

# KSplit: Automating Device Driver Isolation [1]

CNS-1801534: Threat-Aware Defenses - Trent Jaeger (Penn State), Gang Tan (Penn State), Mathias Payer (Purdue/EPFL), Dongyan Xu (Purdue)  
 CNS-1816282: Information Flow Control Infrastructure - Trent Jaeger (Penn State), Danfeng Zhang (Penn State)



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## Introduction

- Device drivers have long been and continue to be a major source of defects and vulnerabilities in modern kernels.
- Previous works on isolating device drivers: (1) significant manual effort and (2) high runtime overhead.
- Recently, some hardware features for efficient isolation have become available (e.g., vmfunc). These techniques significantly reduce the overhead of isolation. [2]
- However, isolating drivers remains hard because of the manual effort required to retrofit the code.

## Motivation and Objective

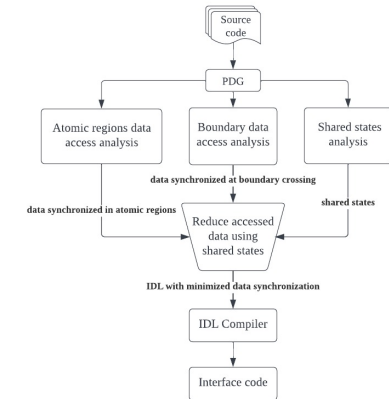
**Motivation:** Reduce the manual work necessary for isolating device drivers as much as possible.

**Objective:** Automate most of key tasks of driver isolation using static analysis techniques and produce warnings for developers to resolve the remaining tasks.

## Main Challenges

- Minimize the data that need to be synchronized across isolation boundary at **cross-domain calls and returns**.
- Correctly handle data synchronization for **kernel concurrency primitives** such as spin\_lock while minimizing the amount of synchronized data.
- Correctly handle data synchronization in the presence of challenging kernel and C language idioms (e.g., pointers to complex struct hierarchies).

## System Workflow

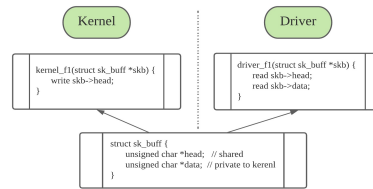


## Compute Shared States

**Shared state:** Shared states are the data structure fields that are accessed by both driver and kernel through the same structure type. This information helps **limit** the amount of data that needs to be synchronized between isolated domains.

### Steps:

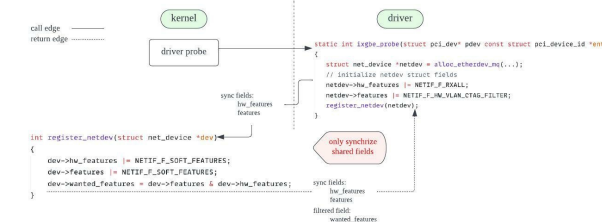
- Compute a set of data structures accessed by both kernel and driver.
- Identify all variables on both sides that match any of the shared data structure types, and analyze the fields accessed through these variables on each side.
- Take the intersection between the accessed fields on both sides to obtain the shared accessed fields.



Compute shared states

## Cross-domain Call Data Synchronization

- Data synchronization for cross-domain calls:** Compute data that needs to be synchronized at domain crossing calls and returns.
- Data access analysis:** For each parameter passed across isolation boundary, use PDG to track the accesses to the parameter.
  - All the data **read** through the parameter during call processing is synchronized at the cross-domain call **invocation**.
  - All the data **modified** through the parameter during call processing is synchronized at the cross-domain call **return**.
- Minimize synchronized data using shared state:** only the shared state is synchronized between the driver and kernel to minimize the overhead of cross-domain calls.



Compute and minimize data synchronized across isolation boundary

## Concurrency Primitives Data Synchronization

- Identify **atomic regions**
  - Find atomic primitives in the PDG (e.g., spin\_lock, mutex\_lock).
  - Use **control flow** in PDG to compute the code within critical sections.
- For each atomic region
  - Compute data **read** within the atomic region and synchronize the data from the other domain **after acquiring** the lock.
  - Compute data **modified** within the atomic region and synchronize the data to the other domain **before releasing** the lock.

```

    struct foo {
        int a; // shared field
        float b; // shared field
    };

    // func is kernel function
    void func(struct foo* f) {
        spin_lock(f->lock);
        read f->a;
        write f->b;
        spin_unlock(f->lock);
    }
    
```

Data synchronization example for critical section

## Evaluation

- KSplit reduces data synchronization by ~30% relative to prior work
- KSplit reduces the manual effort for IDL changes to <60 LOC for the ten drivers isolated and provides concrete warnings for these cases

component	count	number	type	ax	conv	sh_state	mlnet	vmfunc	ax	vmfunc
SLOC	563	194	29K	3K	615	2K	600	54	218	10K
Drv-kern	21	13	132	62	36	14	36	3	15	62
Kern-drv	2	11	85	26	17	1	9	2	5	27
Function	643	1K	5K	3K	512	1K	133	459	1K	

component	count	number	type	ax	conv	sh_state	mlnet	vmfunc	ax	vmfunc
SLOC	1047	2535	13302	896	556	471	1340			
Drv-kern	11	60	25	18	10	14	16			
Kern-drv	10	16	47	4	5	3	13			

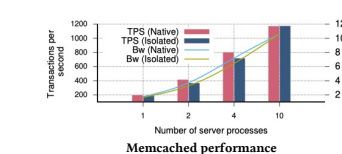
  

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Experiments on 10 automatic isolated drivers



## Conclusion

- Commodity CPUs are converging on a set of practical hardware mechanisms capable of providing support for low-overhead isolation
- The complexity of driver isolation becomes the main challenge for enabling isolation in commodity systems.
- KSplit takes a step forward by enabling isolation of unmodified device drivers in the Linux kernel.

## References

- Huang, Y., Narayanan, V., Detweiler, D., Huang, K., Tan, G., Jaeger, T., & Burtsev, A. (2022, July). KSPLIT: Automating Device Driver Isolation. In *Proceedings of the 16th USENIX Symposium on Operating Systems Design and Implementation (OSDI/22)*. Conditionally accepted.
- Narayanan, V., Huang, Y., Tan, G., Jaeger, T., & Burtsev, A. (2020, March). Lightweight kernel isolation with virtualization and VM functions. In *Proceedings of the 16th ACM SIGPLAN/SIGOPS International Conference on Virtual Execution Environments (VEE'20)* (pp. 157-171). Awarded **Best Paper** of the conference.