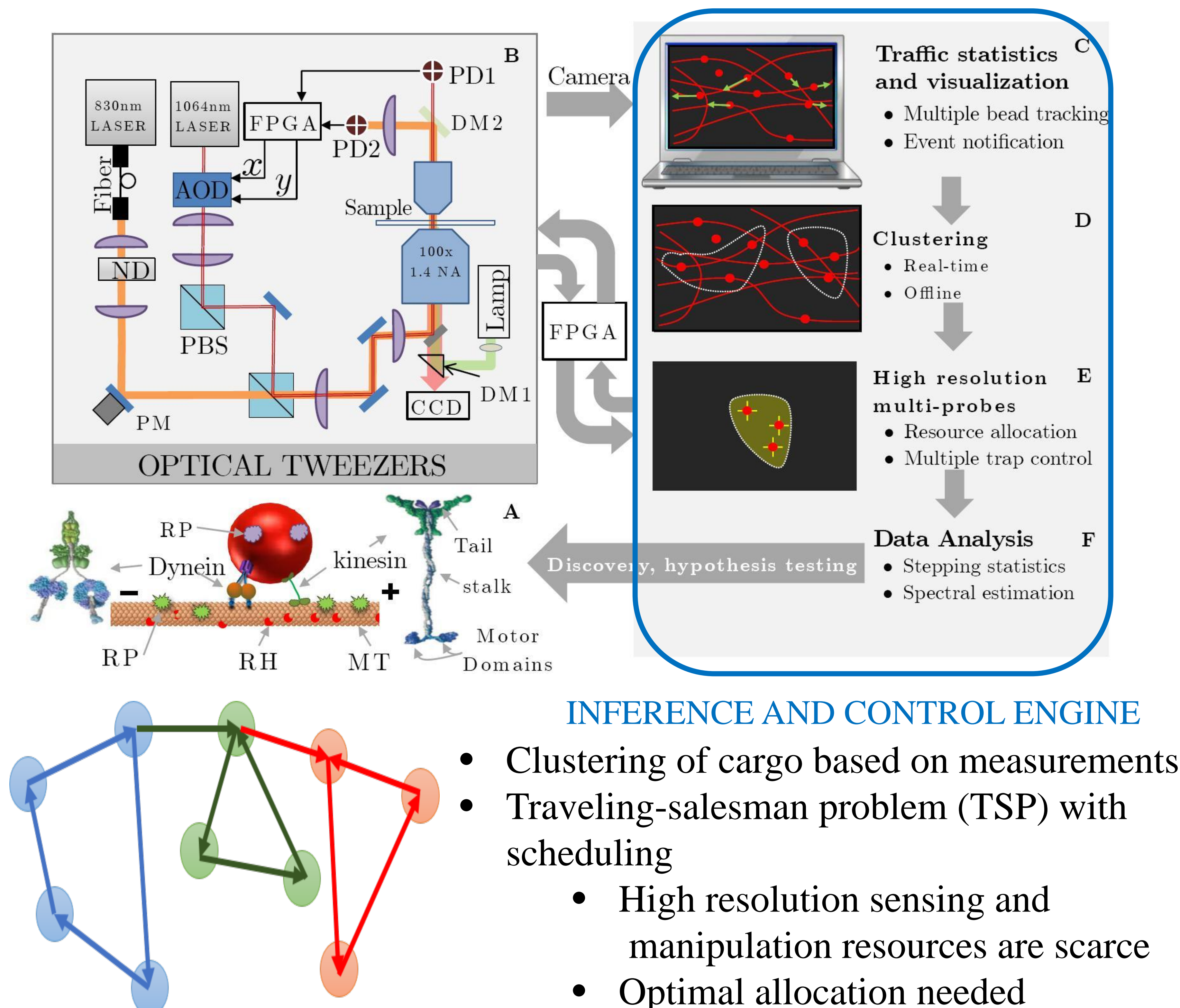
