

NSF CPS PI Meeting, November 17, 2015

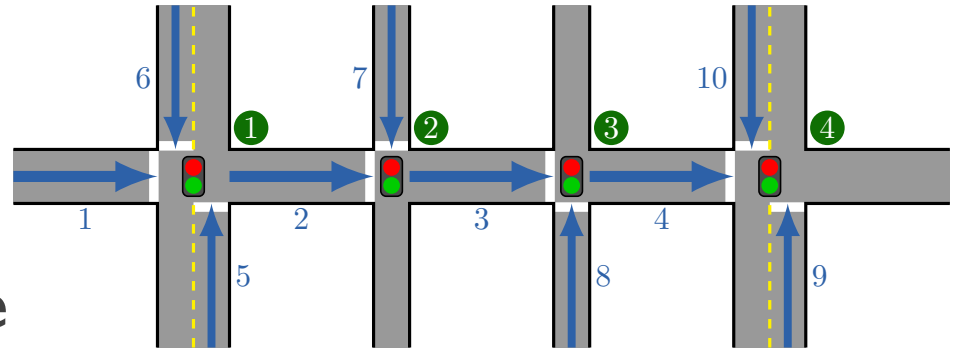
# Formal Synthesis for Traffic Control

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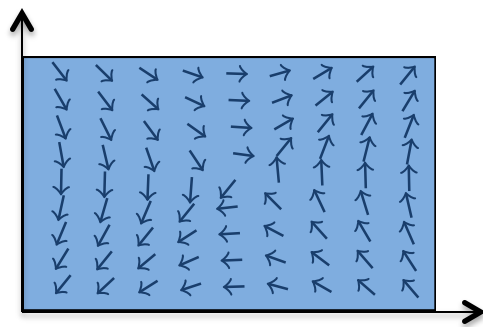
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**Objective:** Traffic control to meet objectives expressed in temporal logic, e.g.,

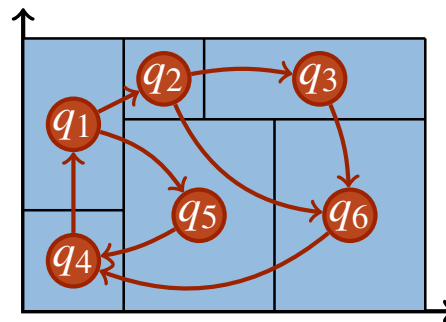
- eventually each link will have  $\leq 30$  vehicles
- upstream link will have low demand until downstream link is no longer congested
- each queue at a junction will be actuated at least once every two minutes.



## Formal Synthesis Workflow



$$x^+ = F(x)$$



finite abstraction



formal methods

## Outline:

### 1. Finite abstraction for formal methods

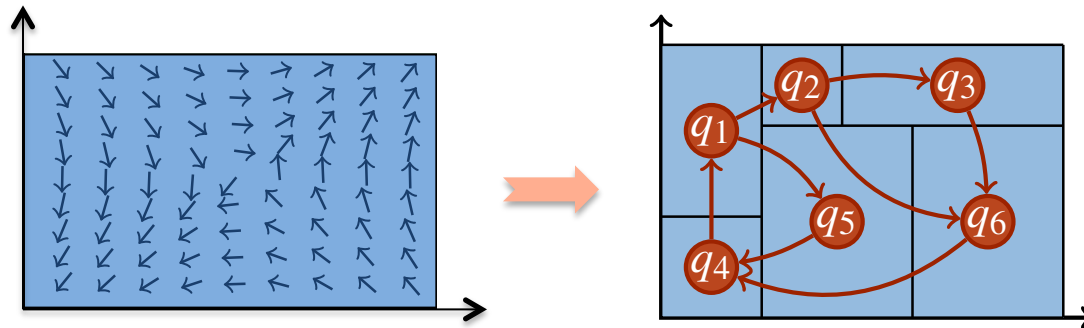
*Exploiting a “mixed monotonicity” property for scalability.*

