Robust Physical Layer Security with Limited Channel Knowledge



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MIMO Wiretap Channel with ISI Heterogeneity

- freedom (SDoF)?



References

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		Perfect	Delayed	No
	No Secrecy	$\frac{K}{2}$	$> \frac{1}{2}(\sqrt{K} - 1)$	1
	Confidential Messages	$\frac{K(K-1)}{2K-1}$	$> \frac{1}{2}(\sqrt{K} - 6)$	0
 1000	External Eavesdropper	$\frac{K(K-1)}{2K-1}$	$> \frac{1}{2}(\sqrt{K} - 3)$	0
	Confidential Messages and Eavesdropper	$\frac{K(K-1)}{2K-1}$	$> \frac{1}{2}(\sqrt{K} - 6)$	0