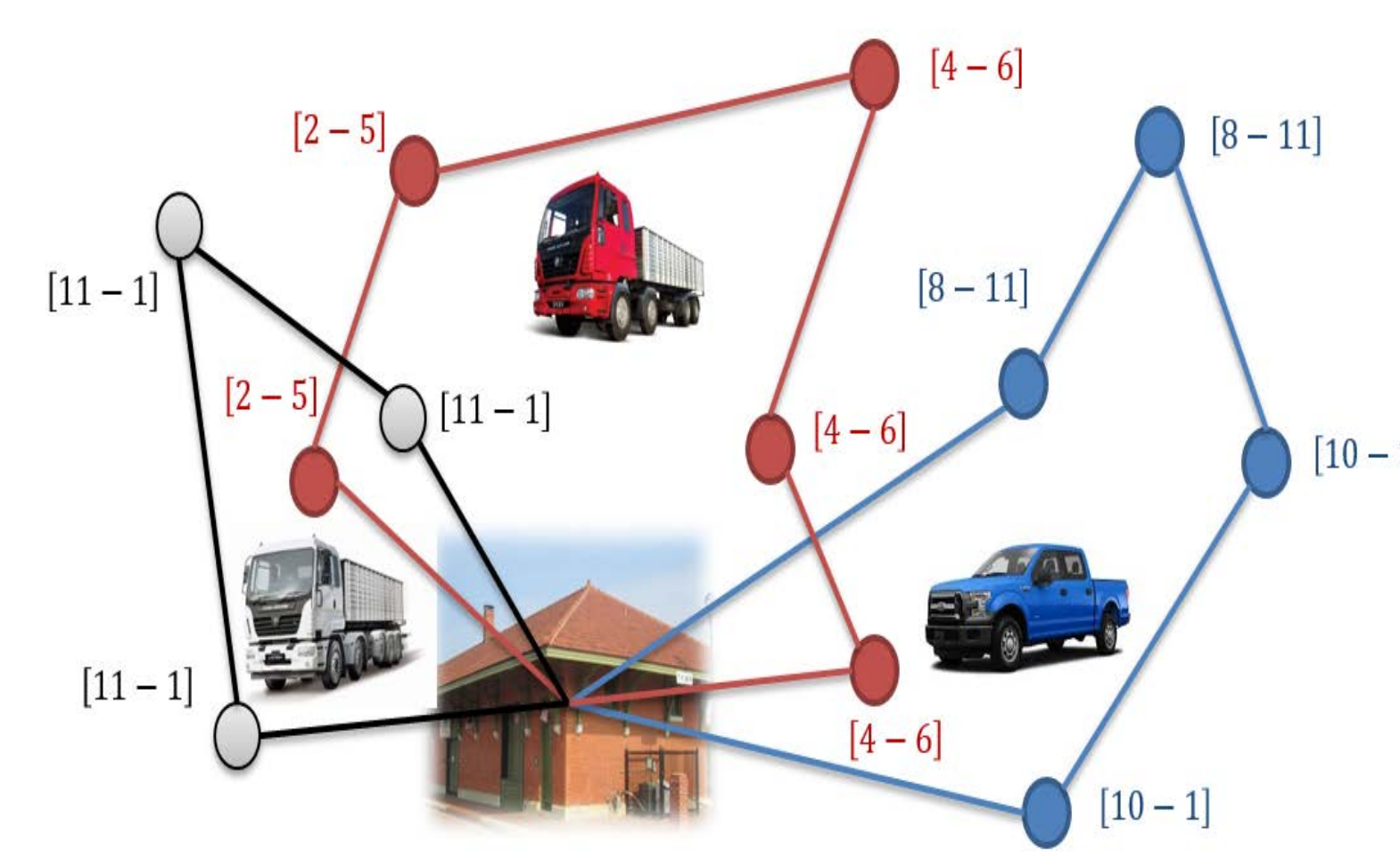