*Application to a macroscopic traffic flow model.*

### 2. Compositional synthesis for large networks

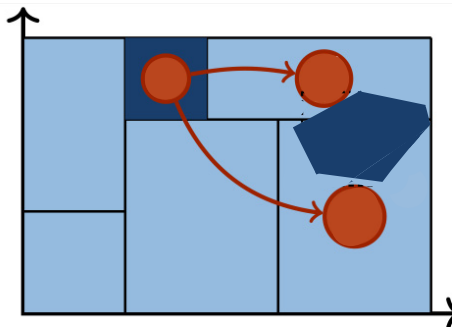
*Decoupled synthesis for subnetworks with supply and demand contracts.*

# I. Finite Abstraction for Formal Methods



Capture the underlying dynamics with a finite set of symbols and transitions between them. Methods exist for classes of systems (Tabuada, Girard, Pappas, Reissig, Abate, Belta, and others...)

**Example:** polyhedral computations for piecewise affine systems (Belta et al.)



# Monotonicity and Mixed Monotonicity

The discrete-time system:

$$x^+ = F(x) \quad x \in \mathcal{X}$$

is **monotone** if

$$x_1 \leq x_2 \quad \Longrightarrow \quad F(x_1) \leq F(x_2)$$

with respect to a partial order (standard order in this talk).

Monotonicity offers strong dynamical properties [Hirsch, Smith, Angeli, Sontag] but is restrictive in practice.

Necessary and sufficient condition for monotonicity:

$$\frac{\partial F_i(x)}{\partial x_j} \geq 0 \quad \forall x \in \mathcal{X} \quad \forall i, j$$

# Monotonicity and Mixed Monotonicity

$$x^+ = F(x) \quad x \in \mathcal{X}$$

is **mixed monotone** if there exists a “decomposition function”

$$f : \mathcal{X} \times \mathcal{X} \rightarrow \mathcal{X}$$

such that  $f(x, x) = F(x)$

$$x_1 \leq x_2 \Rightarrow f(x_1, y) \leq f(x_2, y)$$

$$y_1 \leq y_2 \Rightarrow f(x, y_2) \leq f(x, y_1).$$

A sufficient condition for mixed monotonicity:

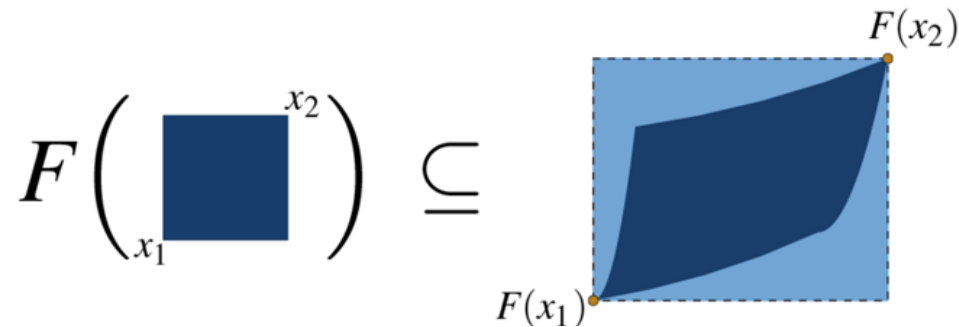
$$\exists \delta_{ij} \in \{-1, 1\} \quad \text{s.t.} \quad \delta_{ij} \frac{\partial F_i(x)}{\partial x_j} \geq 0 \quad \forall i, j$$

Decomposition function:  $F_i(\dots, \cancel{x_j}^{y_j}, \dots)$  if  $\delta_{ij} = -1$

