IskiOS: Lightweight Defense Against Kernel-Level Code-Reuse Attacks

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## Motivation

Commodity operating systems are the most trusted component in the software stack of modern computing platforms, yet, they are vulnerable to code-reuse attacks.

# IskiOS



**Shadow Call Stack** Write-protected **M**Race-free **Efficient** 

Frame 2

**Return Address 2** 

Frame 1

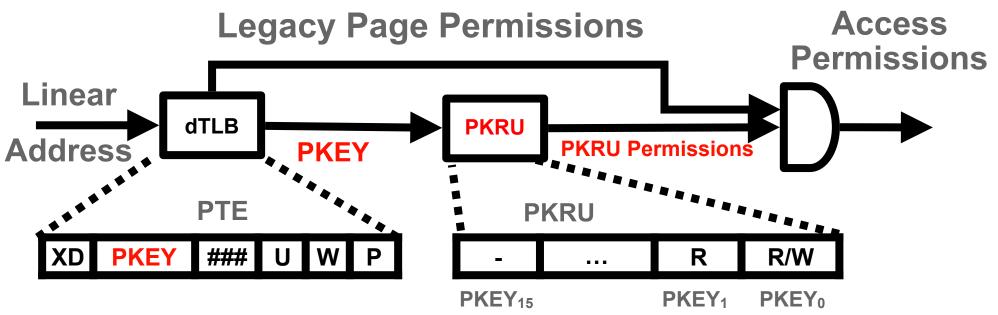
**Return Address 1** 

# Design

### **Intel Protection Keys for Userspace (PKU)**

•**PKRU:** 32-bit user-accessible register

•wrpkru/rdpkru instructions • Up to **16 protection domains** within a singe address space • Applies to PTEs with **U/S bit set** (i.e., user memory only)



#### **Code-reuse Attack**

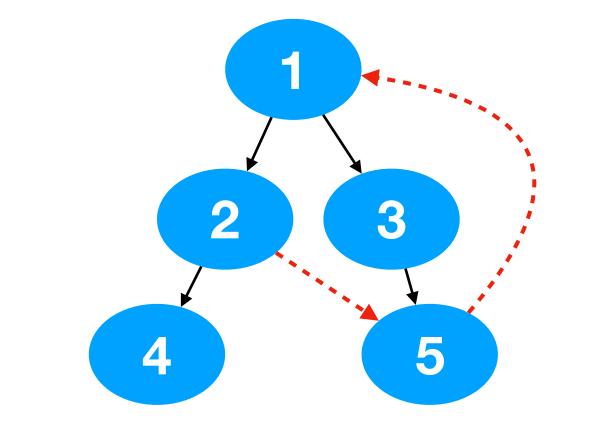
•Goal: control program behavior

• Typically:

-Memory-safety error exploitation (e.g., bufferoverflow)

-Corruption of returnaddress of a function call on stack

- Control-flow redirection to desired code sequence



**Prevents all** *return*-**Return Address 2 Copy** based Return Address 1 Copy code-reuse attacks. **B** Prevents Direct Disclosure of **Code Layout eXecute-Only Memory** (XOM) Code (X) Code cannot be read/ written ☑No layout re-Data (R/W) arrangement **M** Deployed diversification entropy Code (X) is maintained **Prevents all** *just-in*time attacks via direct Data (R/W) reads.

**Kernel Page Table Isolation (KF**  Software-only mitigation for Meltdown attack Separate page tables for user/ke isolation •U/S bit redundant

•Key idea: mark *all* memory as (U/S clear) and rely solely on KF for isolation

 Enables Protection Keys for Kernel (PKK)

## **Evaluation**

Benchmark vanilla	pti	pti+xom	pti+xom+cph	pti+xom+cph+ss-lfo-swo

PTI)	Kernel View	User View
kernel	Kernel Mappings	Unmapped
user PTI	<section-header><section-header></section-header></section-header>	User   Mappings

#### Insufficient Kernel Defenses

 Mostly label-based Controlflow Integrity (CFI) approaches

-Over-permissive controlflow policies due to static analysis imprecision and incompleteness -Lack of return-address

protection



**Trap Padding** foo() Function entries and callq bar callsites are followed/ preceded by random number of traps Pointers in readable memory do not leak information about code layout

Mitigates *just-in-time* attacks via leaked code pointers.

Denemiark	vaiiiia		րո	ритлош	ритхоштери	pu+x0m+cpn+ss-no-sw0
Apache	30131.13	req/s	1.99%	7.93%	31.34%	58.58%
Kbuild	56.93	sec	1.48%	$\sim 0\%$	2.89%	7.20%
GnuPG	15.30	sec	$\sim 0\%$	$\sim 0\%$	1.18%	3.91%
OpenSSL	3814.23	sign/s	$\sim 0\%$	$\sim 0\%$	$\sim 0\%$	$\sim 0\%$
PyBench	1789	msec	$\sim 0\%$	$\sim 0\%$	$\sim 0\%$	$\sim 0\%$
PHPBench	477859	(score)	$\sim 0\%$	$\sim 0\%$	$\sim 0\%$	$\sim 0\%$
PostMark	5210	trans/s	8.91%	8.29%	19.63%	56.90%
SQLite	549.33	sec	$3.87\%^{*}$	$10.45\%^{*}$	5.66%	3.03%
Redis	2.16M	gets/s	4.39%	4.62%	6.19%	20.72%
Nginx	34193.45	req/s	7.28%	9.33%	28.70%	56.09%
Memcached	106973.37	gets/s	9.72%	7.33%	19.60%	49.97%
Geomean			0.92%	0.83%	1.58%	3.92%

\* Indicates that the relative standard deviation in performance among test runs is between 3.5% and 12.8%.

TABLE II: IskiOS runtime overhead (% over vanilla Linux) on the Phoronix Test Suite.

vanilla	pti+xom		pti+xom+cph		pti+xom+cph+ss	
22.55 MB	$\sim 22.55 MB$	$\sim 0\%$	88.04 MB	292%	90.02 MB	299%

TABLE III: IskiOS code size for different configurations and overhead over vanilla Linux.

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### The 4<sup>th</sup> NSF Secure and Trustworthy Cyberspace Principal Investigator Meeting (2019 SaTC PI Meeting)

foo()

....

jmpq 1f

.....

jmpq 2f

callq bar

jmpq 3f

.....

....

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