Overview



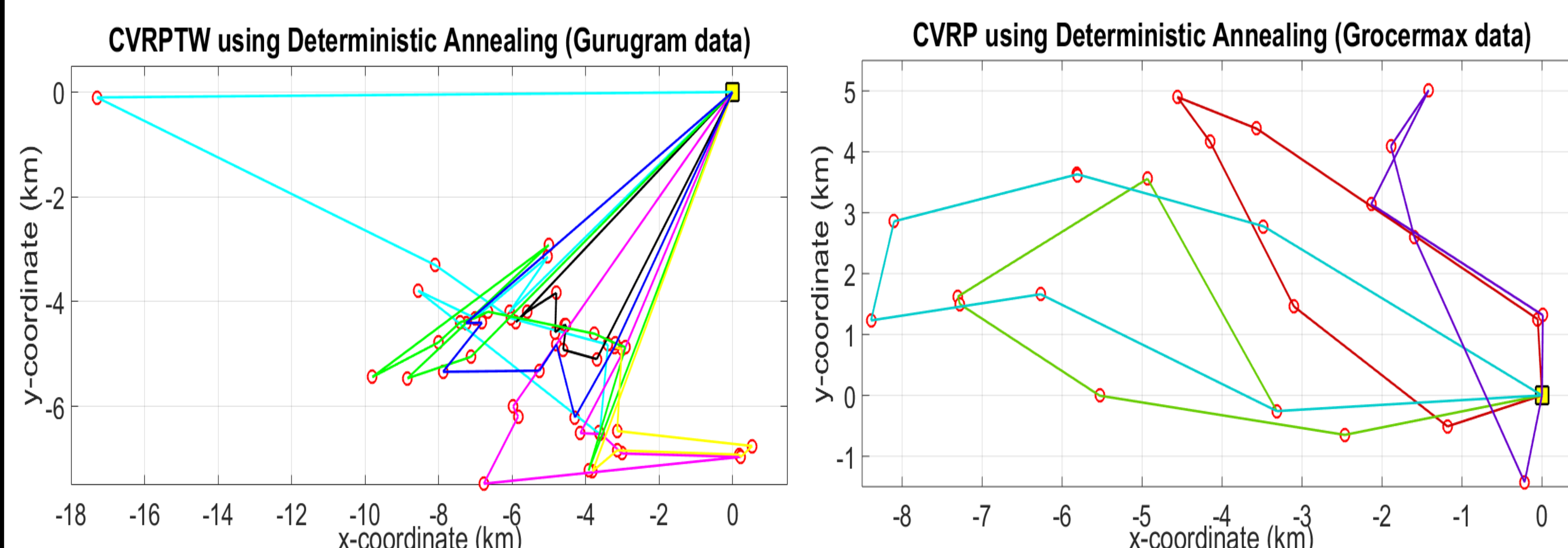
TSP with Time-Windows



- Each customer is equipped with a service time-window $[t_i^s - t_i^f]$
- $d(x_i, y_j)$ in DA is as

$$d(x_i, y_j) = \left\| \begin{array}{c} x_i^{(1)} - y_j^{(1)} \\ x_i^{(2)} - y_j^{(2)} \\ \lambda(0.5(t_i^s + t_i^f) - y_j^3) \end{array} \right\|^2$$

Implementation on Real Datasets from a Logistics Startup



Variants of TSP and Extensions to other related problems

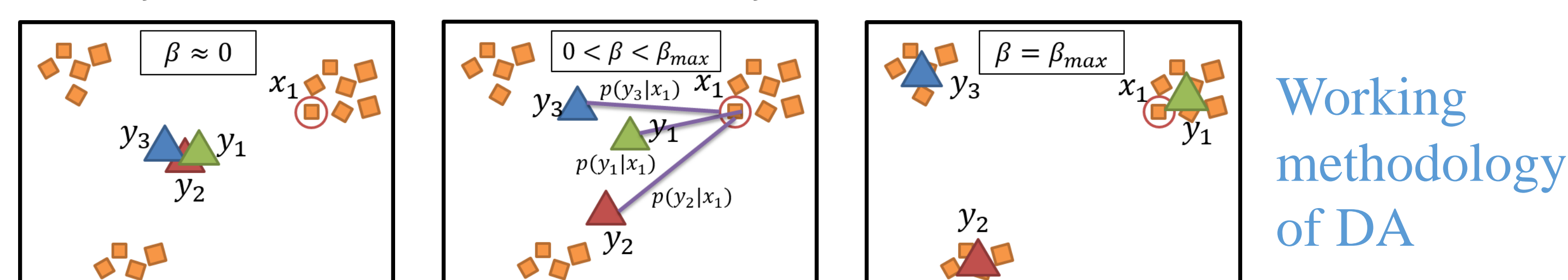
- DA is successfully applied to many NP-hard graph problems