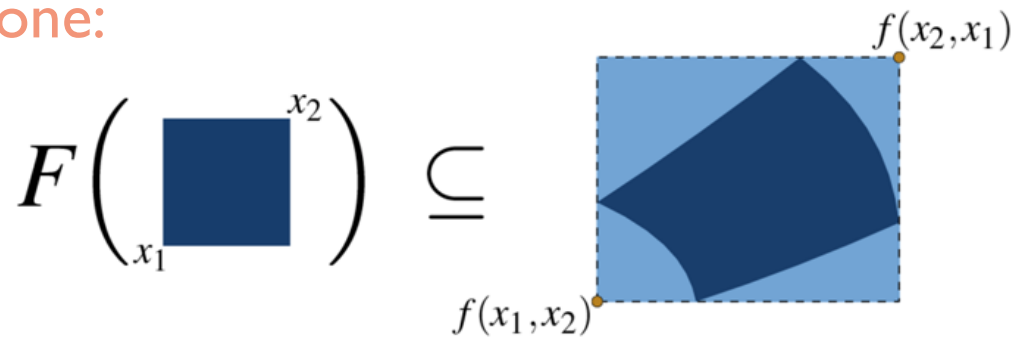
# Mixed Monotonicity Allows Scalable Finite Abstraction

Two function evaluations tightly bound the one-step reach set:

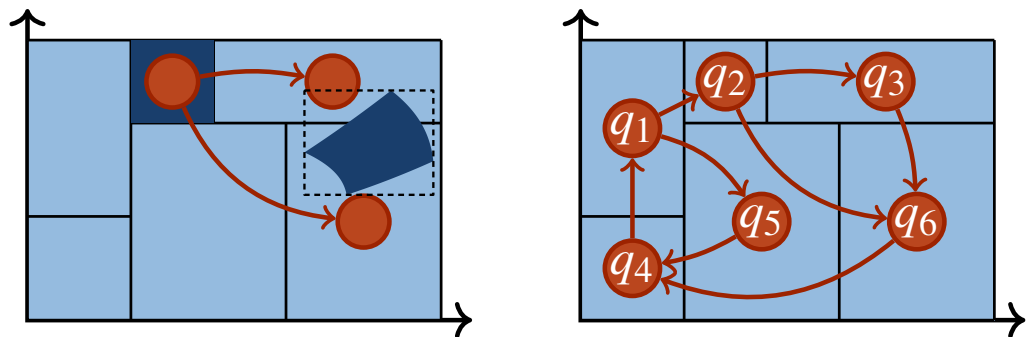
Monotone:



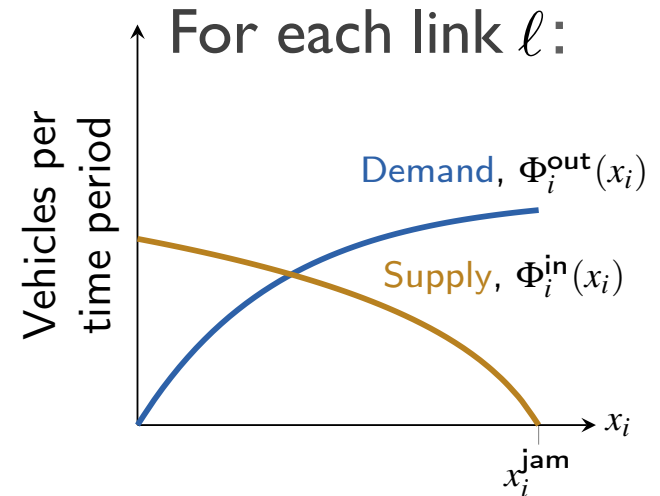
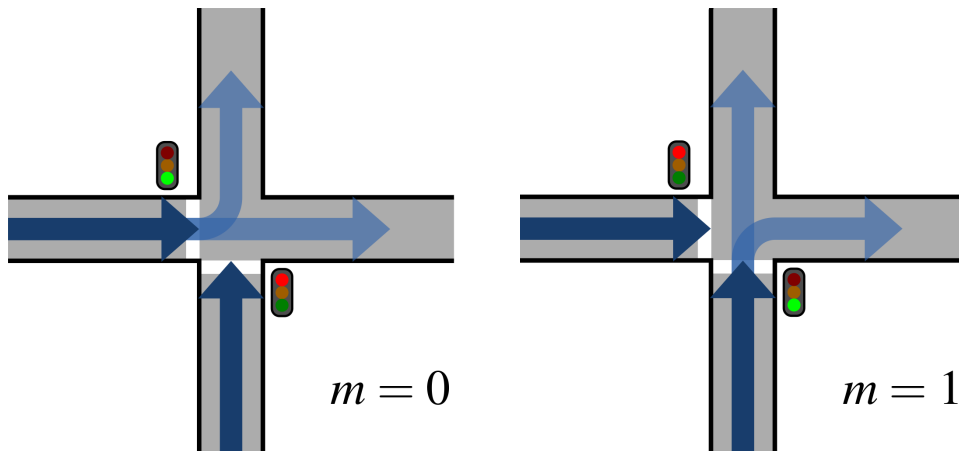
Mixed Monotone:



This allows a scalable abstraction algorithm:



# Traffic Flow: a Macroscopic Model



$$x_\ell^+ = x_\ell + f_\ell^{\text{in}}(x) - f_\ell^{\text{out}}(x) =: F_\ell(x)$$

Outgoing links:  $f_k^{\text{in}}(x, m) = \sum_{\ell \in \text{in}} \beta_{\ell k} f_\ell^{\text{out}}(x, m)$

turn ratio

Incoming links:

$$f_\ell^{\text{out}}(x, m) = s_\ell(m) \min \left\{ \Phi_\ell^{\text{out}}(x_\ell), \min_{k \in \text{out}} \frac{1}{\beta_{\ell k}} \Phi_k^{\text{in}}(x_k) \right\}$$

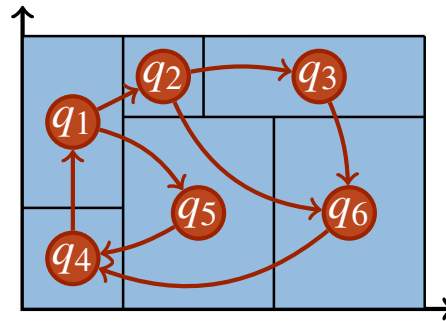
$\{0, 1\}$



# Traffic Flow is Mixed Monotone

$$\delta_{ij} \frac{\partial F_i(x)}{\partial x_j} \geq 0 \quad \delta_{ij} = \begin{cases} -1 & \text{if } i \text{ and } j \text{ share tail node} \\ +1 & \text{otherwise} \end{cases}$$

