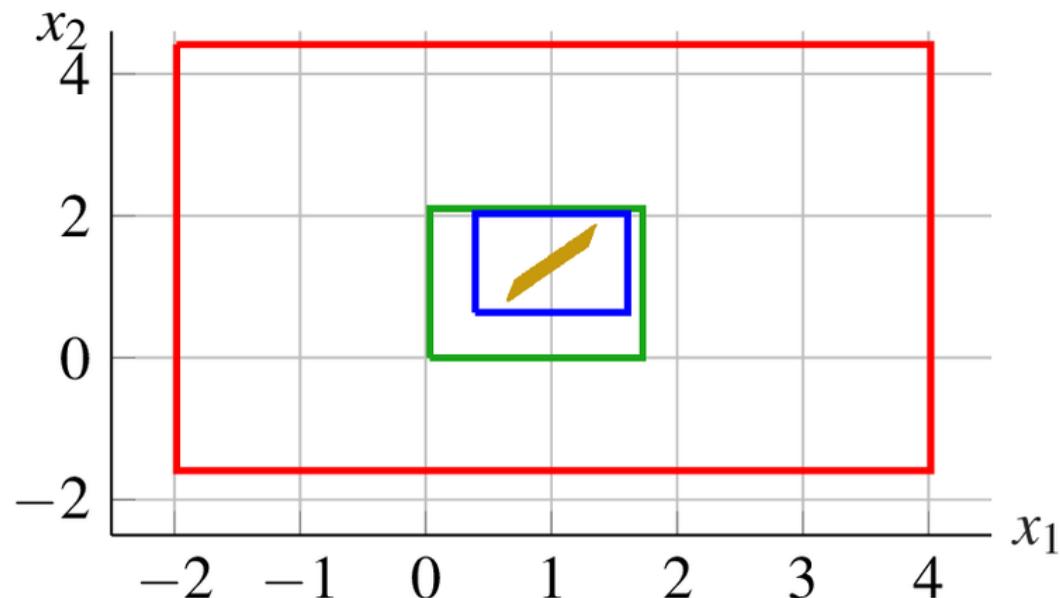


Mixed monotonicity provides computationally and theoretically scalable local robustness certificates

## A simple example

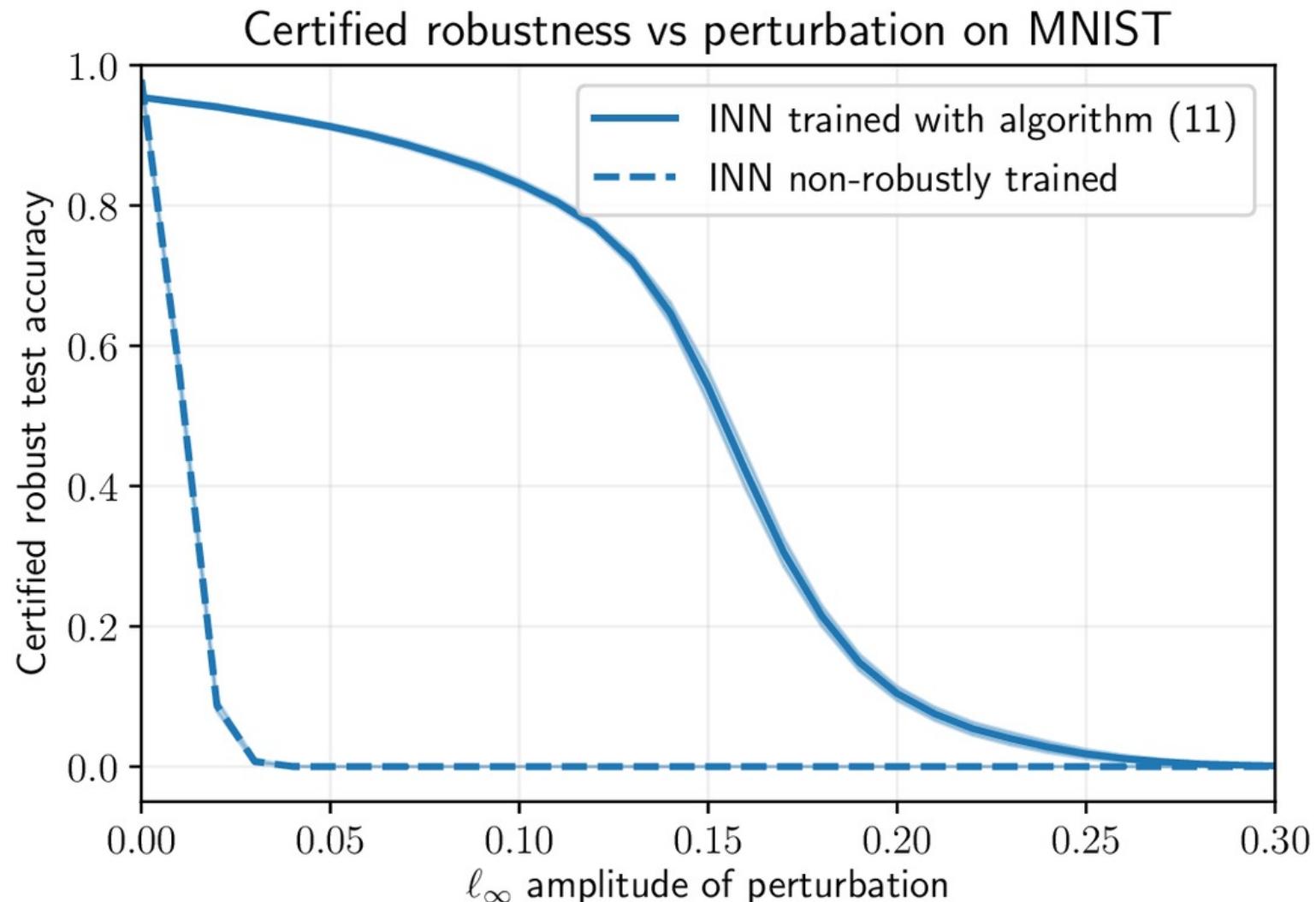
Consider  $x = \Phi(Ax + Bu)$  with output  $y = x$  and

$$A = \begin{bmatrix} -0.25 & -0.25 \\ 0.75 & -0.25 \end{bmatrix}, B = \begin{bmatrix} 0.5 & 1 \\ 1 & 0.5 \end{bmatrix}, \Phi\left(\begin{bmatrix} s_1 \\ s_2 \end{bmatrix}\right) = \begin{bmatrix} \text{ReLU}(s_1) \\ \text{ReLU}(s_2) \end{bmatrix}, \underline{u} = \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \bar{u} = \begin{bmatrix} 1/3 \\ 2 \end{bmatrix}$$



- ▶ Red: Lipschitz bound
- ▶ Green: Interval bound propagation
- ▶ Blue: Embedding equilibrium

# Application to MNIST handwriting dataset



# Concluding thoughts

- ▶ Reachability analysis is a fundamental building block for verified autonomy
- ▶ Need for fast and scalable approximation methods
- ▶ Some challenges:
  - ① NN in closed-loop with control systems
  - ② Accommodating learning of dynamics in the reachable set computations
  - ③ Trade-off between offline computation of safe regions vs. online/runtime intervention
  - ④ Accommodating more complex safety specs than invariance (e.g., temporal logic)
  - ⑤ How to intervene to maintain safety (e.g., CBFs, switch to backup controller)
  - ⑥ Geometries besides rectangles

Papers available at  
[coogan.ece.gatech.edu](http://coogan.ece.gatech.edu)