Deterministic Annealing (DA) algorithm

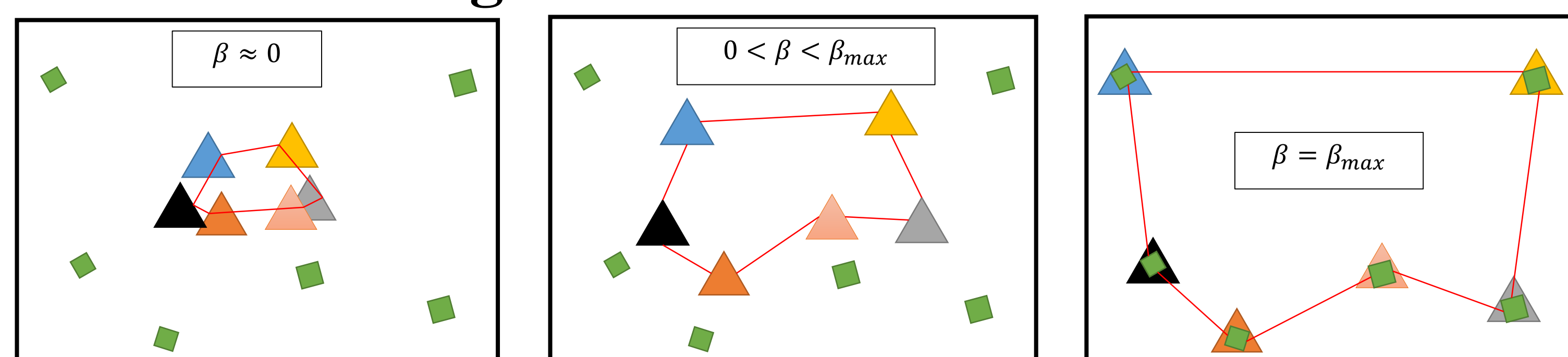
- Given N sites located at $\{x_i; 1 \leq i \leq N\}$, find K facility locations $\{y_j; 1 \leq j \leq K\}$ such that distance of a site i to its nearest facility j is minimized
- DA addresses optimal resource allocation as

$$\min_{\{y_j, p_{ji}\}} D - \frac{1}{\beta} H = \min_{\{y_j\}} -\frac{1}{\beta} \sum_{i=1}^N p_i \log \left(\sum_{j=1}^K e^{-\beta d(x_i, y_j)} \right), \text{ where,}$$

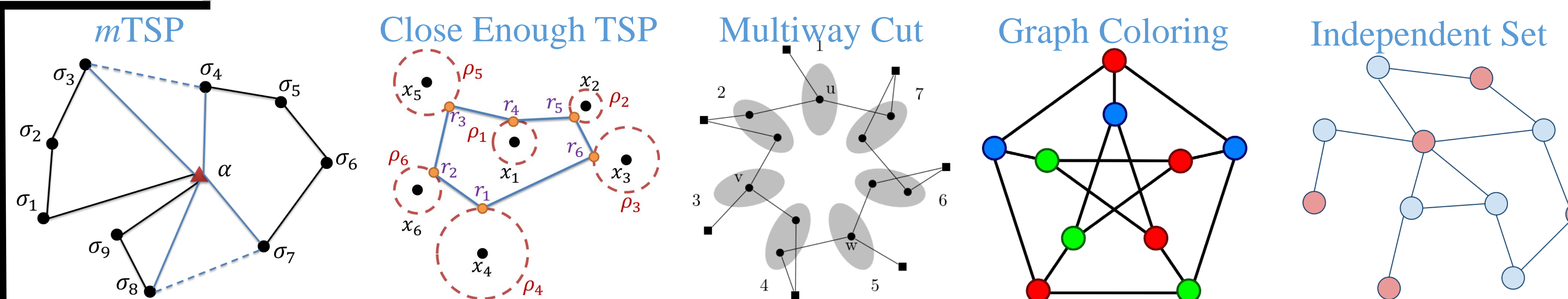
$$D = \sum_{i=1}^N \sum_{j=1}^K p_i p_{ji} d(x_i, y_j); H = -\sum_{i=1}^N p_i \sum_{j=1}^K p_{ji} \log p_{ji}; p_{ji} = \frac{e^{-\beta d(x_i, y_j)}}{\sum_k e^{-\beta d(x_i, y_k)}}; p_i = \frac{1}{N}$$