Apply abstraction algorithm and add signaling states to transition model

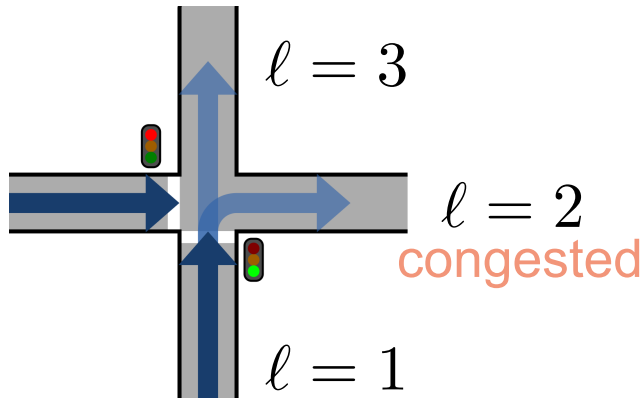


finite abstraction



formal methods

**Note:** Standard monotonicity breaks down at splits

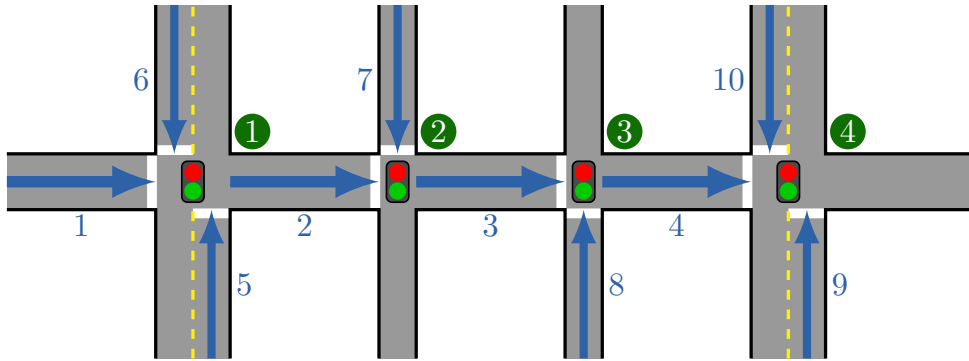


$$f_1^{\text{out}} = \frac{1}{\beta_{12}} \Phi_2^{\text{in}}(x_2)$$

$$f_3^{\text{in}} = \beta_{13} f_1^{\text{out}} = \frac{\beta_{13}}{\beta_{12}} \Phi_2^{\text{in}}(x_2)$$

$$\Rightarrow \delta_{32} = -1$$

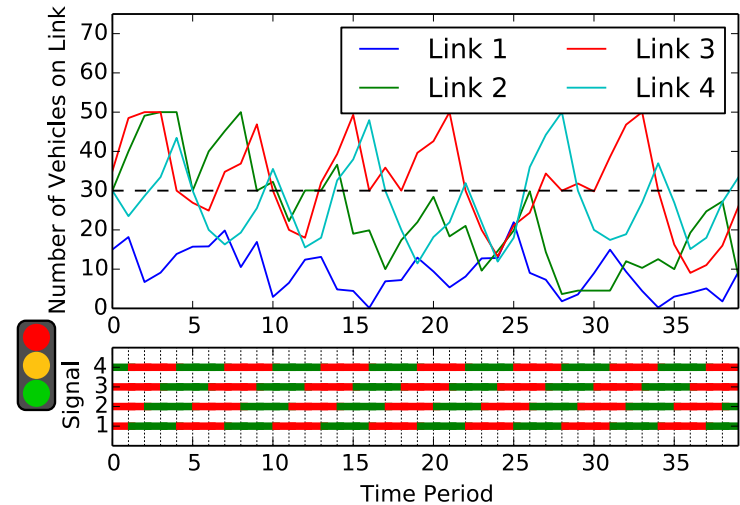
# Example: Signal Control for a Corridor



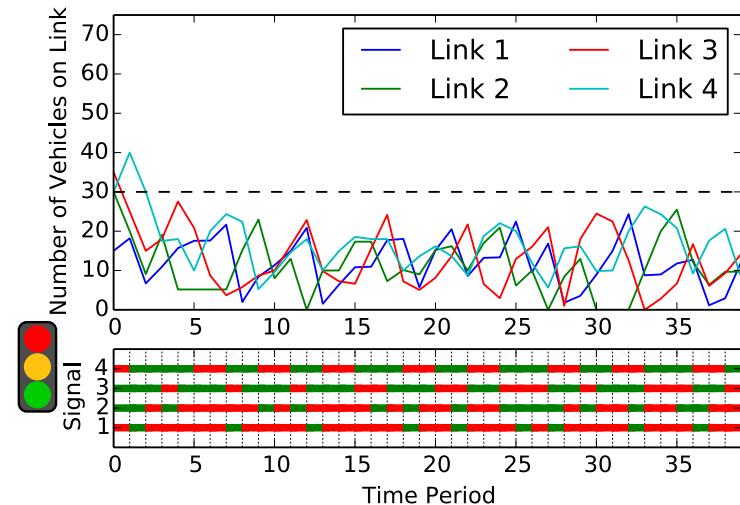
## Temporal Logic Specifications:

- Each signal actuates cross street traffic infinitely often
- Eventually, links 1, 2, 3, and 4 have fewer than 30 vehicles each
- The signal at junction 4 must actuate cross street traffic for at least two sequential time-steps

## Naïve offset optimal policy



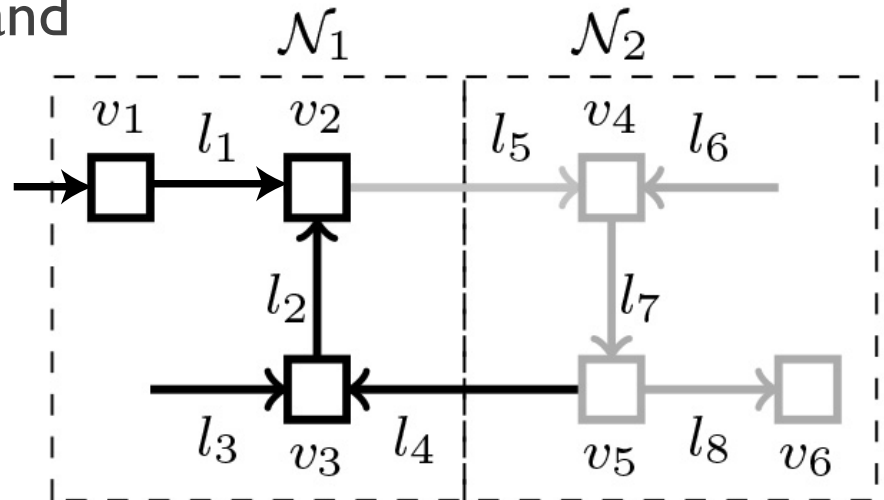
## Correct-by-design policy



## 2. Compositional Synthesis for Large Networks

- “Contracts” between neighboring subnetworks to limit demand and guarantee adequate supply

- Neighbors’ promises allow decoupled subnetwork models with set valued maps

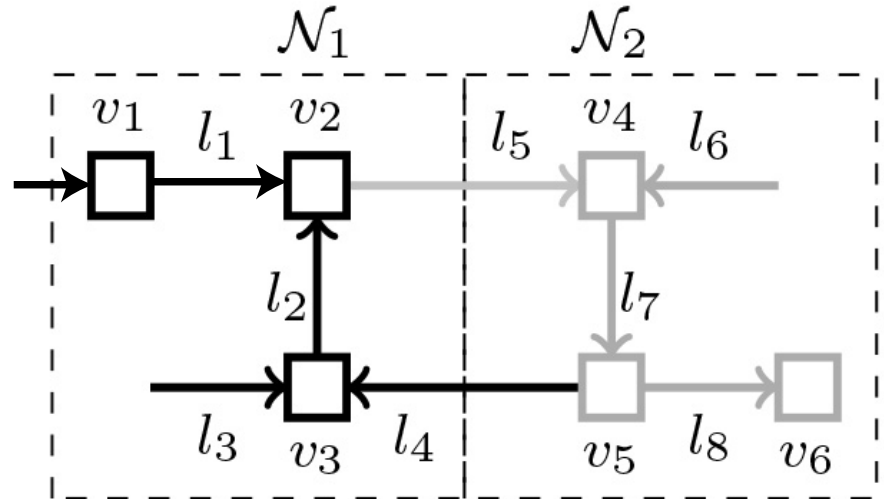


- Augment temporal logic specifications with own promises and synthesize controller for each subnetwork

$$\phi_i^{\text{new}} = \phi_i^{\text{original}} \wedge \underbrace{\phi_i^{\text{supply}} \wedge \phi_i^{\text{demand}}}_{\text{promises to neighbors}} \quad i = 1, 2, \dots$$

# Neighbors' Promises Allow Decoupled Models

Subnetwork 2 promises a minimum supply of  $\sigma_2^{\text{contract}}$  on link 5 and to limit its demand on link 4 by  $\delta_4^{\text{contract}}$  vehicles per period.



Decoupled subnet 1 model:

$$f_2^{\text{out}}(x) = \min \left\{ \Phi_2^{\text{out}}(x_2), \frac{1}{\beta_{25}} \Phi_5^{\text{in}}(x_5) \right\}$$

$$\in \min \left\{ \Phi_2^{\text{out}}(x_2), \frac{1}{\beta_{25}} \sigma \right\}, \quad \sigma \in [\sigma_2^{\text{contract}}, \sigma_2^{\text{best}}]$$

$$f_4^{\text{in}}(x) = \beta_{74} \min \left\{ \Phi_7^{\text{out}}(x_7), \frac{1}{\beta_{78}} \Phi_8^{\text{in}}(x_8), \frac{1}{\beta_{74}} \Phi_4^{\text{in}}(x_4) \right\}$$

$$\in \beta_{74} \min \left\{ \Phi_4^{\text{in}}(x_4), \delta \right\}, \quad \delta \in [0, \delta_4^{\text{contract}}]$$