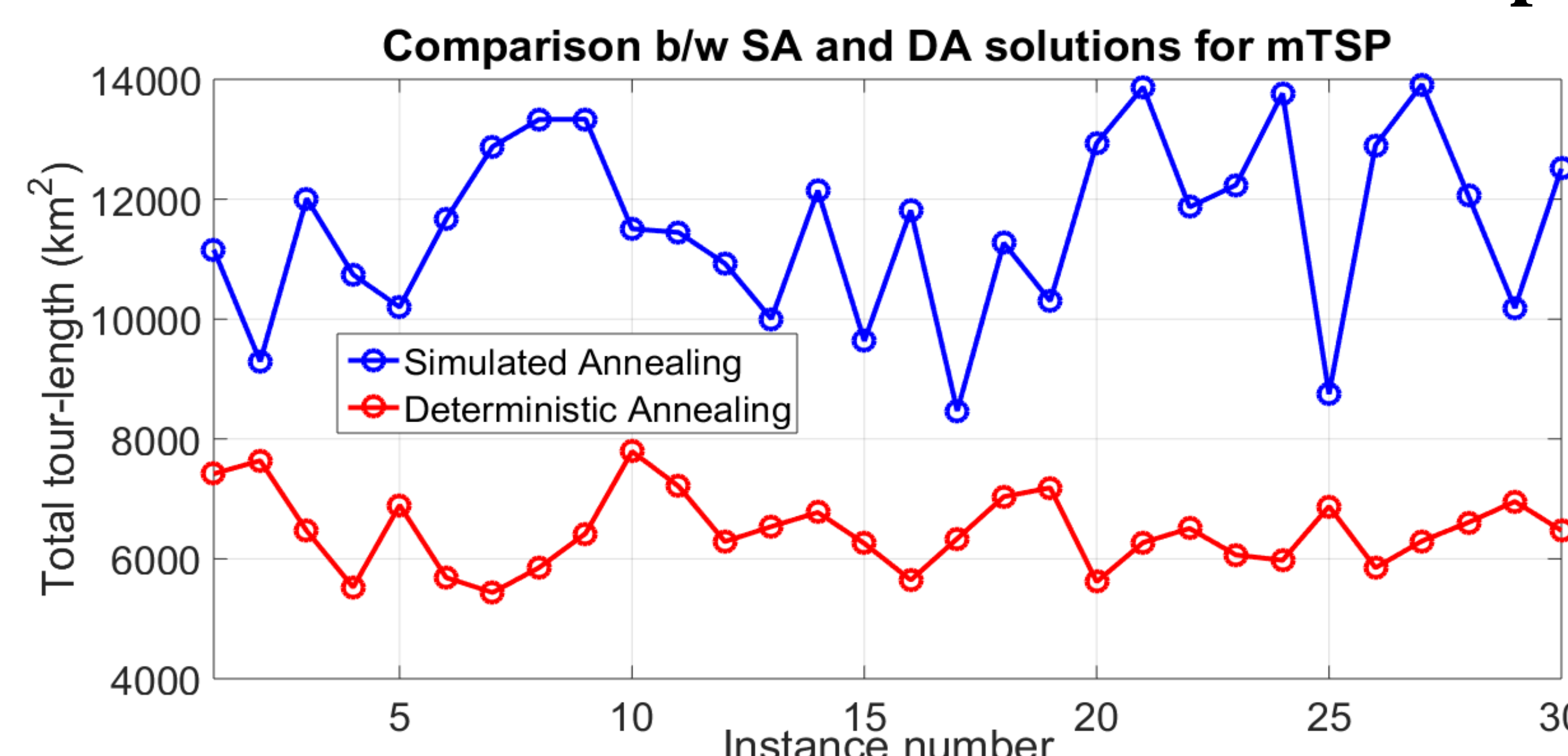
Traveling-Salesman Problem and DA



$$\min_{\{y_j\}} -\frac{1}{\beta} \sum_{i=1}^N \log \left(\sum_{j=1}^N e^{-\beta d(x_i, y_j)} \right) + \left(\sum_{j=1}^N d(y_j, y_{j+1}) \right)$$



Results and Comparisons



- Run-time comparisons:
 - Gurugram data
 - DA: 19.91s
 - SA: 55.40s
 - Grocermax data
 - DA: 5.51s
 - SA: 42.3s

SUMMARY

- Manipulation and sensing using optical fields demand control with resource allocation constraints
- The control problems are studied in the spirit of the TSP with time-windows and its variants thereof
- A Deterministic Annealing (DA) based algorithm is proposed to incorporate several resource allocation constraints