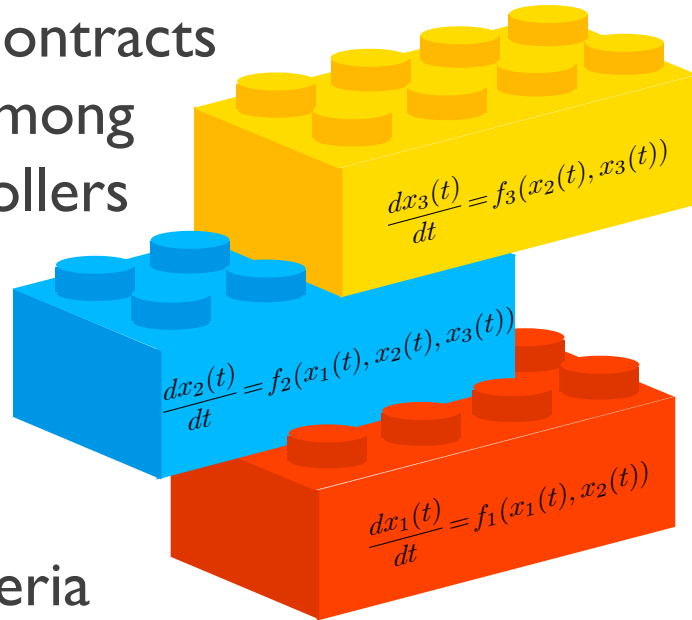
# Ongoing Research

## Compositional synthesis:

less conservative contracts and cooperation among subnetwork controllers rather than fully decentralized control

## Optimality:

add optimality criteria to specifications, e.g., minimize travel time, minimize spatial variations in traffic density, maximize throughput

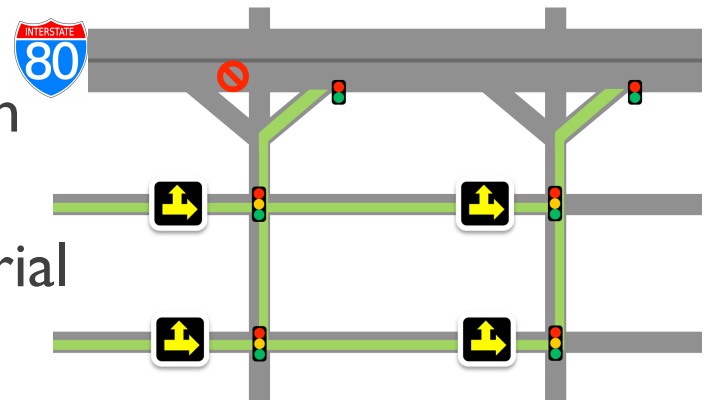


## Probabilistic guarantees:

exploit demand statistics to find transition probabilities and guarantee satisfaction with high probability

## Coordinated onramp metering / arterial signaling

and validation with hybrid freeway/arterial simulation



## Publications This Year

Coogan, Gol, Arcak and Belta “Traffic network control from temporal logic specifications” – IEEE Trans. Control Network Systems, in press

Kim, Arcak and Seshia “Compositional controller synthesis for vehicular traffic networks” – CDC 2015

Sadraddini and Belta “Robust temporal logic model predictive control” – Allerton Conference, 2015

Coogan, Gol, Arcak and Belta “Controlling a network of signalized intersections from temporal logic specifications” – ACC 2015

Coogan and Arcak “Efficient finite abstraction of mixed monotone systems” – HSCC 2015 (Best Student Paper Award)

Sadraddini and Belta “A provably correct MPC approach to safety control of urban traffic networks” – submitted

Coogan, Arcak and Belta “Formal methods for control of transportation networks” – submitted